Big O Notation

Time complexity of an algorithm

"How much time it takes to run a function as the size of the input grows."

Runtime

array

const array1 = [] number of elements n=5

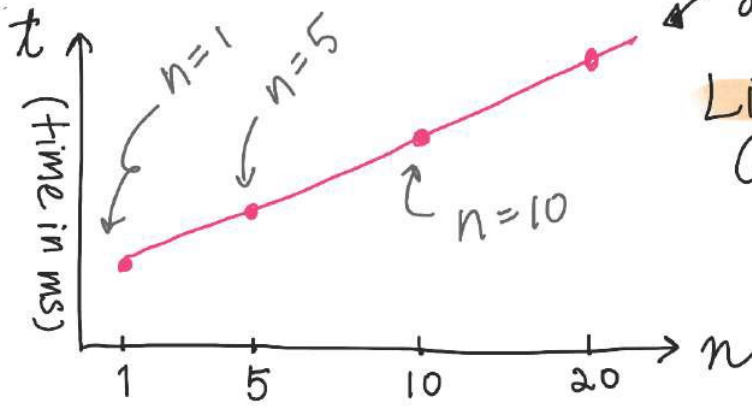
Let's see if there is a needle in the haystack!

JS

```
const numNeedles = (haystack, needle) => {
  let count = 0
  for (let i=0; haystack.length; i++){
    if (haystack[i] === needle) count += 1;
  }
  return count;
}
```

How long does it take to execute when the number of elements (n) is:

♥ execution time grows linearly as array size increases!



this is Linear time Complexity = O(n)

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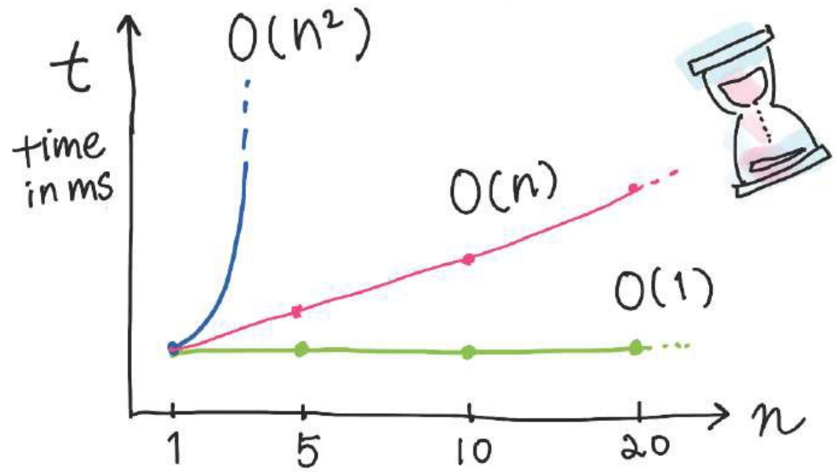
Big O Notation

JS Let's see if we have some function that doesn't actually loop the array:

```
const alwaysTrueNoMatterWhat = (haystack) => {
  return true;
```

n=5
n=10
n=20
⋮

← Array size has no effect on the runtime



☆ Constant time
=
O(1)

☆ Quadratic time = O(n²)

Const array2 = [];

n=5, however the runtime is proportional to n²

```
Const hasDuplicates = (arr) => {
  for (let i = 0; i < arr.length; i++)
    let item = arr[i];
    if(arr.slice(i+1).indexOf(item) !== -1) {
      return true;
    }
  }
  return false;
}
```

