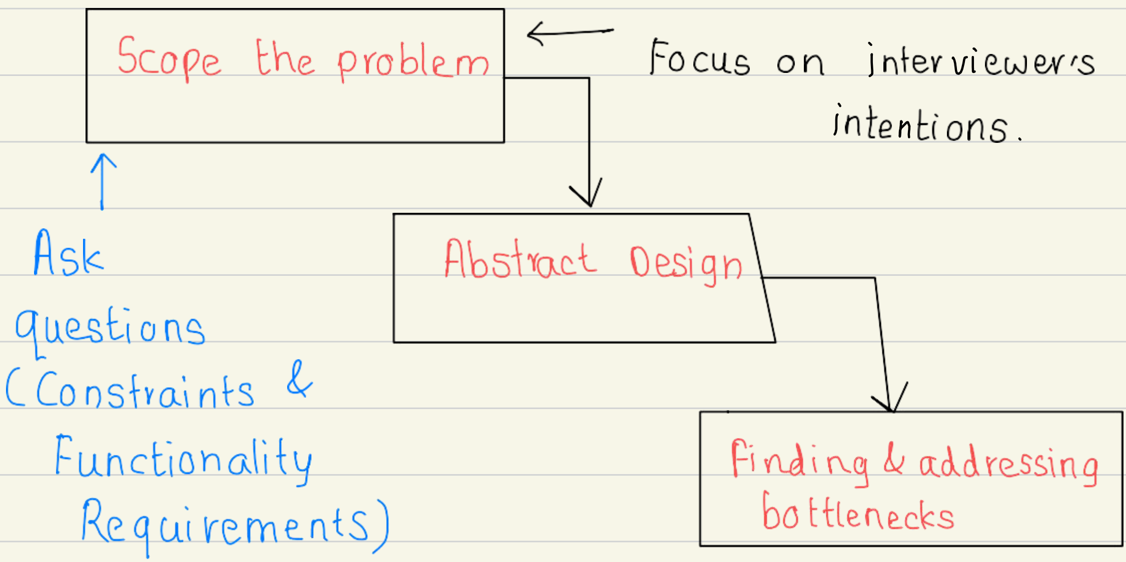
System Design Basics

1) Try to break the problem into simpler modules (Top down approach)

2) Talk about the trade-offs (No solution is perfect)

Calculate the impact on system based on all the constraints and the end test cases.



Rationalize ideas and inputs.

System Design Basics (Contd.)

1) Architectural pieces / resources available
2) How these resources work together
3) Utilization & Tradeoffs

- Consistent Hashing
- CAP Theorem ✓
- Load balancing ✓
- Queues
- Caching ✓
- Replication ✓
- SQL vs No-SQL ✓
- Indexes ✓
- Proxies
- Data Partitioning ✓

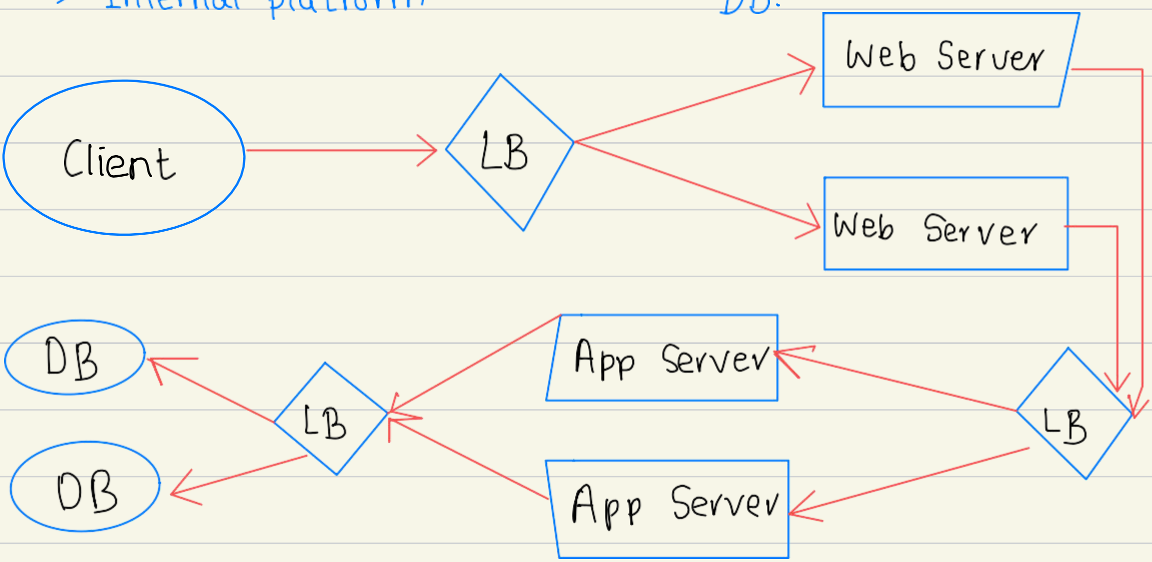
Load Balancing

(Distributed System)

Types of distribution —
- Random
- Round-robin
- Random (weights for memory & CPU cycles)

To utilize full scalability & redundancy, add 3 LB

1) User ⟷(LB1) Web Server
2) Web Server ⟷(LB2) App Server / Cache Server (Internal platform)
3) Internal platform ⟷(LB3) DB.



Smart Clients

Takes a pool of service hosts & balances load.

→ detects hosts that are not responsive

→ recovered hosts

→ addition of new hosts

Load balancing functionality to DB (cache, service)

* Attractive solution for developers (Small scale systems)

As system grows → LBs (Standalone servers)

Hardware Load Balancers:

Expensive but high performance.

e.g. Citrix NetScaler

Not trivial to configure.

Large companies tend to avoid this config. Or use it as 1st point of contact to their System to serve user requests &

Intra network uses Smart clients / hybrid Solution → (Next page) for load balancing traffic.

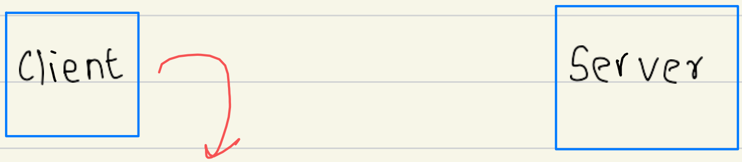
Software Load Balancers

- No pain of creation of smart client
- No cost of purchasing dedicated hardware
  - → hybrid approach

HAProxy ⇒ OSS Load balancer

⇓

1) Running on client machine



(locally bound port)
e.g. localhost : 9000
↑ managed by HAProxy (with efficient management of requests on the port)

2) Running on intermediate server : Proxies running bet^{n} diff. server side components

HAProxy
- manages health checks
- removal & addition of machines
- balances requests a/c pools.

