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(1)

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O O

Language fundamentals

O Adentificans =

(3) Reserved Woods 3

) 3 Data types s

() ① Liteorals

() By Asserte (B) (1)

() (6) Types of Voorsables 22

() To version)

) (Main() method 30

(a) Command-line assignements 33

) Tava Coding Standards 34

) Identifier :-

- A name in Java program is Carled identifier, it can be class name on variable name on method name or label name.

Class Test Class name

p. 8. V. main (Storing [] angs)

int x=10;

yarriable name

yarriable name

* Rules to define identificous: -

1) The only allowed Characters in Java identifier are:

-> If we are using any other character we will get Compiletine Enour

2) identifier Can't Stattes with digit. Ept. X 128total

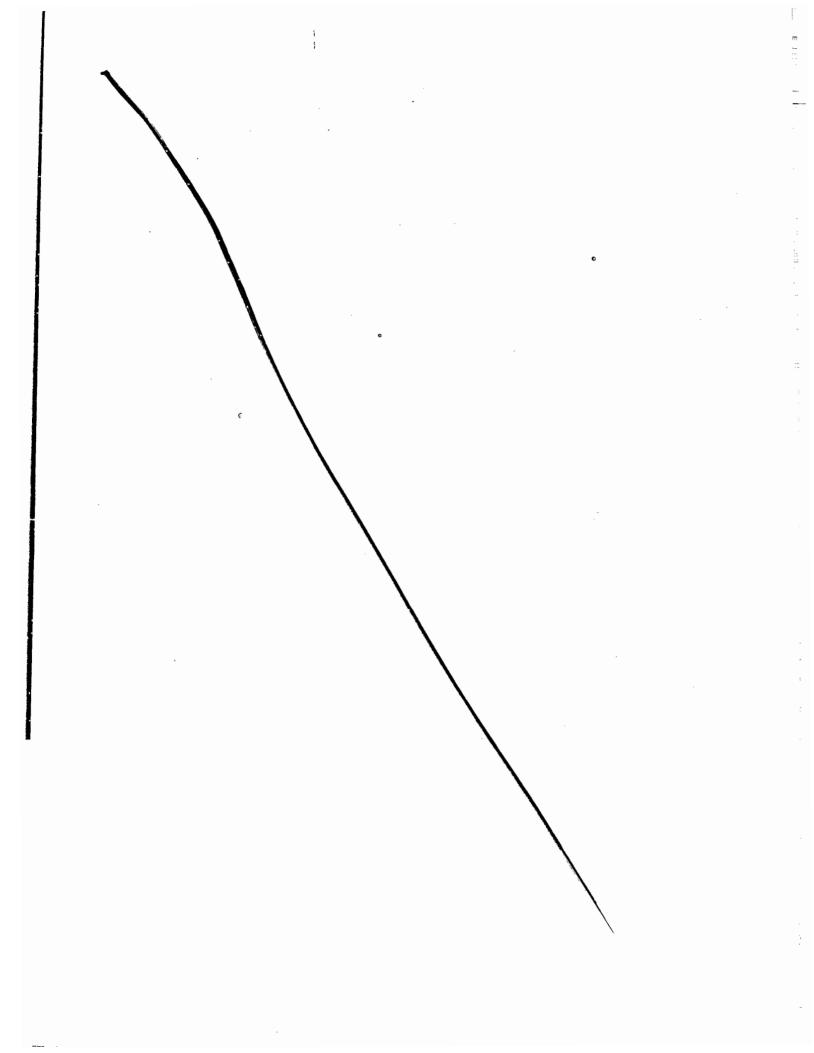
1 total 123.

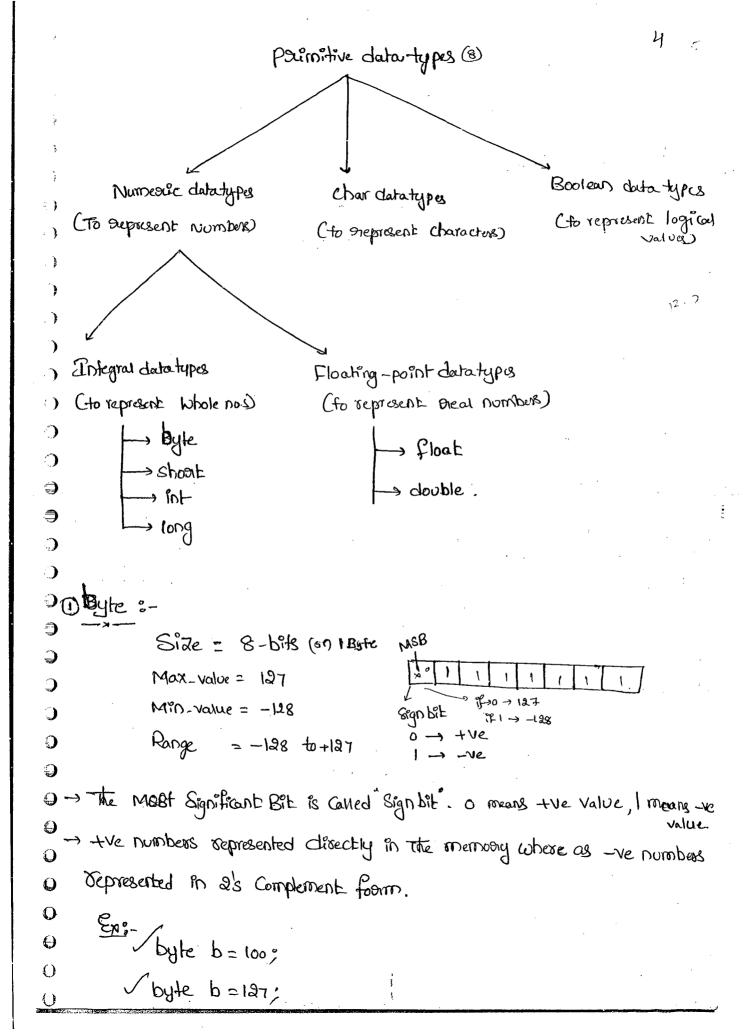
XEnt

V® Integeon

0

O





```
A byte b = 130; C.E.I. possible loss of precision
                               -found: int
                                 Required: byte
 xx byte b = 123.456, C.E. PLP
                                found : double
                                  Required = byte
 X byte b = toue ; CE: Ptp inCompatable types
                              -found: boolean
                                Required: byte
  X byte b="duaga";
                        C.E:- inCompatible types
                                                                       ௗ
                                -found: $ book . lang . Staving
                                                                       )
                                Dequired: byte.
                                                                       ാ
                                                                       )
-> byte datatype is best Suitable if we want to handle data in terms
                                                                       ()
  of Storeams either from the file or from the Network.
& Shoot :-
                                                                       0
                                                                       \odot
          Size: 2-bytes (16-bits)
          Range: -215 to 215-1
                                                                       9
                  -32768 to 32767
                                                                       0
   E.1- Shoak s = 32767
      Shoat 8 = -32768
                                                                       O
                                                                       0
                               c.E!-
     X Shoot S = 32768
                                                                       0
                                     -found : inte
                                      Dregulared I shoot
```

```
5
```

```
C.E. PLP
  X Shoot S = 123.456
                                 found: double
                                  Required : Shont
  X Shoot S = taue
                          C.E: 20000potable types
                                 -found: boolean
                                   Required: Shoot
  -> Most Gragnelly used datatype in Java is Shoot
  -> Shoot data-type is best suitable if we are Using
                                                       16-694 paccessors
   like 8086 but these perocessors are Completly outdated & hence
     Cossesponding Shoot data-type is also out dated.
)
) (3) int :-
   -> The most Commonly used datatype is Pril
9
)
            Size: 4-bytes
)
            Range: -2" to 2-1
•
9
                     -2147483648 to 2147483647)
Э
   Note o-
)
   -> En Clanguage The Side of int is varied from plateform to plateform
()
     -for 16-bit perocessons it is 2-bytes but for 32-bit perocessors is 4-bytes
0
4
    * The main advantage of This approach is aread & woute operation perform
0
     Very efficiently and performance will be improved. But The main
Q .
    * desadvantage of this approach is The chance of freeling a program
0
0
      is very very high if we are changing, plat-form, Flence c-language
0
```