① Loop thru the array

② Another array lookup w/ indexOf method

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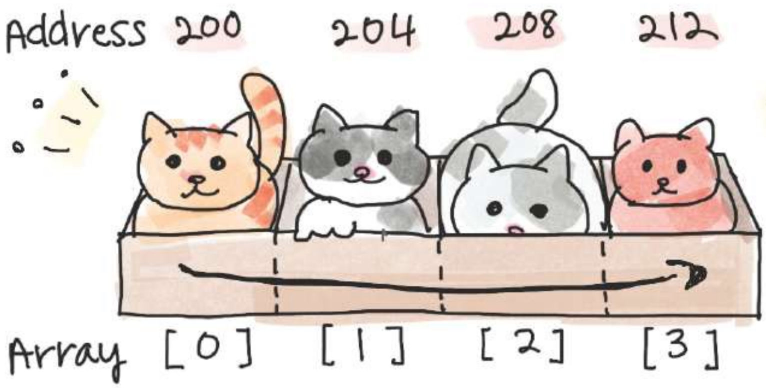
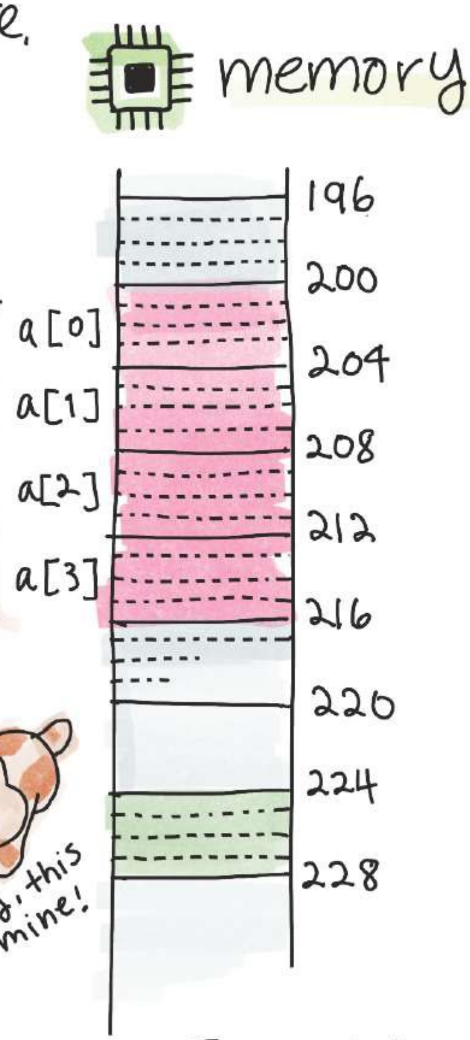
Data Structures
Array & Linked List

Array

a linear data structure, stored in contiguous memory locations.





- Assume each 🐱 is an integer = requires 4 bytes space
- The array of 🐱 must be allocated contiguously!
  → address 200 — 216

yay!

- Can randomly access w/ index
  a[2] → 🐱
- contiguous = no extra memory allocated = no memory overflow

meh!

- fixed size. Large space may not be avail for big array
  🐱 took the space!
- Insert + delete elements are costly.
  → may need to create a new copy of the array + allocate at a new adress.