World of Databases

SQL vs. NoSQL

Relational Database

1) Structured
2) Predefined schema
3) Data in rows & columns

Row ⇒ One Entity Info

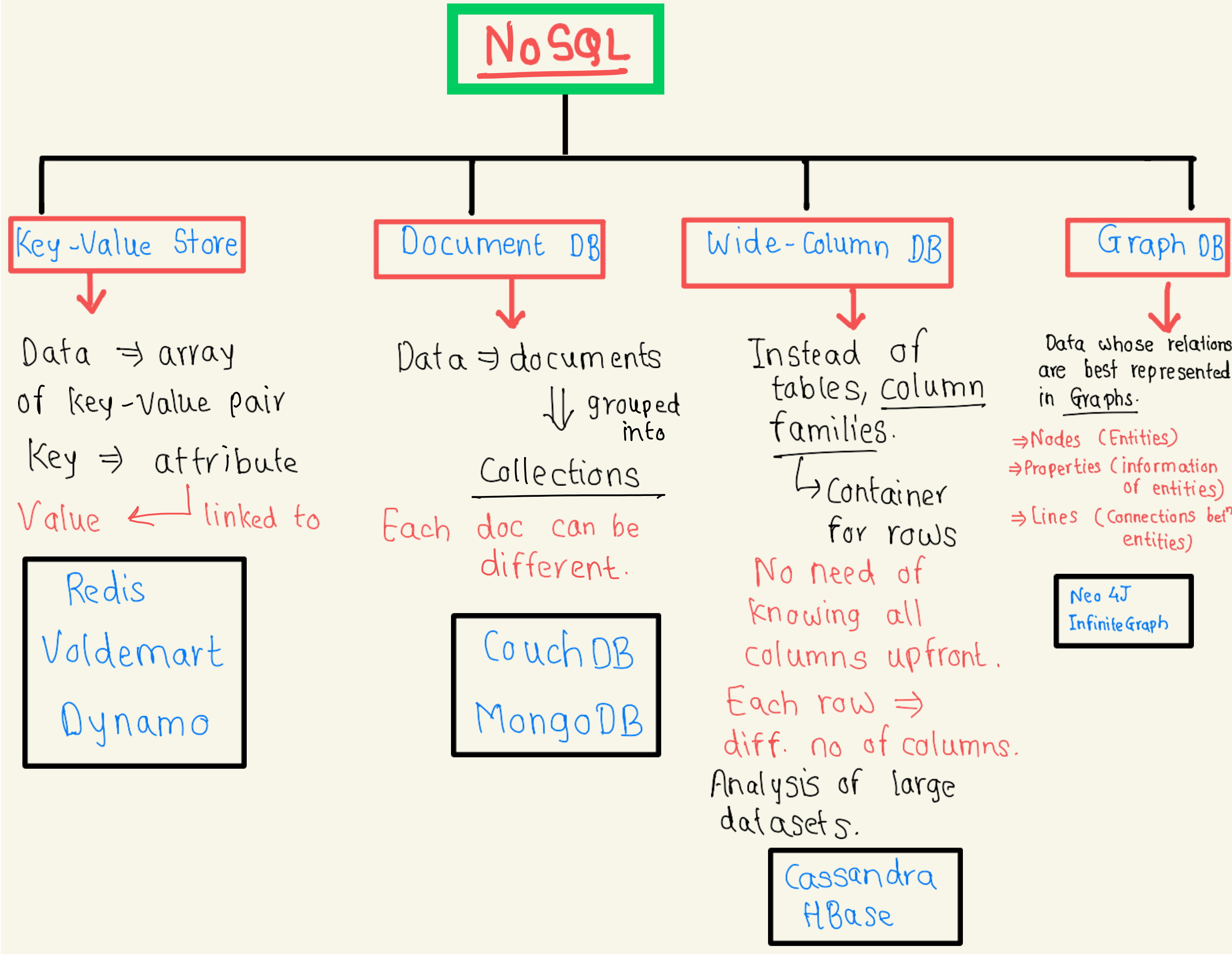
Column ⇒ Separate data points

- MySQL
- Oracle
- MS SQL Server
- SQLite
- Postgres
- MariaDB

Non-relational Database

1) Unstructured,
2) distributed
3) dynamic schema

- Key-Value Stores
- Document DB
- Wide-Column DB
- Graph DB


NoSQL
Key-Value Store
Data ⇒ array of key-value pair
Key ⇒ attribute
Value
linked to
Redis
Voldemart
Dynamo
Document DB
Data ⇒ documents
⇓ grouped into
Collections
Each doc can be different.
CouchDB
MongoDB
Wide-Column DB
Instead of tables, column families.
↳ Container for rows
No need of knowing all columns upfront.
Each row ⇒ diff. no of columns.
Analysis of large datasets.
Cassandra
HBase
Graph DB
Data whose relations are best represented in Graphs.
⇒ Nodes (Entities)
⇒ Properties (information of entities)
⇒ Lines (Connections betⁿ entities)
Neo 4J
InfiniteGraph

High level differences bet^n SQL & NoSQL

Property	SQL	NoSQL
Storage	Tables (Row → Entity, Column → Data point) e.g. Student(Branch, Id, Name)	Diff. data storage models. (Key value, document, graph, columnar)
Schema	fixed Schema (Columns must be decided & chosen before data entry) Can be altered ⇒ modify whole database (need to go offline)	Dynamic Schemas. Columns addition on the fly. Not mandatory for each row to contain data.
Querying	SQL	UnQL (Unstructured query language) Queries focused on collection of documents. Diff. DB ⇒ diff UnQL.
Scalability	Vertically scalable (+ horsepower of h/w) Expensive Possible to scale across multiple servers, ⇒ challenging & time - consuming.	Horizontally scalable. Easy addition of servers. Hosted on cloud or cheap commodity h/w. → Cost effective
Reliability or ACID Compliancy	ACID* Compliant ⇒ Data reliability ⇒ Guarantee of transactions ⇒ Still a better bet.	Sacrifice ACID Compliance for scalability & performance. (ACID - Atomicity, Consistency, Isolation, Durability)

Reasons to use SQL DB

1) You need to ensure ACID compliance:

ACID compliance

⟹ Reduces anomalies

⟹ Protects integrity of the database.

for many E-commerce & financial app^n

→ ACID compliant DB is the first choice.

2) Your data is structured & unchanging.

If your business is not experiencing rapid growth or sudden changes

→ No requirements of more Servers

→ data is consistent

then there's no reason to use system design to support variety of data & high traffic.

Reasons to use NoSQL DB

When all other components of system are fast
→ querying & searching for data ⇒ bottleneck.
NoSQL prevent data from being bottleneck.
Big data ⇒ large success for NoSQL.

1) To store large volumes of data (little/no structure)
No limit on type of data.
Document DB ⇒ Stores all data in one place
(No need of type of data)

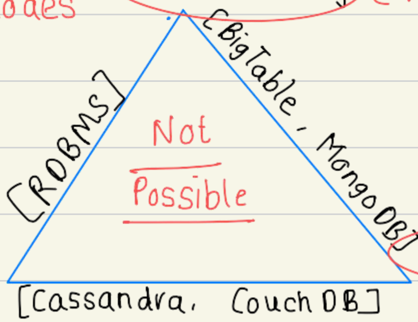
2) Using cloud & storage to the fullest.
Excellent cost saving solution. (Easy spread of data across multiple servers to scale up)
OR commodity h/w on site (affordable, smaller)
⇒ No headache of additional s/w
& NoSQL DBs like Cassandra ⇒ designed to scale across multiple data centers out of the box.

3) Useful for rapid / agile development.
If you're making quick iterations on schema
⇒ SQL will slow you down.