is not Considered as Robust

-> But in Java The Size of int is always 4-bytes innerective of any plateform. * The main advantage of this approach is the chance of failling Java progoram is very very less, if we are changing anderlaying platform? Hence Java is Considered as Robust langue * But The main disadvantage in This approach is Gread & would Operations will become Costly & performance will be neduced. In c-language Java · balkework (BOX) ~ pexformance (JUF) will be neduced 39-bits • 16-bits 16-b9ts 32-bits 4-byles 2-byts * Not 4-bytes 4-bytes Roobust Robust C-proq C-prog Jana prog. P. ong Pes [3/08/11 4) long :--> When ever int is not enough to hora big value their we should To for long data type.) 0 Epas: To Diepoisent The amount of distance travelled by light in () 1000 days int is not enough Compulsary are should go for long type 0 501-1, long l= 1,23,000 x 60 x 60 x 24 x 1000 miles; 9 2 O

Size = 8 bytes

Range =
$$-2^{63}$$
 to 2^{63}

Nok:

) - All the above data-types (byte, shoot, int, long) ment food suppresenting.

Whole values.

) → 2f we want to suppresent steal numbers Compulsary we should go form

Thating point data types.

Joaling Point data types:

floating point data types

C

)

)

•

0

Size: 4-bytes

float

) a) Range: -3.4838 to 3.4838

2) 2f we want 5 to 6 decimal

o places of accusacy then we

O Should go for float

D float fallows Single poecesion

n Size: 8-bytes

25 Range: -1.7e308 to 1.7e308

3) If we want 14 to 15 decimal places of accusacy then we should go for double.

double

4) double follows double precision

```
Boolean data type: -
       Size: Not Applicable (Vistual machine dependent)
       Range: Not Applicable [BUE allowed values are true/false]
whic of the following boolean declarations core valid
     \times 1) boolean b = 0;
                                        in Compatable types
                                 C.Eo-
                                       -found & int
                                         Dieguired: boolean
       S) boolean b = toue;
       X 3) boolean b = True; c.E! Can't find Symbol
                                         Symbol: Variable True
                                          Location: Class Test
      X 4) boolean b= "faise" C-E: in Compatible types
                                         found : Java. lang-stoing
                                         sequired , boolean
        5) boolean Taue = tome
             boolean b = Taue
              P.o. pin (b); take
       int x=0;
                                                                           8.0. Pln("Hello");
                                              d 2.0 pln (" Hello");
                                                                          y
else
       8.0. pln (" Hi");
                         C.E. - in Compatible types
                                                                           0
                             -found : int
                                                                          0
                                                                          0
                              Dequired : boolean
```

-> The only allowed values for the boolean datatypes are "tome" on "false" where Gase is impositant.

Chan data-type :-

→ En could languages like c { C++ we an use only AscII characters and to shepswent all AscII chasacters 8-6945 are enough. here chasacters 8-6945 are enough. here chasacters 8-6945 are enough.

But in java we can use unicode Chasacters which Covers would wide all alphabetes sets. The no of unicode chasacters is >256 & hence 1-byte is not enough to represent all chasacters Compulsary we should go for 2-bytes.

⇒ Size: 2-bytes

⇒ Rapa: 0 to 6

)

)

Range: 0 to 65535

.	Summany of pourmitive data types:			Connesponding	
9	datatype	Size	Range	Worapper classes	default Value
9	byte	1-byte	-27 to 27-1 [-128 to 127]	Byte	٥
)	Shoat	a-byta	-2 to 2 -1 [-32768 to 32767]	Shont	0
9	int	4-bytes	-21 + 2-1 [-2147483648 to -] Integen	0
9	long	8-bytes	2147483647	Long	0
0	float	4-byles	-3.4638 to 3.4638	Float	0.0
U	double	8-bytes		Double	0.0
0		Ì			Company bloom
9 0	Chao _l	2-bytes	· .	Character	o [stepresents blank Space]
O	boolean	NA	NA (time false are allowed)	Boolean	I false (Theoret)

Literials 8-	
A Constant value which can be assign to the Vasciable is Call	1.4
Literal"	ieg
\mathbb{Z}_{0} : int $x = 10$;	3
	•
Constant value Literal.	' '
datatype/keyword	٠, ٠
1 Keyword	ر ، بر
Name of vascable identified	; ; }
Tal la	, ()
Integral Literals:	Ó
-> foor the Entegoral data types (byte, shoot, int, long) the following a one	0
has the following able	\odot
Various ways to specify Literal value	\bigcirc
1) decimal literals:	•
20	()
allowed digits agre 0 to 9)
En:- Pot x = 10;)
	0
8) Octal literals:	
-, allowed digits are 0 to 7	0
-> literial value should be perefixed with "o" [zero]	٩
	9
Epo. int x = 010;	0
3) Hexadecina literals:	9
→ allowed digits are o to 9, a tof 6, A to F	Ü
	O
→ for the Cxtona digits we can use both upper case & lower case.	•
This is one of very few places where Java is not case sensitive	Q O

The state of the s

```
8
```

```
-> Literia value should be priefixed with ox 600 0X
```

```
Ex= int x = 0 \times 10

int x = 0 \times 10
```

```
- these agre the only possible ways to specify integral literal.
```

```
Exe- class Test

\frac{1}{1}

\frac{1}{1}
```

) subset of the following declarations are Valid.

) int z=10;

) VE POT x = 066;