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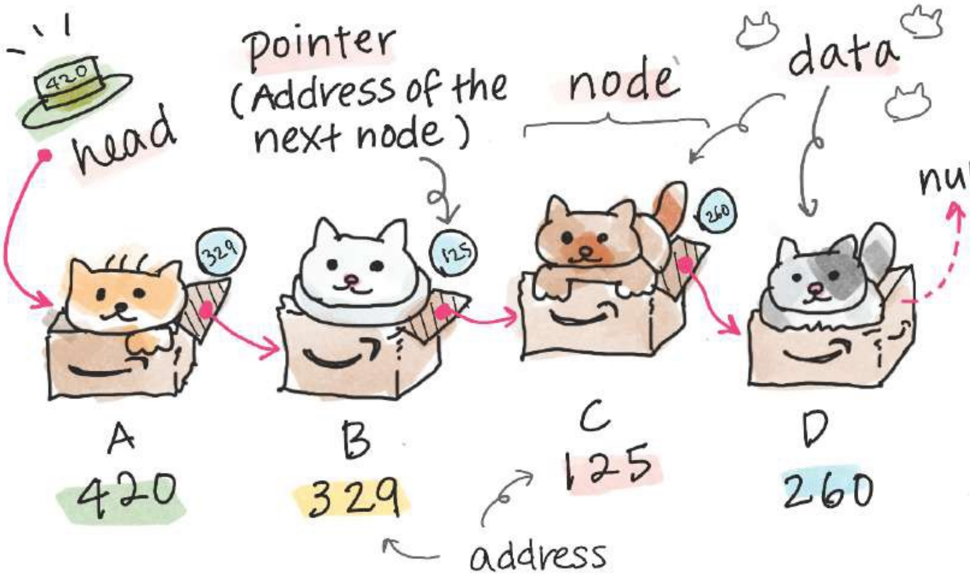
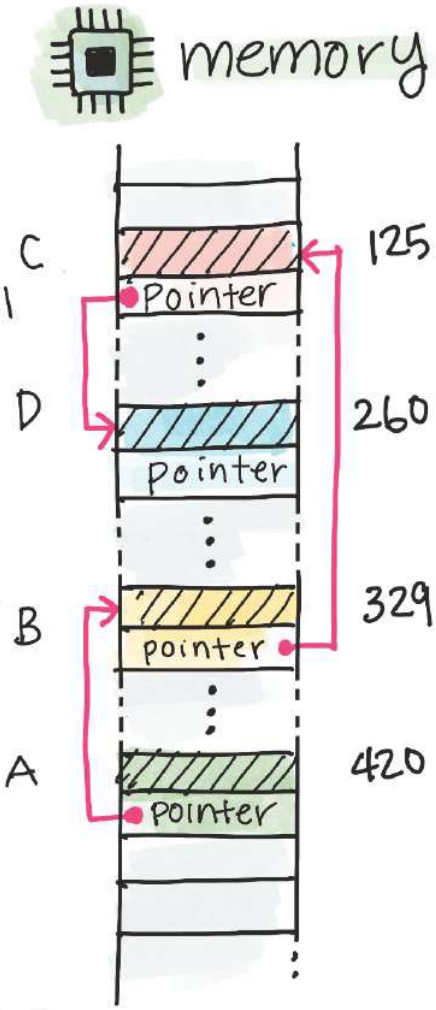
Data Structures

Array & Linked List

Linked list

- ★ a linear data structure
- ★ each element is a separated object & elements are linked w/ pointers





- ★ Unlike an array, LinkedList elements are not stored in contiguous locations.

yay!

- ♥ Dynamic data = size can grow or shrink
- ♥ Insert & delete element are flexible.
  → no need to shift nodes like array insertion
- ♥ memory is allocated at runtime

meh!

- No random access memory.
  → Need to traverse n times
  → time complexity is O(n). array is O(1)
- Reverse traverse is hard

Data Structures

Stack & Queue

LIFO

FIFO

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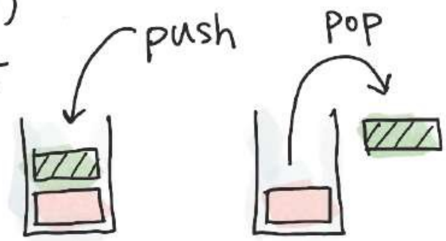
A stack is a LIFO (Last-in-First-out) data structure, where an element added last (=push) gets removed first (=pop)



just like a Stack of ice cream scoope!

omg, the bottom one is always melting !!

Stack as an array in JS

arrays in JavaScript are dynamic!

```
let stack = [ ];                      stack is:
stack.push('mint choc');              // ['mint choc']
stack.push('vanilla');                // ['mint choc', 'vanilla']
stack.push('strawberry');             // ['mint choc', 'vanilla',
                                          'strawberry']
let eaten = stack.pop();              // eaten is
                                          'strawberry'
                                       ['mint choc', 'vanilla']
```

Time complexity is O(1) for both pop + push.

Data Structures

Stack & Queue

LIFO FIFO

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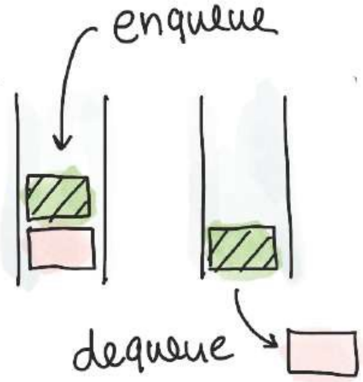
A queue is a FIFO (First-in-First-out) data structure, where an element added first (= enqueue) gets removed first (= dequeue)



just like waiting in line at a popular restaurant!