CAP Theorem

Consistency (All nodes see same data at same time)

Achieved by updating several nodes before allowing reads



Availability

⇓

Every request gets response (success / failure)

Achieved by replicating data across different servers

Partition Tolerance

⇓

System continues to work despite message loss/partial failure.

(can sustain any amount of network failure without resulting in failure of entire network)

Data is sufficiently replicated across combination of nodes / networks to keep the system up.

> It is impossible for a distributed system to simultaneously provide more than two of three of the above guarantees.

We cannot build a datastore which is :

1) Continually available
2) Sequentially consistent
3) partition failure tolerant.

Because,

To be consistent ⟹ all nodes should see the same set of updates in the same order

But if network suffers partition,

update in one partition might not make it to other partitions

↳ client reads data from out-of-date partition After having read from up-to-date partition.

Solution: Stop serving requests from out-of-date partition.

↳ Service is no longer 100 % available.

Redundancy & Replication

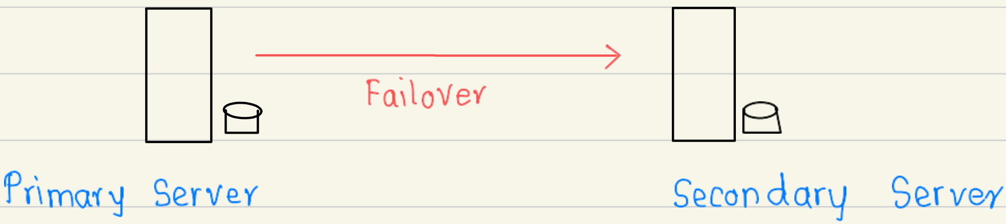
⇒ Duplication of critical data & Services

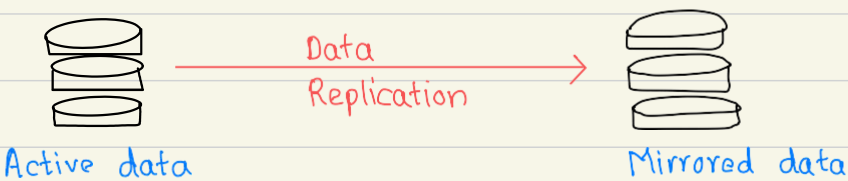
↳ increasing reliability of system.

For critical services & data ⇒ ensure that multiple copies / versions are running simultaneously on different Servers / databases.

⇒ Secure against single node failures.

⇒ Provides backups if needed in crisis.





Service Redundancy: Shared - nothing architecture.

Every node ⇒ independent. No central service managing State. → Helps in Scalability

Helps in Scalability → New servers addition without Special conditions → More resilient to failures → No single point of failure

Caching

Load balancing ⇒ Scales horizontally

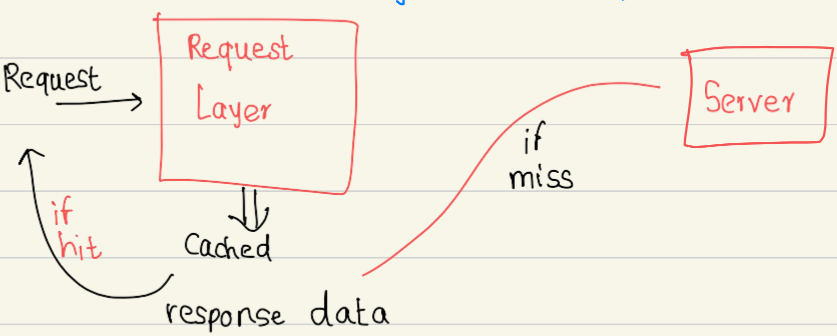
Caching : Locality of reference principle

↑ Used in almost every layer of computing.

1) Application Server Cache:

Placing a cache directly on a request layer node.

↳ Local storage of response



Cache on One Request layer node
can be located

- Memory (very fast)
- Node's local disk (faster than going to network storage)

Bottleneck : If LB distributes requests randomly

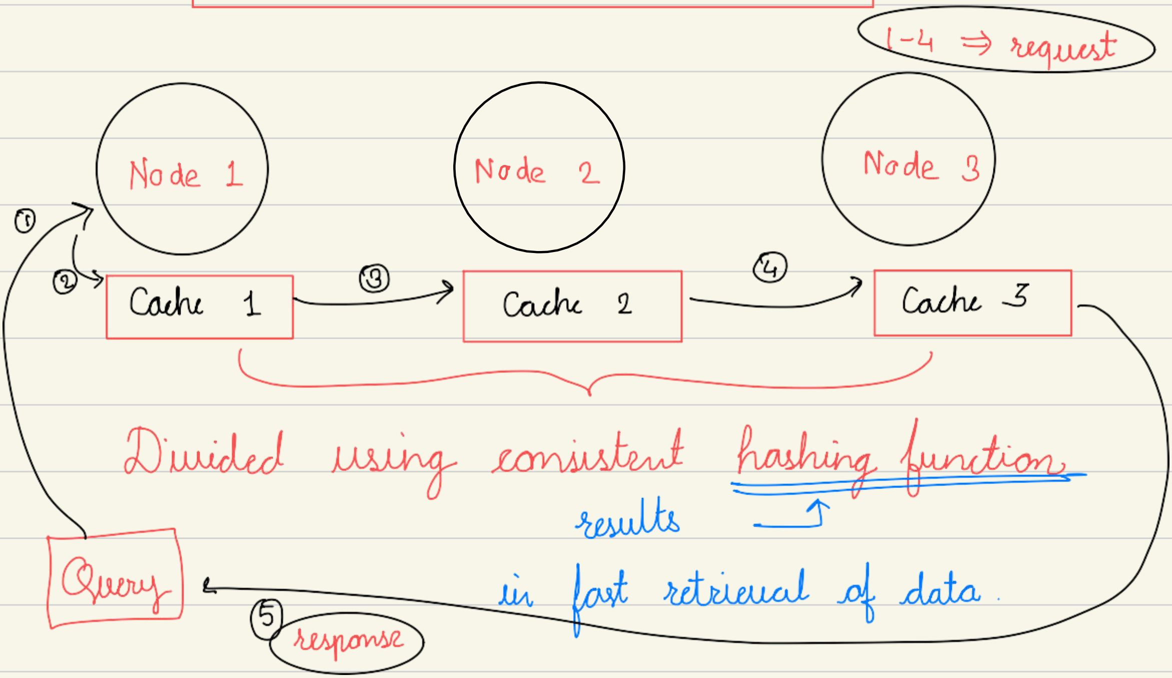
↳ Same request ⇒ different nodes

More cache miss

can be overcomed by

1) Global Caches
2) Distributed Caches

Distributed Cache



Divided using consistent hashing function
results in fast retrieval of data.

Easy to increase cache space by adding more nodes

Disadvantage : Resolving a missing node

storing multiple copies of data on different nodes ← can be handled by

⟶ We're making it more complicated.

Even if node disappears ⇒

request can pull data from Origin.