) X 3 Pot x = 0786; C.E: Poteger number too large

) int x : 0x FACE; 64206

) X (5) Pn + x = 0xBEER\$ C-E1- (after Bed : Excepted

0 80 x = 0xBea; 3050

```
-> Bythefault Every integral literal is of int type but we can
   Specify Explicitly as long-type by Suffixing with L or L.
    8x8-
0 int i = 10;
        X 2) int i = 101;
                               C.E ! PLP
                                     forma : long
        3) Long 1 = 101;
                                      Suguired : Int
        4) Long 1 = 10;
-> There is no way to Specify integral literal is to byte & short
   types Explecitly.
                                                                       )
-> 8f we are assigning integral literal to the byte variable & That
                                                                      •
   integral literial is with in the Drange of byte then it torects as
                                                                      -)
   byte literial automatically, similarly short literial also.
                                                                      En! -
             byte b=10 ,
             byte b = 130; X
                              C.E. brb
                               -found: int
                                Required : byte
 Floating point Literals:
                                                                      )
-> Every flooding point literal is bydefault double type & hence
                                                                      ()
   We Can't assign disrectly to float variable
                                                                      ()
- But we Can Specify Explicitly floating point literal is the float
   type by Suffixing with posi F.
     gel > float f = 123.456;
                                     found 1 double
          Ploat f= 123.456f;
                                     Tragined: Ploat
                                                                      0
             double 1 = 123-456;
```

```
-> We Can Specify floating point literal Explicitly as double-type of
      by Suffexing with doal.
          goi. double d = 123-45670;
               x float f= 123.4567d; C.E: PLP
                                              -bond : double
                                               Required: float
   → We Can Specify floating point literal only in decimal from &
     We Can't Specify in Octon & Hexa decimal form.
         D double d = 123.456;
         2) double d = 0123.456; 0/P!- 123.456
        X 3) double d = 0x123.456; c.E:- malformed floating point literal
9
) which of the following floating point declarations are valid?
     X) froat f=123.456;
     (2) double d= 0123.456;
     X 3) double d = 0x123.456;
)
     4) double d = 0xfae; /64206.0
                                    Because these 3 are not floating point
    5) froak f = 0xBea;
                                     So, that values age taking ent type.
    6 float f = 0642; /418.0
0
We Can assign integral literal directly to the Hoaling point datatypes
     4 That integral Literal Can be specified either in decimal-form or
O
    Octal form or hexa decimal from.
0
```

```
double
> But we can't assign floating point literals directly to the integral
                                                                      ( )
  types.
                                                                      ( )
                                                                      ()
        N int 1 = 183.456;
                                   PLP
                                 -found: double
                                                                      ()
                                  Icaquised: int
                                                                      ()
        V dauble d=1.203;
                                                                      ()
              0.00El; CD019.0.2
                                                                      ()
                                                                      - we Can Specify floating point Literal even in Scientific form
                                                                      also [exponetial form]
                                                                      ()
   en! 1) double d = 1.2e3;
               S.o.pln(d); 1200.0
     X 2) float f = 1.2e3; c.e. plp
                                     found = double
    3) float f= 1.2e3f;
                                       Dequired: floats
                              0/p!- 1200-0
Boolean Literals!
                                                                      0
-> The only possible values for the Boolean data types are true/falls
a) which to the following Boolean declarations are valid.
X O boolean b = 0; C.E. Incompatible types
                               -- found : int
X @ boolean b= Tome;
                                  Dreguired: boolean
                         C.E! Can't find Symbol
                                                                      0
 13 boolen b = tome;
                                Symbol : vaguable True
                                                                      O
A @ boolen b="true"; CEL- Encompatible tapes
                                                                      0
                                Looks: Java-lang-Storing Drequired: boolean,
```

```
J & int x=0;
           - TF LXS
                                             while (1)
             2-0. PIP ("Heno");
                                             S.o.pin("Heno");
             eise
            2.0.pln (" ++;"),
                          incompatible types
                  C.E :
                          found : int
                           Drequired: boolean
    Er0:-
                  \times
     int x =10;
                            10t x=10;
_
                                                     boolean b = toue,
                                                                     boolean b = toue;
     if(x = 20)
                            If (x == 20)
9
                                                   If (b = False)
                                                                     if (b== tone)
     8.0. pin (Heuo);
                            S. o. PID ("Hello");
                                                   8-0-p10("Heho");
                                                                      S.o.pln("Heno");
)
     else
                                                                     f
Э
                            else
     ¿ .o.pln("44");
)
                                                   ese
                                                                     else
Э
                            8.0.pln ("Hi").
                                                   20.b/U (4 th,)
Э
                                                                     S.o.blu("Hi"):
      C.E. IT
(
                            Olb. Hi
          f: int
9
                                                  0/P1- ++;
          R: boolean
                                                                    0/P!- +k110.
()
(
0
O
Ô
0
0
```

```
Chan Literals :-
 A Chan literal Can be suppresented as Single chanacter with in
    Single codes
       Col- chan ch = 'a';
          X Chan ch = a; c.E: Can't find Symbol
                                 Symbol: variable a
         X Chan Chz'ab',
                                location: Class XXXX
                            > C.E: Unclosed Character literal
                              C.E: Unclosed
                              C.E: not a Statement
32/08/11.
A chasi Literial can be deposesented as integral Literial which
                                                                     9
                                                                     €
    Suppresents conscode of that characters.
                                                                     0
- we can specify antegral literal either in decimal form or Octal form
                                                                     1
                                                                     ()
   001 Hexa decimal from. But allowed Drange 0 to 65535.
                                                                     •
                                                                     ാ
     &1-1) Chan Ch = 97;
                                                                     0
              8-0-P(ch); a
                                                                     0
      (1) chaon ch = 65585;
                                                                     0
               So.pln(ch);
                                                                     0
      X3) Chan ch = 65536; C.E. PLP
                                                                    0
                                 -found: int
                                 Deguired: Chair
     4) chan ch = OXFACE;
     (5) chan ch = 0640;
```

```
3, A chas literial Can be supresented in unicode supresentation
     Which is nothing but | Uxxxx 74-digit hera decimal no.
    En : ) chas ch = ' [ 10061';
                S-op (ch); a
     X 2) chasi ch = 'luabed; - semiscolon missing
     3) chan ch = 'luface';
     ×4) chan ch= 11 beat;
  4: Every escape character is a characterial
     En: D' chast ch = 'In';
       a) chan ch= '(t';
      X3) chan ch= 1/1/3.
.)
Э
    escape Character
                       Meaning
.)
9
       10
                    new line
       \t
                    hoocizental tab
       127
                    Cappiage Return
       16
                    Back Space
        16
                    form feed
                    Single acads
0
                   Double Quads
```

)

)

)

0 **9**

4

9

0

0

0

0

Back slash

```
of which of the following agre valid chan declasiations.
  1) chan
              ch = Oxbeaf;
 Ma) Chan ch = lubeaf; because
 (3) chan ch = -10;
× 4) chast ch = 1/x/;
       chaon ch = 'a';
 Staing Literals :-
 -> Any Sequence of characters with in "_" (double codes) is
   Called String Literial.
          String s= "fava";
-> The following powersoftions will be performed automatically by the Compiler.
                                                                 •
             -> Shoot
                                                                 9
                                 -long ----- float ----- double
```

- 1. Assay declaration
- 3 Array Creation.
-) 3. Assay Institutization.
 - 4. Declaration, Concation, Enitialization in a Single Line.
 - 5. length vs length ()
- G. Annonymous Assay
- 7. Assay element assaynments
-) 8. Assay Vascable assignments.
- Assay-

9

- An Assay is an Endexed Collection of fixed no of homogeneous
- I data elements.
-) The main advantage of assoray is we Can supresent multiple values
- Onder the Same name. So, that shadability of Code Proposoved.
- 3 -> But the main limitation of array is one we concated an array
- There is no chance of incoreasing / decaessing size based on own
- 3 Siequisiement. Hence memory point of view arrays Concept is not
- ne Commanded to use.
- 0 we can resolve this paroblem by using Collections.

```
1) Agray declarations;
 (a) Single d'imenshional Assay declasiation:
             intes a;
     √ a) int a[];
      √3) int [Ja;
                                                               . )
 - 1st one Succommanded because Type is adeasily seperated from
                                                               1 3
                                                               The Nome.
                                                               \bigcirc
-> At the time of declaration we Can't Specify the Size.
     ex:- XD int[6] a;
(b) 2D Assay declasion: =-
 / ) inf[][] a;
 va) not
                 [][]a;
     3) int
               atill;
     4) int[] a[];
 √ 5) int []
               []a;
  √ 6) int []a[];
                                                               0
                                                               ()
                                                               0
```

```
13
   c) 3D - Array declarations:
        interest a;
     Ø
         int
              ac1[1[1;
     9)
        int
              [][][]a;
        nt[]
                [][Ja;
         int[]
                acici:
        inf[]
                 []al];
         [][] interior
                 []0[]
         [][] 4ni
                 arj;
               [][]a[];
        int
         int
               [][[][]
9) Which of the following agre valid declarations.
     1) "MECI a,b; a -11
    (a) inter acrib; a \rightarrow 2
    3) PNETJ []a,b; a→2
b→2
    4) int[] []a,b[]; a-12
b-3
   X5) int[1 []a,[]b; a→20 (E:-
  If we want to Specify the demenshion before the variable
   It is possible only for the first variable.
 E01-
       inec] []a, []b,
                      not allowed;
```

}

.)