Stack as an array in JS

wrong! if you queue.unshift('bad kitty'), instead of push(), then the cat cuts in to the front of line!

```
let queue = [ ];                                  queue is:
queue.push('Simba');   // ['simba']
queue.push('Nyan');    // ['simba', 'nyan']
queue.push('maru');    // ['simba', 'nyan', 'maru']
let eater = queue.shift(); // eater is 'Simba'
                                  queue is ['nyan', 'maru']
```

Time Complexity should be O(1) for both enqueue + dequeue but JS shift() is slower!

Data Structures #: Hash Table

- A hash table is used to index large amount of data
- Quick key-value look up. O(1) on average
  - ↳ Faster than brute-force linear search

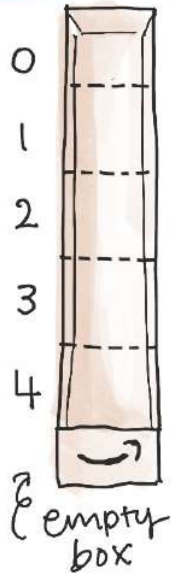
① Let's create an array of size 5.

We're going to add 🐱 data.

key = "Tabby"

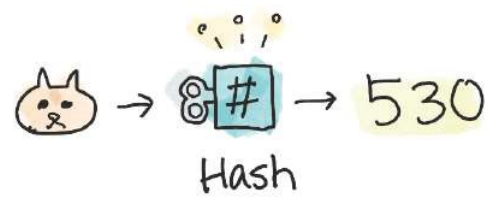
value = "pizza"

↖ some data. Let's say, favorite food!

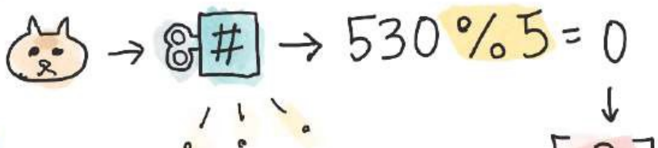


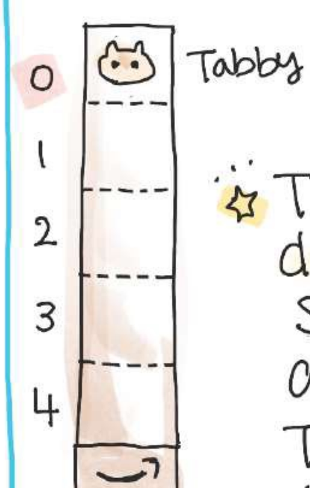
② Calculate the hash value by using the key. "Tabby".

e.g. ASCII code, MD5, SHA1



③ Use modulo to pick a position in the array!

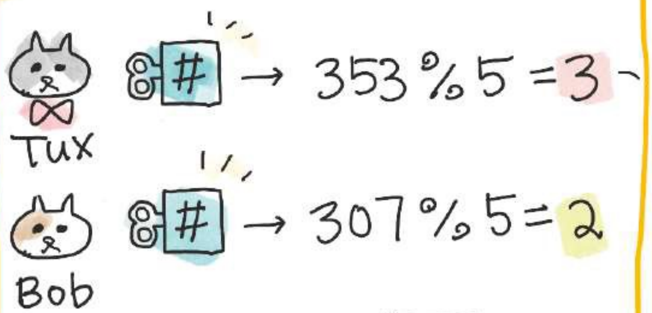




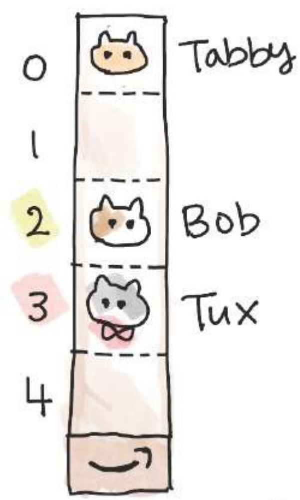
☆ The hash is divided by the size of the array.

The remainder is the position!

④ Let's add more data.