Global Cache

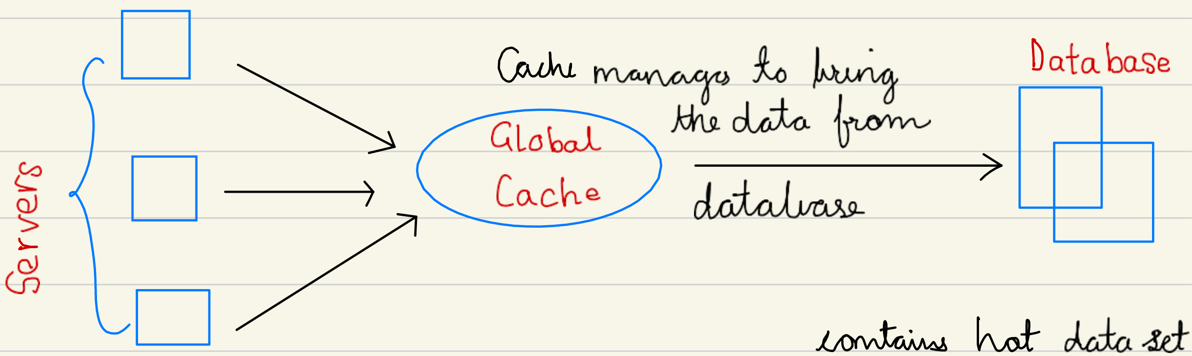
Single cache space for all the nodes.
↳ Adding a cache server / file store (faster than original store)

Difficult to manage if no of clients / request increases.

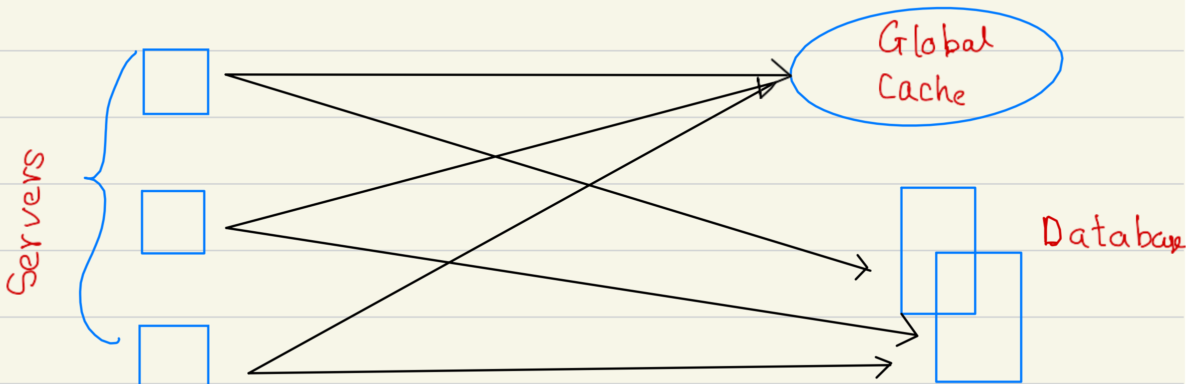
Effective if

1) fixed dataset that needs to be cached
2) special H/w ⇒ fast I/O.

Forms of global cache :



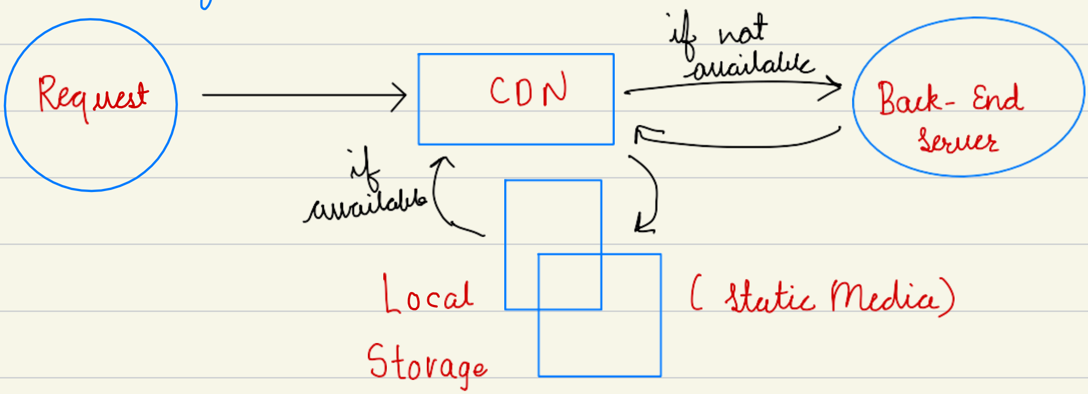
contains hot data set



Appⁿ logic understands the eviction strategy / hot spots better than cache.

CDN : Content Distribution Network

↳ Cache store for Sites that serves large amount of static media.



If the site isn't large enough to have its own CDN

↳ for better & easy future transition

↳ Serve static media using separate subdomain (static . yourservice . com)

↳ using lightweight nginx server

↳ cutover DNS from your server to a CDN later

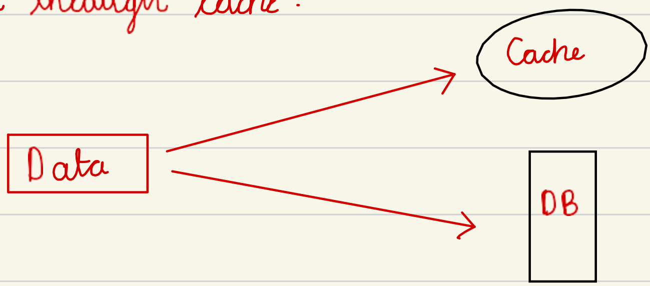
Cache Invalidation

Cached data ⇒ needs to be coherent with the database

If data in DB modified ⇒ invalidate the cached data.

3 schemes :

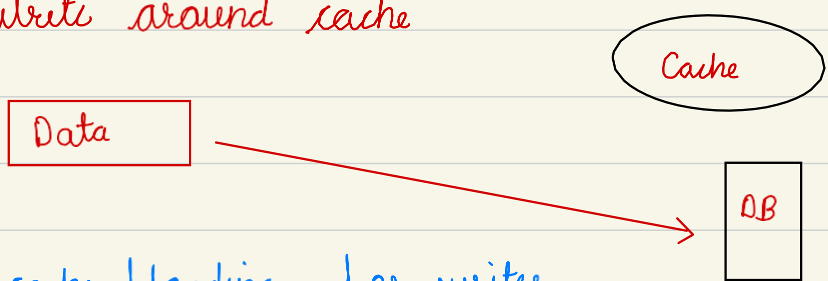
1) Write-through cache :



Data is written same time in both cache & DB.