)

•

Э

0

3

3

)

)

 \mathbf{C}

1

)

)

Ó

9

9

0

0

0

0

0

0

Allowed

9) Assiry Construction;						
-> Every agray in Java is an object. hence we can Greate by						
Using new ancentral						
	otijest					
These classes none Dob analysis of these classes are available but						
These classes rare not applicable for programmer Level.						
	1					
-Abronay type	Coonesponding classname					
⊙ intc]	[]C					
(3) SUF [][]	[[10=					
3 double[]	0 @41					
))						
,						
-> At the time of Construction Compulsory use should specify the						
Size Otherwise we will get C.E.						
ENI ÎNTE] a = DEM	inf[]; X C.E!					
intil a = new						
-> Ph la loom -h hour in	Oppose with Size of the byla.					
·	O					
epol PotEI a = ne	ev int[o];					
- If we are Specifying assor	au Rizo as —ve int value, we will get					
The we come Specifying associal still as we not value, we will get of Stuntime Exception Souther to New York Association to						
Tieganverissay Sate Maphon.						
Sol Inter a = new inte-6]; R.E.I. Negative Assay & Reception						

=

Chasi. If we are using any other type we will get C.E. en of pote 1 a= new Pote [a]; A=65 byte-short @ byle b=10; **()** M INE[] Q= new int[b]. int) - long - front Shoot 5=20; (3) · Winter. a = new notes); X INE[] a= Dew int(101]; ()× int[] a= new int[10+5]; **()** ()Mofe: -> The max. allowed assaysize in java is 2147483647 (max. value of **_** int datatype) Coreation of 2D-Assays:) - In java morti dimenshional agrays agre not implemented in materix from they implemented by Using Assay of Assay Concept. **£**)) -> The main advantage of This appeach is memory officiation will be improved. ·) ex: int[][] a = new int[][]; **(** a [o] = new int[a]; \mathbf{Q} a[i] = New int[i]; a[2] = new int[3]; 0 O Note: _ 2n C++, Q91 0 Ð

0

-> To Specify, agray Size The allowed data-types agre byte, short int

<u>en</u> 91. intelled a = new intelled a [o] = new int [3][]; a [o][o] = new int[i]; a [o] [i] = New int [2]; outo][2] = new int[3]; a[1] = new int[2][2]; 9: which of the following Assay declarations are valid) X O gnt[] a = new int[]; √② PNE[][] a = New int[3][2]; √ (3) M=(1[] a = n, w 9nx[3][]; > nt[][] a = new int[][]; Con Con Con (5) int[][] a = new int[3][4][5]; (6) Pot [7[1] a = new sot [3][4][7; XA INFCICICI a = New INFCICION Assay Pritializations--> Whenever use one Coreating an abonay automatically every element is initialized with default values. O exco). int[] a= new int[3]; Sophola), [1@3=25a5 hashcode S.o.ph (acol), o 0 0 Note: Whenever we are trying to paint any object reforence internally to Storings, will be call which is implemented as failous.

```
Ex (2):-
                            WE BILDI
        interes a = new
1
          S.o.pln(a);
                         [[]a_____
20
          8.0.pln(aca); [IQ 4567
                                                                       0/23
                                                         a[0]
                                                                  Jall
          S.o.pin(acoicoi); a
                                            int (3)[2]
)
. }
   Ex(3):-
        intejej a = new integiej,
1
          S.o.plo (a); [[10----
)
つ
         Sopho(aw); null
9
         S-o-PIn (alostos); R.E. NPE
3
•
  - Once we careated an assnay Every element by default initialized
(
    with default values. If we agre not Salfsfy with those default values
(
    Then We Can override, Those with our Customized values.
)
)
      Ex %-
\bigcirc
             int[] a = new int[5];
()
             a[0] = 10 =
0
             ali] = 20%
                                                         o to
0
             (02) = (13)
0
             a[50] = 50; R.E. AIOBE
0
             a[-60] = 60; ________ R.E.: Alobe
0
0
                      -> C-E:- PCP, found = double, thequited = int.
0
   -> 8f we are trying to access an array with out of Trange index we will
     get RuntimeException Saying AIOBE.
```

```
Alononay declaration, Construction & Britialization in a Single Line:
 → We Can
              declasie, Construct & Britialize an asistay into a
   Singletine.
   Excu!
     int[] a;
                                 int[] a = \ 10, 20, 30, 40 };
       Q[0] =10;
       a [ i] = 20 ,
       a[2] = 30;
       0 [3] = 40;
  Ex(8)!
            Chan[] ch={'a', 'e', 'o', 'u'};
                                                                    Э
            Stocy[] S=/"Soon", 'Ravi', "Laxmi', "Lundan');
                                                                    )
→ we Can Extend This shookCut Even foor multidimenshional
  asistays also.
  Ex(B) :-
      int[][] a= { {30,40,50}, {60,70}}, a
                                                                    0
                                                                    0
- We can External This Shoot Cut Even for 30 agarage
                                                                    0
                                                                    0
 ۔ ام
   ENECJCICI a= // 210,20,30/, 240,50/, 260/, {20,80/, 190,00/, (110)}
                                                                    0
                                                                    ()
```

B: int[][][] a = // (10,20,30), (40,50), 760), , of 70,803, 140, 100), 21105) S.o.plo(aci)[2][3]), RE:- AIOBE S.o.pln(aco][i][o]); 40 S.o.plo(a[i][i][o]); 90 S.o.pln (a[3][1][2]), RE: AZOBE 2.0.pln (a[2][2](2]); REJ. ABOBE [100 20 30] 20 Pln (a(1)(1)(1); 100 06 ((1)(0)(0)/d 02 20 Pln (aci][0] (2]); A.E!- AZOBE -> If we want to use Shookcut Compulsory westrould perform declasiation, Construction & instialization is a Single Line. → If we are Using multiplelines we will get Compile time Esonor. Ex: intix = 10,20,30%:-POF X=10; 3-V PAECT X; int x; X x = {10,20,30}, C.E! - allegal Start of Expression.