Use the same method to add more 🐱



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Hash Table

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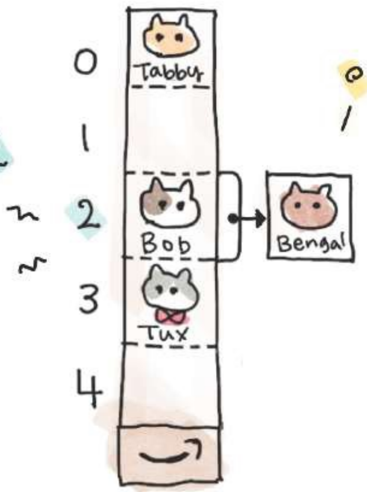
Collision!

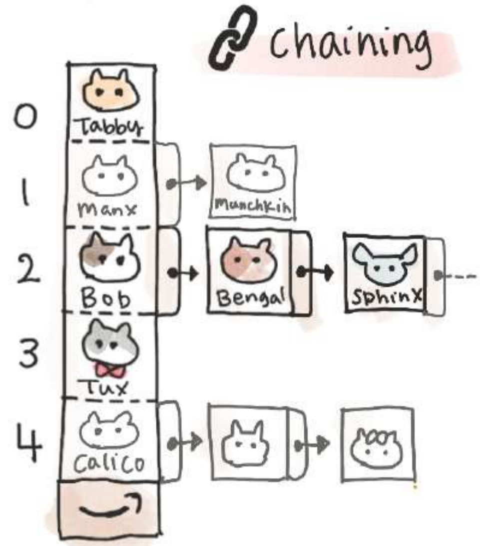
Now we want to add more data.
Let's add "Bengal".



"Bengal" → # → 617 % 5 = 2

But [2] slot has been taken by "Bob" already! = collision!

so let's chain Bengal next to Bob! = chaining



Key: "Bengal" Value: "Dosa"

"Sphinx" "Fish + Chips"

Keep adding data

Searching for data

☆ Let's look up the value for "Bob"

① Get the hash → 307

② Get the index → 307 % 5 = 2

③ Look up Array [2] → found!

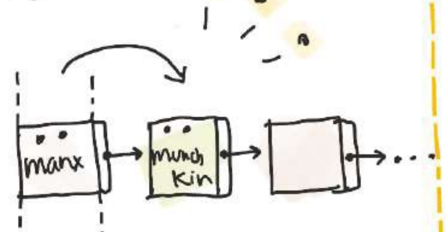
← O(1)

☆ Let's look up "munchkin"

① Hash → 861

② Index → 861 % 5 = 1

③ Array [1] → "manx"



④ Operate a linear-search to find "munchkin"

↳ Average O(n)

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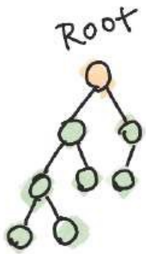
Data Structures: Binary Heap



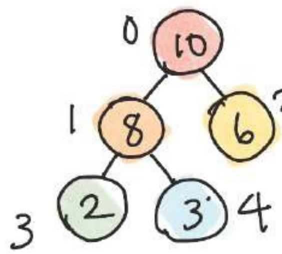
Binary tree → Binary search tree / Binary heap

- tree data structure
- each node has at most 2 children

Binary heap

- Complete tree
- min heap or max heap
- used for priority queue, heap sort etc.

☆ Max heap



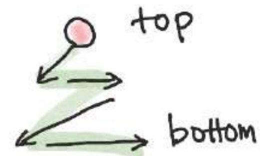
← the root has the largest

(min heap is the opposite!)

in array →

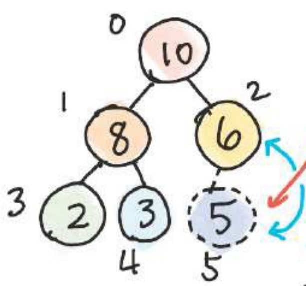
10	8	6	2	3
[0]	[1]	[2]	[3]	[4]

- each node has 0 – 2 children
- always fill top → bottom, left → right



☆ Insertion

Let's add 5 to the heap!