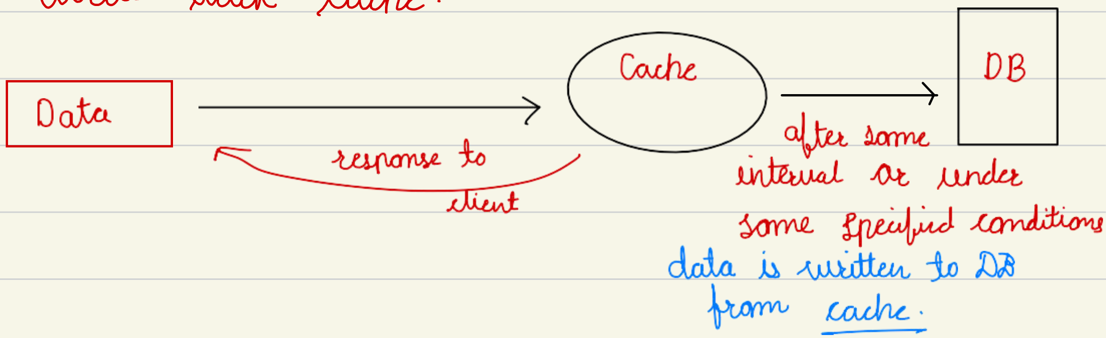
+ Complete data consistency (Cache ≡ DB)
+ Fault tolerance in case of failure (↓↓ data loss)
- high latency in writes ⇒ 2 write operations

2) Write around cache



+ No cache flooding for writes
- read request for newly written data ⇒ Miss → higher latency

3) Write back cache:



+ low latency & high throughput for write-intensive app^n

- Data loss ↑↑ (only one copy in cache)

Cache Eviction Policies

1) FIFO
2) LIFO or FILO
3) LRU
4) MRU
5) LFU
6) Random Replacement

Sharding || Data Partitioning

Data Partitioning : Splitting up DB / table across multiple machines ⇒ managealility, performance, availability & LB

** After a certain scale point, it is cheaper and more feasible to scale horizontally by adding more machines instead of vertical scaling by adding beefier servers.

Methods of Partitioning:

1) Horizontal Partitioning : Different rows into diff. tables

Range based sharding

e.g. storing locations by zip

Table 1 : Zips with < 100000
Table 2 : Zips with > 100000
and so on

⇒ different ranges in different tables

** Cons: if the value of the range not chosen carefully
⇒ leads to unbalanced servers
e.g. Table 1 can have more data than table 2.

Vertical Partitioning

\# Feature wise distribution of data
↳ in different servers.

e.g. Instagram — DB server 1 : user info
— DB server 2 : followers
— DB server 3 : photos

** Straightforward to implement
** low impact on app.

⊖⊖ if app → additional growth
need to partition feature specific DB across various servers
(e.g. it would not be possible for a single server to handle all metadata queries for 10 billion photos by 140 mill. users

Directory based partitioning

⇒ A loosely coupled approach to work around issues mentioned in above two partitionings.

** Create lookup service ⇒ current partitioning scheme & abstracts it away from the DB access code.

Mapping (tuple key → DB server)

Easy to add DB servers or change partitioning scheme.

Partitioning Criteria

1) Key or Hash based partitioning :

Key attr. of the data → Hash function → Partition number

Effectively fixes the total number of servers / partitions

So if we add new server / partition → change in hash function → downtime  of redistribution

→ **Solution : Consistent hashing**

2) List Partitioning : Each partition is assigned a list of values.

new record → Lookup for key → store the record (partition based on the key)

3) Round Robin Partitioning:

uniform data distribution

With 'n' partitions

⇒ the 'i' tuple is assigned to partition (i mod n)

4) Composite Partitioning:

combination of above partitioning schemes

Hashing + List ⇒ Consistent Hashing

⇓

Hash reduces the key space to a size that can be listed.

Common Problems of Sharding:

Sharded DB: Extra constraints on the diff. operations

⇓

operations across multiple tables or multiple rows in the same table → no longer running in single server.

1) Joins & Denormalization :

Joins on tables on single server ⇒ straightforward.

* Not feasible to perform joins on sharded tables

↳ Less efficient (data needs to be compiled from multiple servers)

Workaround ⇒ Denormalize the DB

so that the queries that previously reqd. joins can be performed from a single table.

cons : Perils of denormalization

↳ data inconsistency

2) Referential integrity : Foreign Keys on sharded DB

↳ difficult

* Most of the RDBMS does not support foreign keys on sharded DB.

If appⁿ demands referential integrity on sharded DB

↳ enforce it in appⁿ code (SQL jobs to clean up dangling references)

3) Rebalancing:

Reasons to change sharding scheme:

a) non-uniform distribution (data wise)

b) non-uniform load balancing (request wise)

Workaround: 1) add new DB

2) rebalance

↳ change in partitioning scheme

↳ data movement

↳ downtime

↓

We can use directory-based partitioning

↳ highly complex

↳ single point of failure (lookup service / table)

Indexes

⇒ Well Known because of databases.

⇒ Improves speed of retrieval

- Increased storage overhead

- Slower writes

  ↳ Write the data

    ↳ Update the index

⇒ Can be created using one or more columns

* Rapid random lookups

  & efficient access of ordered records.

Data Structure

Column ⟶ Pointer to whole row

→ Create different views of the same data.

  ↳ very good for filtering / sorting of large data sets.

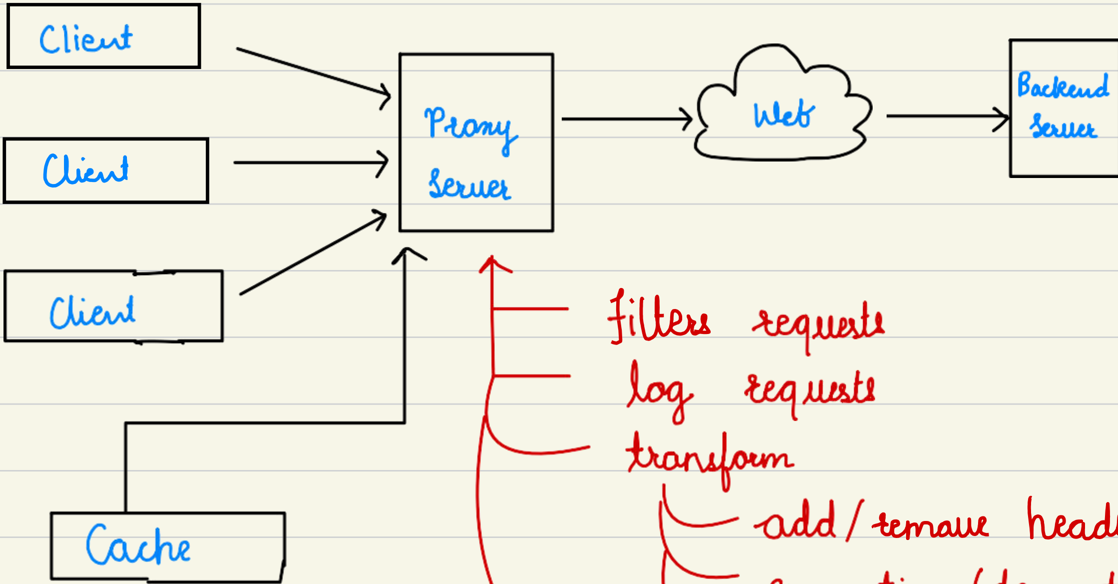
    ↳ no need to Create additional copies.

Used for datasets (TB in size) & small payload (KB)

spred over several physical devices ⟶ We need some way to find the correct physical location i.e. Indexes

Proxies

- useful under high load situations
- if we have limited Caching
  - ↳ batches several requests into one



Cache: frequently used resources

Proxy Server:
- filters requests
- log requests
- transform
  - add / remove headers
  - encryption / decryption
  - compression
- request co-ordination (request traffic optimization)

↑

we can also use spatial locality ← Collapse same data access request into one.

⇒ Collapsed forwarding

↳ collapsing requests for data that is spatially close

↳ minimize reads from origin.