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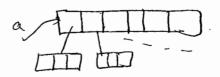
O

```
length vs length :-
  length:-
-> ZE is a final variable applicable only for assays.
-> 2L Stepaesents the State of assay
        int[] a = new int[10];
                                           Cannot find Symbol
                                           Symbol: method length
           S.o.pln (a.length); 10
S.o.pln (a. length()); C.E
                                                                     • )
                                           location: class int[]
                                                                     )
                                                                     Э
It is a final method applicable only for Strong Objects
                                                                     )
-> It Steppresents The no. of Characters procesent in String.
                                                                     •

    9:-
          Strang s = daga";
                                                                     ာ
             S.o.pho(S.length()); 5
            S.o.pln (s. length);
                    See: Cannot find Symbol
                               Symbol: Vascable length
                               location: java. lang. Storing.
                                                                     €
                                                                     0
    In multidimenshional assorages length variable stepresents only
                                                                     0
                                                                     0
    base Stae, but not total Stae.
```

=g!- int[][] a = new int[6][3];

S.o.pln (a.length), 6 S.o.pln (acoj.length), 3



Notes-

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Is applicable for String objects.

Annonymous Asissay:

> Sometimes we can create an array with out name also

Suchtype of nameless assays are Called "Annonymus assays".

→ The main objective of pannonymous Dooray is Just foor instant use.

(only ording)

) -> We can Create Annonymous Arroray as fallows.

New interdio, 20, 30, 40}

> At the time of Annonymous Assay Coneation we Can't Specify the Size, Otherwise we will get Compiletine Esister.

En! - x new int[u] { 10,20,30,40}

O Eg: Class Test

لا 10 د

P.S. v. main (Storing [] args)

```
Sum (new int[] 10,20,30,40}).
        Public Static Void 1800 Sum (int () x)
           int total = 0,
          for (int x1:x)
             total = total +x;
          S.o.pin (" the Sum: "+ total);
                                                                     )
Based on over siequisiement we can give the name for Annaymous
   assiay, then it is no longer Annonymous,
                                                                     a
                                                                     )
  £9!-
           Stainger & = new Stainger / A", "B");
            - S.o.p/n(s[o]); A
            So.pln(S[i]); B
            Sopho (Slength); 2.
                                                                     0
                                                                    0
```

```
Agrony element assignments:
Case(1) ;
-> for the posimitive type associates as Aspray elements we can provide
  any type which Can be paromoted to declare type.
O Eg: - 1, for the list type assays, The allowed Element types are
       byte, shoot, chase, int. if we are providing any other type
     , we will get Compiletime Esonosi.
     (1) :-
               int[] a = new int[10];
                √a[0] =10;
                ~acij='a';
                    byte b=10;
                 ~a[2] = b;
                   8hoot 5=20;
                   (8 = 1830
                   a[4] = (02) C.E! - PLP
                                       found! horny
              X a[5] = 10.5; C.F! - PLP, -Bund idouble
                                           Dequired; PDL
Egra): for the float type array, The arrowed Element types
     agre byte, short, chag, int long, float.
       byte -> short
                             int -- long -- floating double
              Chan
```

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```
-> En the Case of Object type assorage as assorage doments we can
    posovide either declared type on its child class Objects.
   eg1:.
          Number [] n = new Number [10];
         /DEO] = New Integera (10);
         VD[i] = Dew
                         Doubk (10.5);
       >> D[2] = Dero Stacing ("dworg"); → C.E: - ZnCompatiable types
                                             -found: String
                                              Stequired: Number.
         Object []
                  a = new Object[10];
                  = new Object();
          /a[6]
                                                         Object.
          ✓ a[1] = Dew Enteger (10);
                                             Priveto
                                                        Number (abstract)
          √a[2] = new Double (105);
          Va[3] = new Storing ("dogge");
                                                            100F
Case(3):-
                                                                     )
                                                                     ()
  -> En the Case of abstract type arrays as array elements we
                                                                    • )
    Can perovide its child Go Class Objects.
                                                                     ()
                                                                     0
            Number [10];
                                                                     0
                                                                     €
            ~ D [0] = Dew Enseger (10);
             > D[1] = New Number(),
                                                                     0
                                                                     0
                                                                     •
                                                                     0
```

Case(8):-

```
46)
    case 4!
     -> In the Case of Interface type assay, as assay element we
       Can perovide it's implementation class Objects
        옑.
              Runnable[] 91= new Romable[10];
                .9[0] = new Thread(); L
                                                                 Ronnable (I)
            × 9, [1] = new Strang ("dorga"); C.E! - Encomptatabletype)
                                                                 Thread (c)
                                              -food; Stoing
7
                                               Required ! Runnable
    Note:-
7
Assay type
)
                                  -Ollowed element-type
)
   1. Poumstive type arrays
                            -Any type which an be implicitly promoted
€
                           to declared type.
)
   2. Object type assays
                            Either declared type Objects on it's child
)
9
                                  Objects
                            class
       abstract class type
Э
                            86's child class objects are allowed.
         Shores
)
()
       anterface thipe
                                                 class Objects age allowed
                            945
                                 implementation
Q
       0,3191a ys
```

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()

```
-Poronay vasiable Assignment:
```

Case(1): -> Element lever peromotions agre not applicable at agrangement A chas value Can be paromoted to into type. Chasianay (chanis) Can't be Promoted to intil type. 10 POFET a= 110,20,30,40); chants ch= la', b', c'b Vintea b=a; In Compatiable type pintel c=ch, CEL -found : charc] Organized: int[] a) which of the following peromotions agre valid. O chan - int yo @ chan [] _ nt[] √3 int ____ slong × 9 int[] - long[] X 1 long - int X @ long[] - double[]

ì

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Staing -> Object (parent)

(chilly)

(Staing[] -> Object()

eg: Child type assay we an assign to the passent type vastable o

```
ar a
```

-> Child-type assay we an assign to the passent type vasciable.

Eges String [] S= /4 A7, "B", "B", "",";

√ Object() a = 8;

(ase(3):-

)

9

)

)

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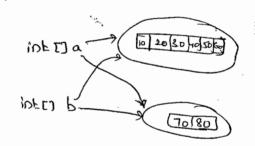
O

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→ When even we asse assigning one annay to another annay only sufference variables will be reassign but not underlying elements. Hence types must be matched but not Sezed.

eg: - (1) POEE] a = {10, 20, 30, 40, 50, 60};



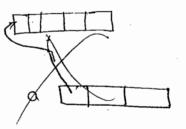
Eg(3). Int[][] a= new int[3](2],

a [o] = new int[s];

aci] = new int(y];

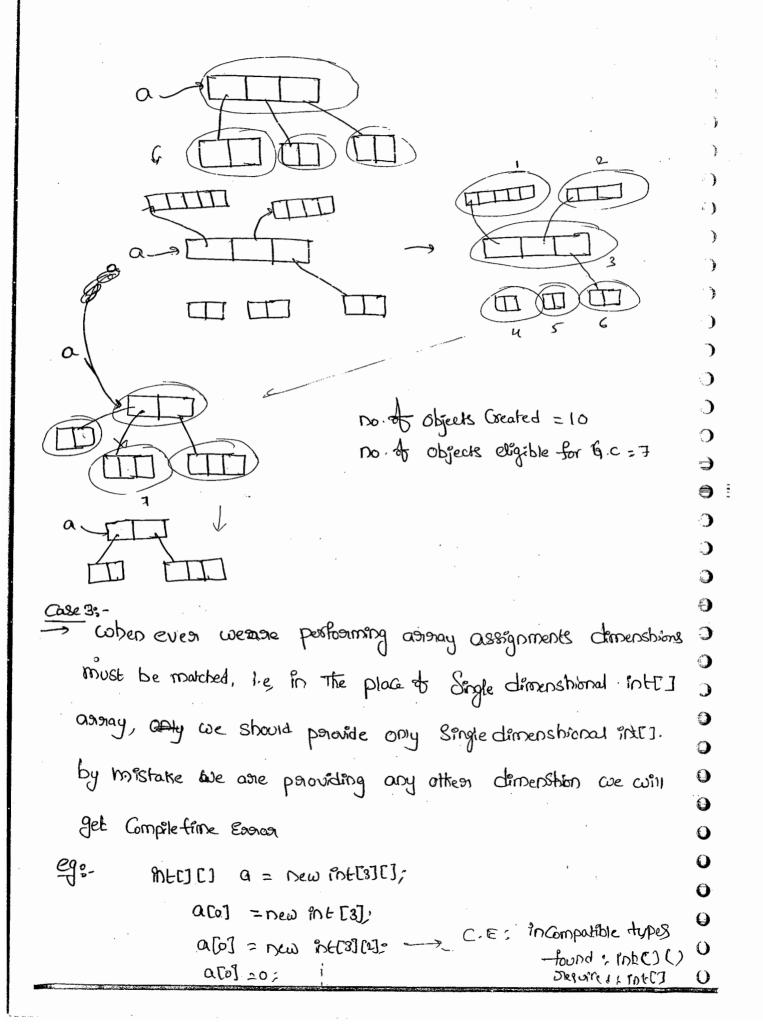
a = new int[e][3].

a[o] = new int[2];



no of objects Consisted = 10

no of objects eligible for G.c=7.