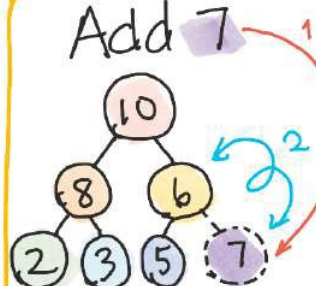
1. add to the next node.
2. compare w/ its parent
3. the parent is greater.

Cool, it's done!

Let's add more!

10	8	6	2	3	5	…
[0]	[1]	[2]	[3]	[4]	[5]	

Add 7

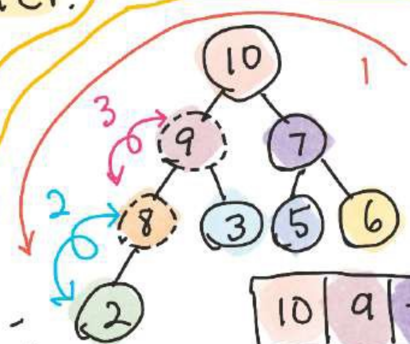


1. Add to the next node
2. Compare w/ parent.

Oh, no!

the parent is smaller than its child! Swap them!!!

10	8	6	2	3	5	7	…



Add to the next node & repeat the process!

10	9	7	8	3	5	6	2	…

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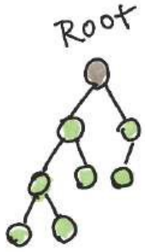
BST

Data Structures

Binary Search Tree



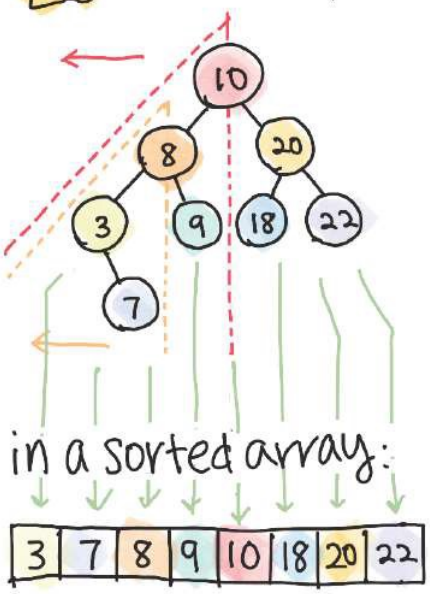
Binary tree → Binary heap / Binary Search Tree

- tree data structure
- each node has at most 2 children

Binary Search Tree

- a.k.a. Ordered or Sorted binary tree
- fast look up
  e.g. phone number lookup table by name

Rule of thumb



★ each value of all nodes in the left subtrees is lesser

△ (10)'s left subtrees: 8, 3, 9, 7

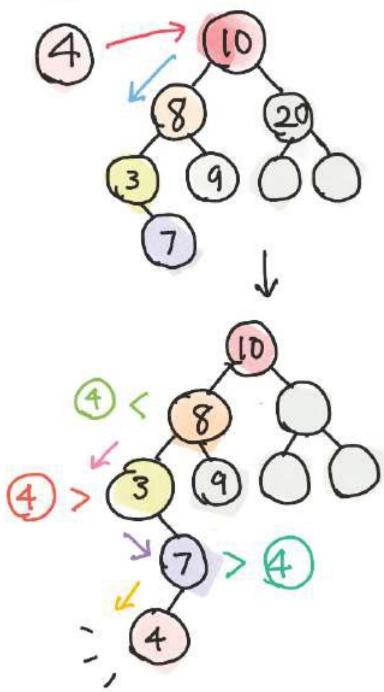
△ (8): 3, 7 ← smaller than parent

★ each value of all nodes in the right subtrees is larger

★ no duplicate values

Insertion

→ Always add to the lowest spot to be a leaf. No rearrange!



Let's add (4)

1. Compare w/ the root first.
2. (4) < (10) so go left.
3. then compare w/ the next, (8)
4. (4) < (8) so go left
5. Compare w/ the (3)
6. (4) > (3) so go right.
7. Compare w/ the (7)
8. (4) < (7), so add to the left! Done.

Complexity:

Ave. $O(\log n)$

Worst. $O(n)$