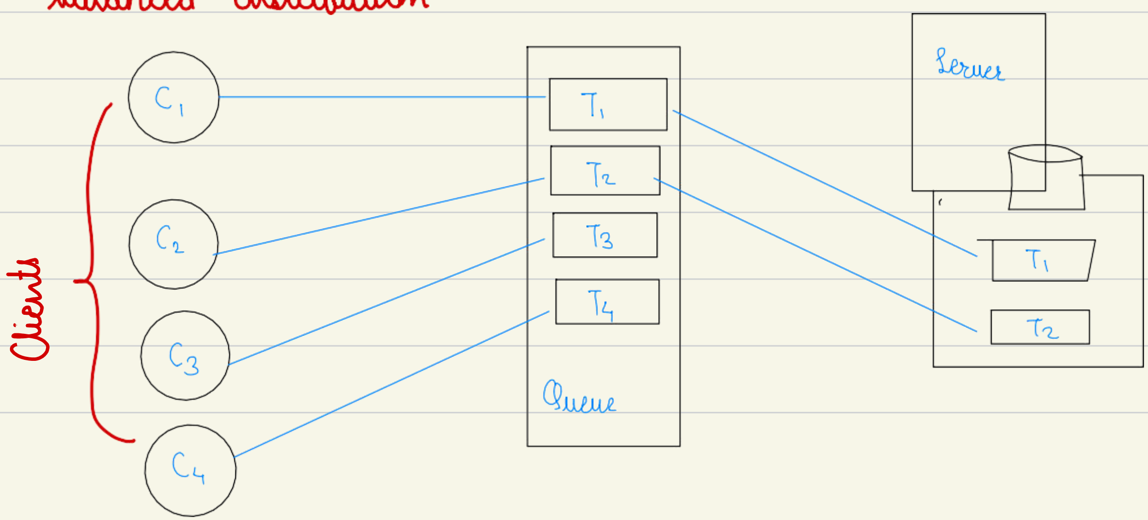
Queues

⇒ Effectively manages requests in large-scale distributed system

→ In small systems → writes are fast.
→ In complex systems → high incoming load & individual writes take more time

* To achieve high performance & availability
  ↳ system needs to be asynchronous
    ↳ **Queues**

Synchronous behaviour → degrades performance
↓
can use Load balancing ← difficult for fair & balanced distribution



Queues : asynchronous communication protocol

↳ Client sends task

↳ gets ACK from queue (receipt)

↑ serves as reference for the results in future

↳ Client continues its work.

Limit on the size of request & number of requests in queue

Queue : Provides fault-tolerance

↳ protection from service outage/failure

↑ highly robust

↳ retry failed service request

↑ Enforces Quality of Service guarantee (Does NOT expose clients to outages)

Queues : distributed communication

↳ Open Source implementations

↳ RabbitMQ, ZeroMQ, ActiveMQ, BeanstalkD.

Consistent Hashing

Distributed Hash Table

index = hash_function (Key)

Suppose we're designing distributed caching system with n cache servers

↳ hash_function ⇒ (Key % n)

Drawbacks:

1) NOT horizontally scalable

↳ addition of new server results in

↳ need to change all existing mapping.

(downtime of system)

2) NOT load balanced

(because of non-uniform distribution of data)

↓

Some caches : hot & saturated

Other caches : idle & empty

How to tackle above problems?

Consistent Hashing

What is consistent Hashing?

→ Very useful strategy for distributed caching & DHTs.

→ minimizes reorganization in scaling up / down.

→ only k/n keys needs to be remapped.

$k \Rightarrow$ total number of keys

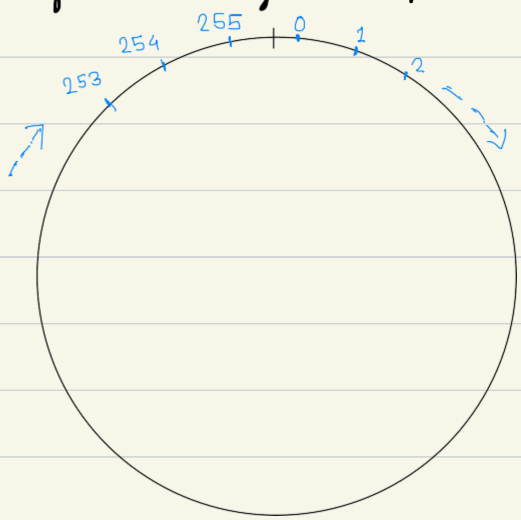
$n \Rightarrow$ number of servers

How it works?

Typical hash function suppose outputs in [0, 256)

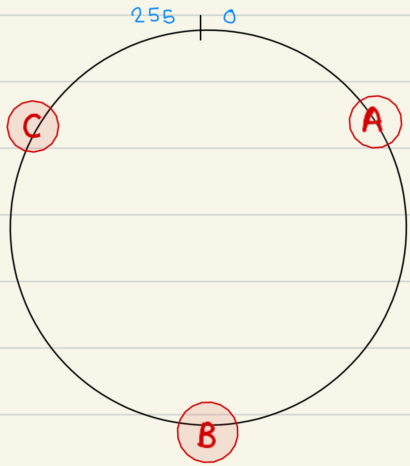
In consistent hashing,

imagine all of these integers are placed on a ring.



& we have 3 servers : A, B & C.

1) Given a list of servers, hash them to integers in the range.

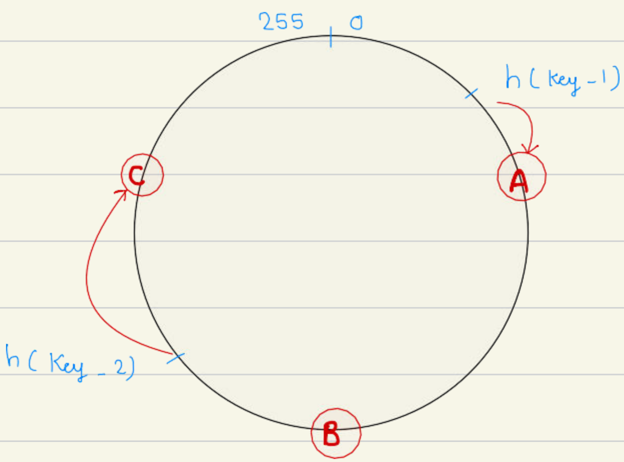


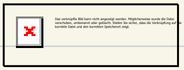
2) Map key to a server:

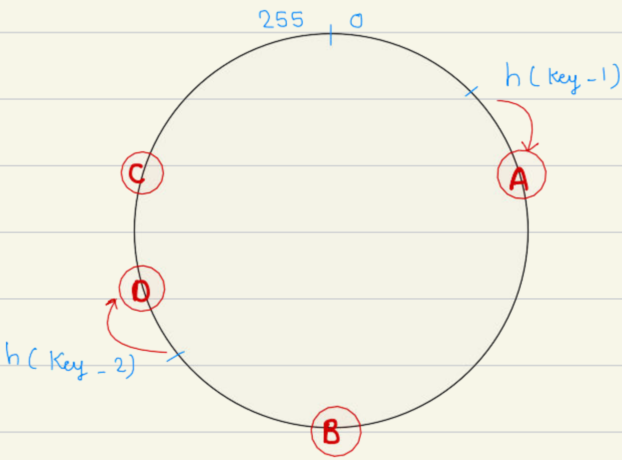
a) Hash it to single integer

b) Move CLK wise until you find server

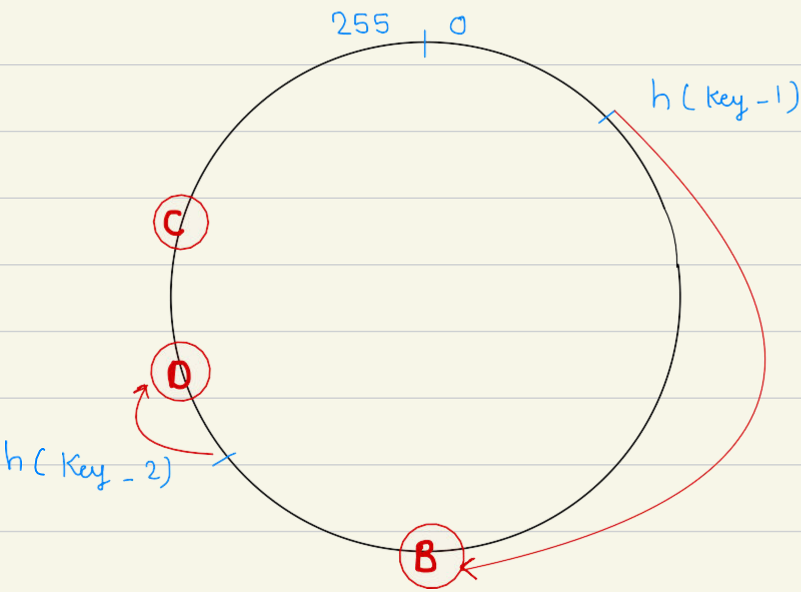
c) map key to that server.



Adding a new server '**D**', will result in moving  to 'D'



Removing server '**A**', will result in moving the 'key_1' to '**B**'



Consider real world scenario

data → randomly distributed

↳ unbalanced caches.

How to handle this issue?

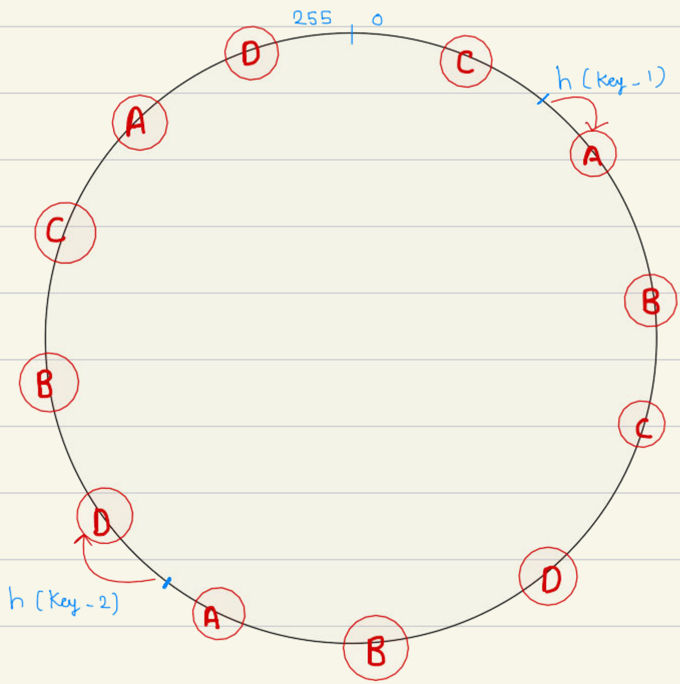
Virtual Replicas

⇒ Instead of mapping each node to a single point we map it to <u>multiple</u> points.

↳ (More number of replicas

↳ more equal distribution

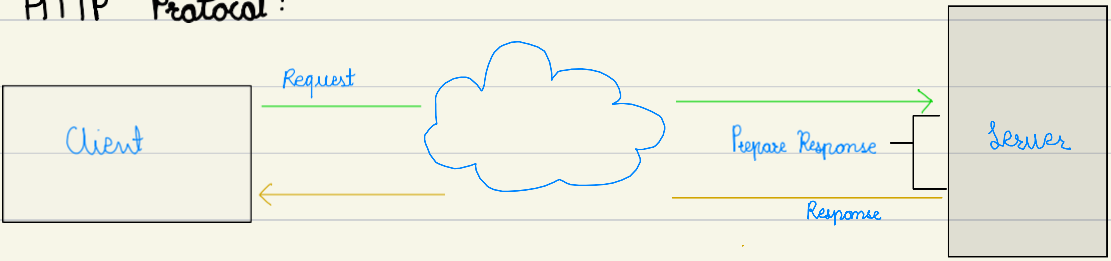
↳ good load balancing)



Long-Polling vs WebSockets vs Server-Sent Events

↳ Client-Server Communication Protocols

HTTP Protocol:



AJAX Polling:

Clients repeatedly polls servers for data

Similar to HTTP protocol

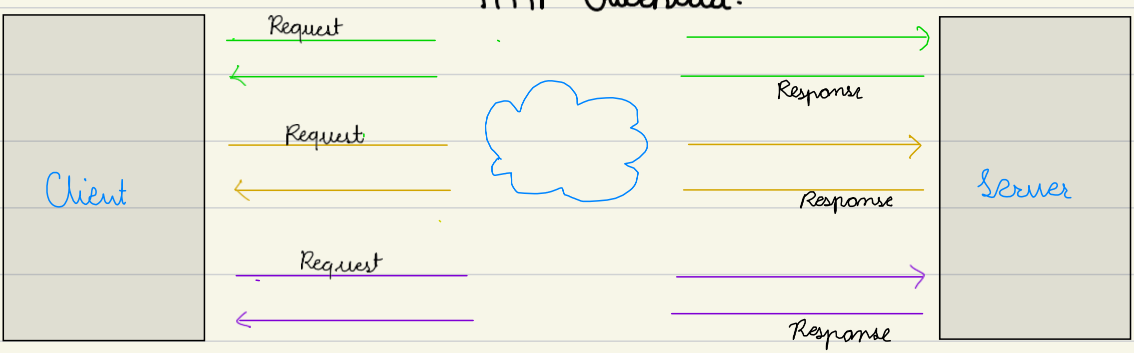
↳ requests sent to server at regular intervals (0.5 sec)

Drawbacks:

Client keeps asking the server new data

↳ Lot of responses are 'empty'

↳ HTTP Overhead.



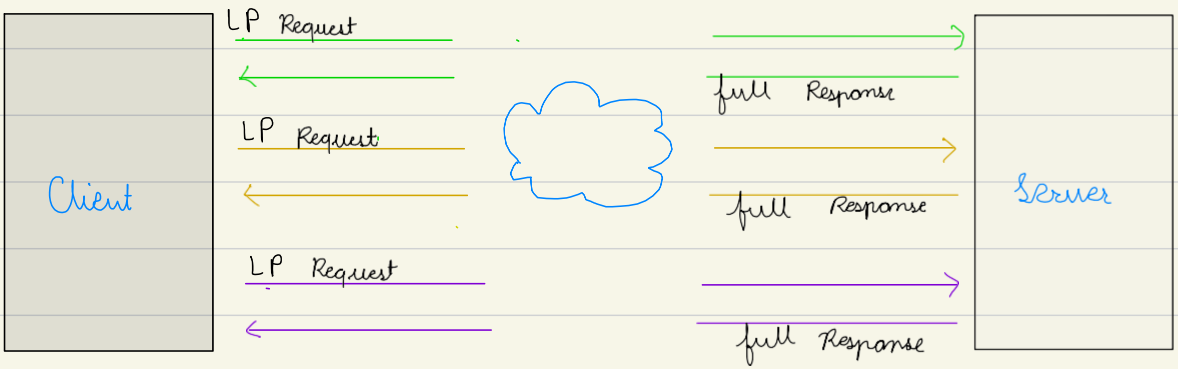
HTTP Long Polling : 'Hanging GET'

Server does NOT send empty response.