```
a[a] =10; C.E!. in Compatible types found: int

9 required: int[]
```

Types of Vasiables

- Based on the type of value neponesented by a variable, an variable, an variable, and v
 - (i) painitue Vasiables
 - (i) The fearence vasciables
- (1) Pormitive Vasciables 1

.)

_)

9

 \odot

)

0

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-) → Can be used to suppresent positive values
 - exi: int x = 10;
-) (ii) greference variables!
-) → Can be used to nefeon Objects
 - Cal- Student & = new Student();

g es Defexence vasiable

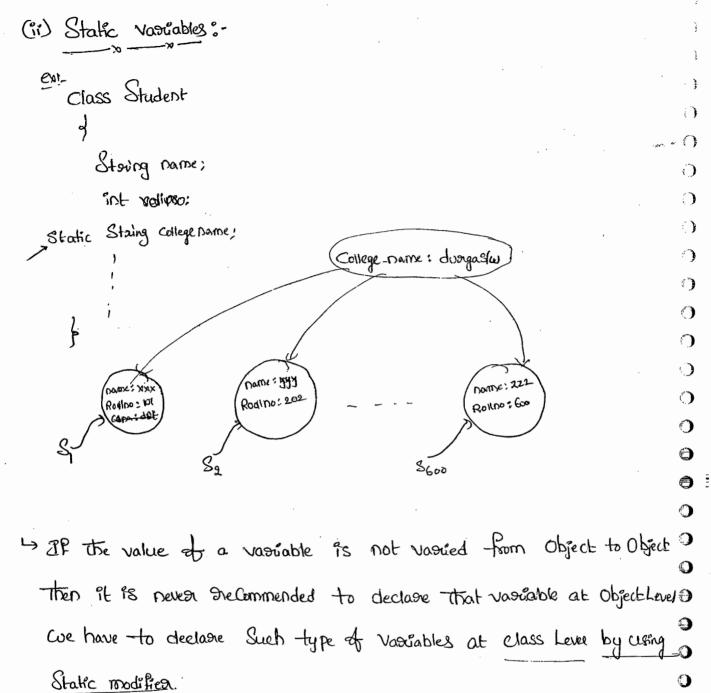
Desired on the purpose & position of declaration all variables of declaration all variables of declaration all variables

- (i) instance vaxiables
- O (") Static variables
- O (ii) local variables.

```
(i) instance vasciable:-
-> Af The value of a variable is varied from Object to Object
   Suchtype of variables are Cared instance variable.
-> from every Object a Seperate Copy of instance variable will be
   Coeated.
 -> the Scope of instance variables is exactly same as the Scope
   of the Objects. because Instance vasciables coil be Coneated at the
                                                                       )
   time of Objects Coneation & clastoney at the time of Objects destruction?
-> Instance vasiables will be stoored as the past of Objects
                                                                      3
- Instance variables should be declare with in the class directly,
                                                                      9
                                                                       Э
   But outside of any method on Block on Construction,
                                                                      )
instance vasiables Cannot be accessed from insta Static asrea
                                                                      ₹)
  districtly we an access by using Object diefeorence.
                                                                       Э
-> Bost from instance asear we can access instance members directly
                                                                       )
                                                                      0
 Ep1.
           Class Test
                                                                      ()
                                                                      0
              int &= 10;
             (2012 1) privet 2) m.v.2.9
                                                                      0
                                        non-Static variable & Cannot 0
               S.o.pln(x); -> C.E:
                                        be referenced from Static Contact
```

```
B
```

```
Test t = new Test();
              Sophotas; 10~
              Public usid mich
• }
. )
S.o.pln(x); Vio
( )
( )
   -> foor the Phalance variables it is not stequioned to perform
     initialization Explicitly, Jum will provide default values.
    £g:-
               class Test.
)
                                                         "restan Graziches
                                                 Ep!.
)
                 Starny S;
$
.)
                  intx)
Э
                  boolean b;
)
                                                       Students objects, In that
               P.S. v.m (Storings) args)
)
                                                          name, Rollinos are instance
                                                           variables, BCZ, These values
•
                                                          are varied from Object to
)
                  Test t = new Test(),
                                                           Object
C
                   Sorph (t.s); null
0
                   S.o.pln(E.x); 0
9
0
                   S.o.pln(t.b); false
0
0
0
       Instance Vastiables also known as "Object level vastiables" of
O
0
       attributes.
0
```



In the Case of instance variables for every Object a Seperate Copy will be Created, But In the Case of Static Variable Single Copy will be Created at class Level & the Copy will be Shorted by all Objects of that class.

Static variables will be created at the time of class Loading & destory of at the time of class unloading. Hence the Scope of the Static variable is o

Exactly Same as the Scope of the class.

ì	Note- Java Test I execution proass is
: }	1) Start Jum
1)	2 Coreate main thread
)	
.)	(3) Locate Test-class (4) Locate Test-class (5) Static Vasiables Coneation
1)	(4) Load Test-class Static variating Cheanon
)	6 Execute main() meltod of Test-class
)	
•	© Unload Test-class — Static variables destruction
)	1 Destanoy main Thread
) 2	8 ShutDown Jum
))	
3	-> Static variables should be dealers will sent a standing
9	Static variables should be declare with in the class directly
Э	(but outside of any method on Block on Construction), with Static-
ϵ	modifica.
Ĵ	
•	-> Static variables can be accessed either by using class name or by
•	Using Object and a second of the control of the con
()	Using Object Deference, but DeCommended to Use Class name.
0	- With in the Same class event it's not required to use class name.
9	also
9	also we can access disrectly.
Ð	En!- class Test
O)
Q	Static int $x=10$;
Q	D. S. 24 mg 20 C 61 4
€	= $=$ $=$ $=$ $=$ $=$ $=$ $=$ $=$ $=$
0	0-0-pin(Test.x); /10 &o.pin(t.x); /10
O	8-0.pln (x); ~ 6

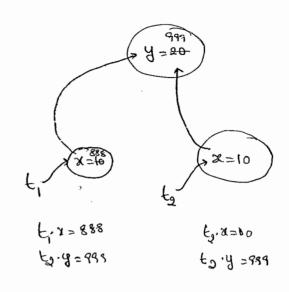
```
Static variables are Coreated at the time of class loading ie,
     (at the begining of the powgram). Hence, we can access from both
    instance & Static asseas dissectly.
             Class
                    Test
                 Static int x=10;
                 P. S. V. M (Stating[] ange)
                    S.o.pln(x);
                  Public void mill)
                    8.0. Pln (x);
                                                                             .)
-> - Food the Static vascables it is not sequioned to perform initialization
                                                                             )
  Explicitly, Compulsary Jum will provide default values.
               Class Test
                 Static int x;
                 P.S. V. m (Staing[] args)
                   S.o. pln(x); o
```

→ Static Vascables will be Stored in method-asea. Static vascables also known as class-level vascables or fields

; } Class Test () 1) int &=10; Static int y=20; P.S. V.M (Stockall args) \odot Test ti= new Test();) t.x =888; Э 0 t1.4 =999; **3** Test to = new Testo; **9** S.o.pln(to.x+"---"+b.y) Э) 0

)

0



The coe Performing any change for Postane variables these Changes wont be reflected for the remaining Objects because, from every Object a Seperate Copy of instance variables will be their.

But, if we are performing any change to the Static variable, these Changes will be reflected for all objects because to are maintaining a Single Copy.