Pushes response to clients only when new data is available

1) Client makes HTTP Request & waits for the response.
2) Server delays response until update is available or until timeout occurs.
3) When update → Server sends full response.
4) Client sends new long-poll request
   a) immediately after receiving response
   b) after a pause to allow acceptable latency period
5) Each request has timeout.

Client needs to reconnect periodically due to timeouts

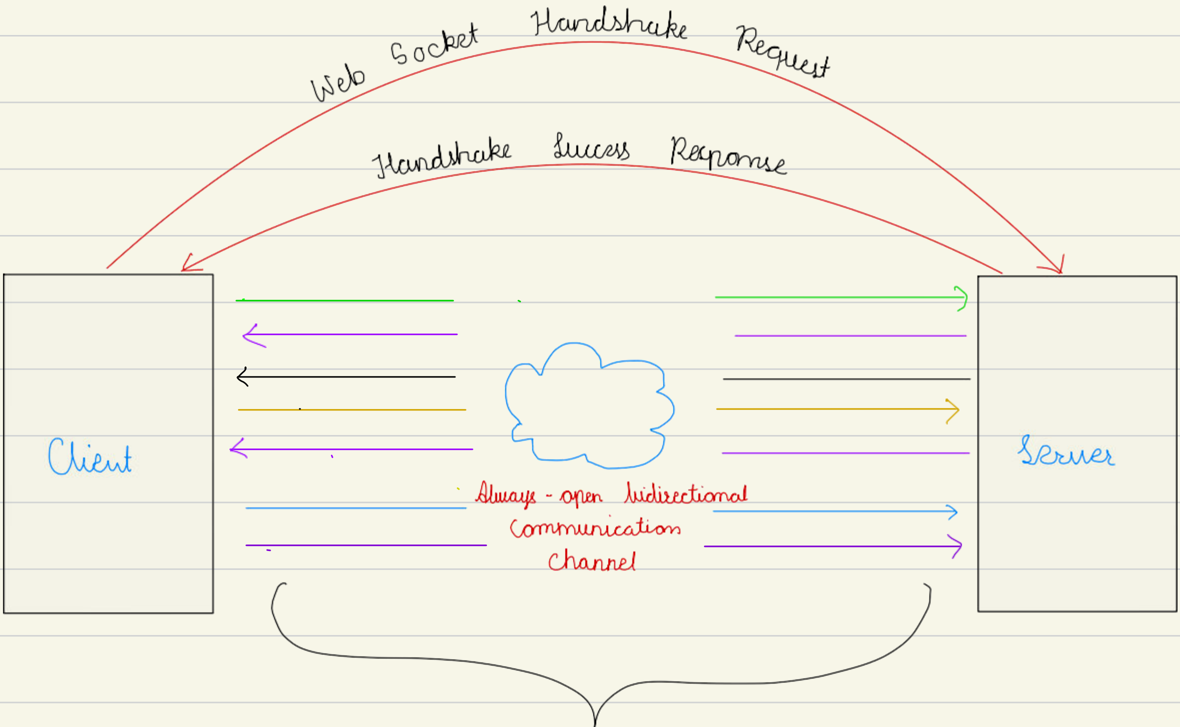


Web Sockets

→ Full duplex communication channel over single TCP connection.

→ Provides 'persistent communication'
(client & server can send data at anytime)

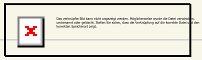
→ bidirectional communication in always open channel.



→ Lower Overheads

→ Real time data transfer

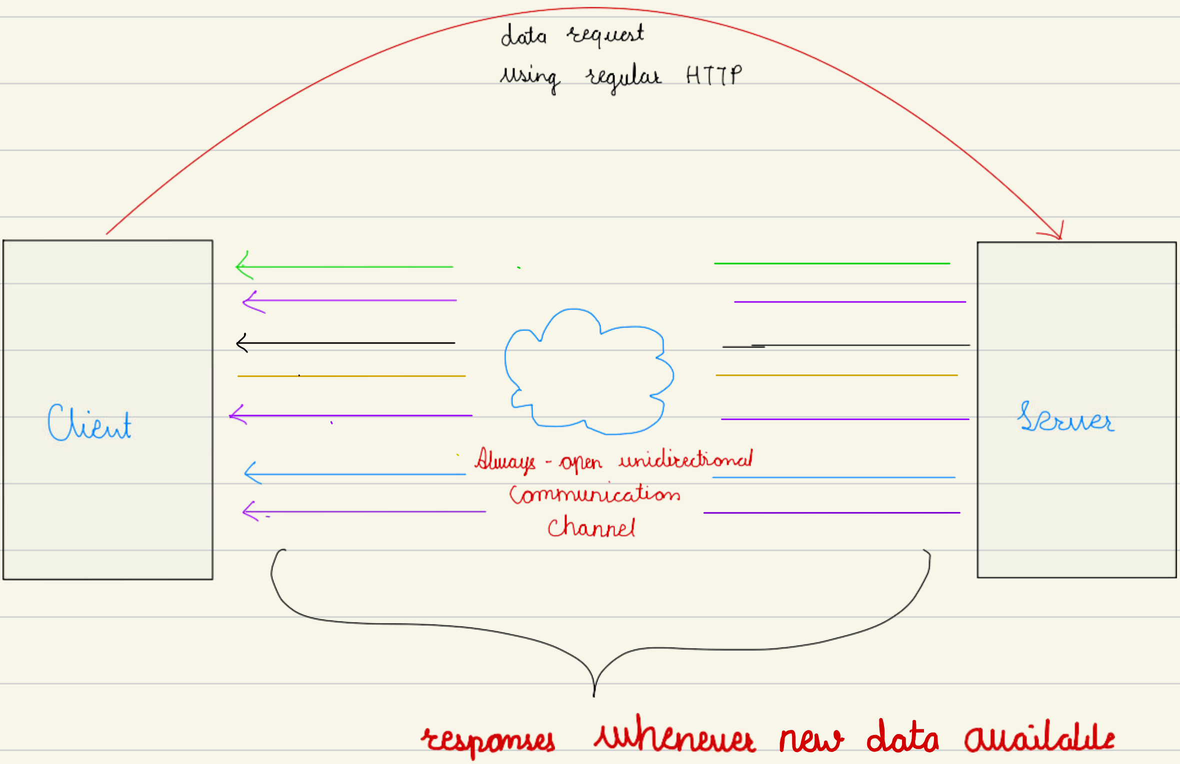
Server - Sent Events (SSE)

Client establishes persistent &  connection with server

Server uses this connection to send data to client

** If client wants to send data to server

↳ Requires another technology / protocol.



→ best when we need real-time data from server to client

OR server is generating data in a loop &

will be sending multiple events to the client.