```
Local Variables:-
-> To meet templasing sieguisnements of the paragramer sometimes
                                                                        .
  We have to Coreate Vasciables inside method on Block on Construction.
                                                                        , }
  Suchtype of variables are Called Local variables.
                                                                         ()
-> Local variables also known as Stack variables or Automatic variables
                                                                        )′
    or temporary variables.
                                                                        .
                                                                        -> Local vasciables Will be Storied inside a Stack.
                                                                        -> The Local variables will be created while Executing the block
                                                                        i 🕽
   in which we declared it & destoryed once the Block Completed. Herre,
                                                                        \odot
  The Scope of Vasiable is Executly Same as the Block in which we
                                                                        \odot
                                                                        declared it.
 Cx;-
            Class Test
              p. s. v.m (Storings)
                int 1=0;
                for ( int j=0; j=3; j++)
                   ? = ĭ+j;
                                                   Can't find Symbol
                                                                        €
```

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Symbol: vascable J

Location! Class Test

```
26
```

```
- For the Local vascables Jum won't provide any default values,
     Compulsably we should perfoom initialization Explicitly, before using
     That Vascable.
1
()
    [일:- 0
           Class Test
                                             class Test
.)
             P·S·V·M(Storing[] args)
                                              P·S·v·m (Storing[] args)
.
               Pot x;
                                                  int x;
           S.o.pln (Hello);
                                                  8.0.pln(x);
Э
)
           %P!- Heno
3
                                               C.E.
                                            Vasiable & might not have
•)
                                             been initialized.
Ç
    eg(a) ;-
9
               Class Test
.
)
                  P. S. v.m (Storing [] angk)
0
                  int a;
9
                  if (asys. length >0)
()
2c = 10;
0
O
                  S.o.pin(x);
\Theta
                                    variable & might not have been initialized
0
```

```
Class Test
    G3!
                 f
                   P.S. V. m (Storing[] args)
                                                                               ()
                                                                               į j
                                                                               0
                      int x;
                                                                               ()
                     TF (asigs.length >0)
                                                                               ()
                                                                               x=10;
                                                                               0
                       else
                                                                               \bigcirc
                                                                              \mathbf{O}
                         x=20;
                                                                              \bigcirc
                    Soph(x);
                                                                              ()
                                                                              3
                                                                              a
                                                                              ()
            P/9 !
                    Java Test +
                                                                              )
                      20
                                                                              0
                                                                              \Theta
                    Java Test X y
                      10
                                                                              0
                                                                              )
- Nole !-
                                                                              0
 -> BE is not succommended to perform initialization of Local variables
                                                                              9
                                                                              0
    inside logical blocks because there is no garantee exembion of these
                                                                              •
                                                         blocks at Juntine.
-> BE is highly enecommended to perstoom initialization for the local variables
   at the time of declaration, at least with default values
                                                                              0
                                                                              Û
                                                                              O
                                                                              \mathbf{O}
```

```
-> The only applicable modifies for the local vasiables is final.
     If we are using any other modifier we will get Compile-time Error.
;
   Eg:
0
              Class Test
• )
               P. S. V. m (Storing C) args)
\cdot
                   pouvate int x=10;
1
                                           C.E! - Illegal Stabit of Expression.
                   Public int x=10;
)
7
                    Protected int x=10;
Э
               × Static int x=10;
0
                  -final int x=10,
3
3
1
)
0
0
    Un Britialized - Assays:

)
       Class Test
)
        POFES as
        P.S. V.M (Stocking F) argo
()
0
         TEST t, = new TEST().
          S.o. pln(t, a);
                            neel)
0
          S-o-pin (t, -acos); Nunpointer Exception
O
O
```

```
instance level:
                        S. 0. P(Obj. a)
                                       Duli
    Por [0] a ;
                                                                           ( )
                        S-o-p(obj a bal) Null pointer Exception
     i'e a=nun
                                                                           ( )
                                                                           ( )
                           S.o.p(obj.a) [1@1a2b3
   Mr[] a = new int[3];
                                                                           ()
                           o ([0] a (do) 9.0.2
                                                                           \odot
                                                                           ()
Staticlever !
                                                                           )
                                                                           Static intil as
                       S.o.p(a);
                                   nuy
                                                                           \bigcirc
                       Soplawi); NDE
                                                                           )
                                                                           \odot
 Static int[] a = new int[3];
                                  S.o.p(a); [[01234
                                                                           \bigcirc
                                   S. o. pcasos); o
                                                                           ə :
Explanation:
  int[]a; - here The assay (i.e object) Sie feven a is Coeate but its not
  initialized (i.e object is not) Caeated. So JVM parvides nun value to
                                                                           ()
                                                                           4
 The Variable a.
 Potes a = new int[3]; -> here becor of new operator we agre Greating
                                                                           )
 an Object and jum by default provides o value in assnay
                                                                           0
Local Level:
                                       C.E.: Nascable a might not have
                         S.O.P (a)
                                                                           9
                                                been initialized
  int []a ;
                         Co.p(a(d)
                                                                           0
  intel a = new int(3); So.p(a)
                                                                           0
                                    EI@1234
                        (600) q. o.2
                                                                           0
                                                                           0
     Once an Assay is Careated all its elements asse always
                                                                           0
                                                                           0
initialized with default values is spective weather it is Static or
                                                                           0
 instance og Local assay,
                                                                           0
```

```
Van-ang methods (1.5 version)
```

```
-> Until 1.4 version we can't declare a method with variable no. It advanguements, if there is any change in no. It assignements Compulsory we should declare a new method. This approach in weases length the Code & Deduces Dreadability.
```

FO Diesolve These peroblem Sun people interodocad vasi-arg method in 1.5 version. Hence from 1.5 version onwards we can declare a method with variable no of aprogrements Such type of methods one called vari-arg methods.

> We can declarae vari-arg method as fallows.

```
m (int... x)
```

→ We can invoke this method by possing any no. of int values
 including zero no. also.

```
m, (10); /
m, (10); /
m, (10); /
```

) <u>Call</u>:-

7

 \odot

)

Э

9

)

)

9

0

0

0

0

0

0

0

0

mg (10, 90);

My (60, 20, 30, 40)!

```
Class Test

2

P.S. void mi ("nt... i)

P.S. void mi ("nt... i)

P.S. v. m (String[] args)

Inscript
```

```
-> Enternally van-any method is implemented by using single dimenshime
   advays Concept. Hence with in the Vari-any method we can differencial
  agranguements by using index.
  Coj.
            Class Test
                                                                     )
             Public Static void Sum (int...x)
               Pot total = 0;
              for (Post y: x)
                 total = total +y;
                S-o-pin (" The Sum: "+ total);
                P.S.V.M (Storng [] avgs)
                  Sum ();
                  (10, 20);
                  Sum (10, 20, 30) 60
                  Sum (10, 20, 30, 40); 100
            The Sum: 0
             The SUM: 30
              The Sum: 60
             The Sum: 000
```

```
Z
```

```
Case(1):-
   (a) which the following was any method declarations are Valid.
             mi (int ... x) ~
\langle \cdot \rangle
             mi (int
                     x...) X
             m) (int
                       ...x) L
              micent.
                      ..x) X
              m, (int .x..) x
7 Case 21.
) - we can mix vasi-any parameter with normal parameter also.
            mi (int x, Storing...y)
→ Case 89.
  -> Ref core one miving vari-arig parameter with general parameter
     Then Van-any parameter should be last parameter.
      Sol.
              mi($ht... x, Storing y) X
  Case 4:
  -> En any Vasi-asig method we can take only one Vour-arg
      Pagameter.
0
              m1 (Pot... &, Storg... y) X
0
  Caset: -
            Class Test
                                             p. s.v. m (Storing [] angs)
              p.s.v.mi(Pot i)
0
                                                 mill; var-ang
               S-opin ("General method");
                                                mi (10); General (only)
0
             P-S.v.mI (Ast... i)
0
                                                 101(10/30); Non-and
```

(Pero-nov") 19.0.2

```
→ 2n General Vari-arg method will get Least priority 1.e
      if no other method matched then only van-any method
      Will get chance. This is Similar to default case inside Switch
  Case 6:
     ex:
            Class Test
            d
p-s-v.mi(int[] x)
                 8-0.pln(" "n+c7");
              P-S-v.mi (int ... x)
                S.o.pln( " int...");
         City- Cannit declare Both mi (Potci) and mi (Int ...) in Test
 Vasi-ang Vs Single dimenshional annays.
  Case(1):
-> Where ever Single démenshional assoray present we can replace
  with van-any panameter.
               m_1(\hat{n}_1 + in_2 + in_3) \Rightarrow m_1(\hat{n}_1 + in_4 + in_3)
              Main (Storing [] args) => main (Storing... x)
                                                                          0
 Casea
-> where ever vari-and parameter present we Con't replace with Bingle
                                                                         0
  dimenshional agoray.
                                                                         0
                   m \cdot (up \cdot m \cdot m) \Longrightarrow m \cdot (up \cdot m \cdot m)
                                                                         0
                                                                         O
```

manl)!

→ Wheather the Class Contains Main() on not & wheather the main() is peroperly declared on not, these checkings are not responsibilities of Compiler. At nurtime, TVM is nesponsible for these checking.

→ 28 the TVM unable to find nequired main() then we will get nurtime Exception Saying NoSuch Method Englos: main.

Class Test

compile Javac Test. Java ~

But X Java Test ___ R.E? No Such Method Eggo : Main

) - Jun always Searches for the main() with the following Signature.

Public >

Static Void main (Staing 17 angs)

To Call by Jum

without existing

from any where

Object also Jum

has to can this method

Mame of method Which is Configured inside JVM

, Command-Line

assiguements.

main method Coon't Stetusin anything to Jum

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```
-> If we are performining any change to the above Signature
   We will get 9 vontime Exception Saying "No Such Method English: main".
- Any whope The following changes are acceptable.
   (1) we can change the Order of modifiers. i.e instead of
       public Static We Can take Static public
  (2) We Can declasse Storing 1 in any valid from
                  Sprea [] privets
                  Storing
                          E7 asigs
                  [Japen prijet]
  (3) Instead of aggs we Can take any valid Java identifier.
                                                                      Э
  (4) Enstead of Stornges we can take Vasi-ang Storing panameter.
        Staring...
                                                                      .)
         main (String[] angs) ==> main (String... angs)
                                                                      9
      main () Can be declassed with the following modifiers also
           (i) -final
           (ii) Synchronized
           (iii) Stockfp.
 ex1.
                 Test
          Class
                                                                      0
           Final Static Statictiff Synchronized public Void main (String... A)
            S. o. pln (" Hai dunga");
                                                                      0
                                                                      O
```

```
Which of the following main() declarations are valid
  (i) Public Static int main (String[] args) \chi
          Static public Void Main (Storing E1 args) X
     (iii) Public Synchronized Structfp final void main (Structg args) >
      (M) Public Final Static Void main (Strong ange) X
   (V) public storictip synchonorized Static void main (Storinger angs)
   a) In which of the above Cases we will get Compiletime Esmost.
            No where, All Cases will Compile.
 ) -> Enhercitance Concept is applicable for static methods including
      main() also. Hence if the child class doesn't contain main() then
      Passent Class main() will be executed while executing child class.
\Rightarrow
     Cx!-
            Class P
Ō
             Public Static void main (Storing C) args)
-_)
(.)
              S. opho (" ELU duonga slw");
`)
 .)
•
           class c extends P.
(پ
         Javac p. java
( )
          Java p
0
          of 1- PLU dvaga S/W
· {-)
         java C
()
                ZLU desiga Slw
```

```
Cx 21.
             Class P
              P. S. v. m (Strung [] args)
               2.0.pln (" I Love");
              class cextends p
               p.s.v.m (Storing[] args)
                S.o.pln(" dungas/w");
            JONAC P. java
            Java P
             OP! ILove
             olp! - duorga Slw.
-> 21 Seems to be oversiding Concept is applicable for Static
 methods, but it's not oversiding but it is methodhidding
-> Overloading concept is applicable for main() but JVM always Calls
   Storing Established any. The other method we have to
                                                                         )
  all Explicitly.
                                                                        Ĵ
        class Test
                                                                        9
 ep!-
          (typo 13 grivets) m.v. 2 . 9
                                         O/P:- dungasin.
           2.0.pln("duagas/w"),
          P.S.V.m (inter args)
           2.0. Pln (" & good");
```

```
a) Instead of main is it possible to Configure any other method
     as main method?
         Les, But inside Jum we have to configure Some changes
       Then "E is possible.
  a) Explain about So pln?
) <del>-}</del>)
                                                 Class System
      Class
              Test
                                                    Static PountStoream Out;
        Static Strung name = "dunga";
      Test. Dame. length U
                                                   System. out. pointln ()
                                                                       21 is a method
                              TE is a method
                                                                        present in
                              Possent in
                                                                        DrintStream
  Pt is a
            Static vásúable of
                               Storing class
                                              71 Psa
                                                              Static vasuable of
            type Stating present
  Class-
                                                Class Name
                                                              type pountstream
              Po Test class
                                               present in
.)
                                                              present in System
                                               Java. lang
)
                                                               Class
)
)
```

•

)

)

1

•

.)

()

()

1

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Commandine assignments:

- The assignments which asse passing from Commandporompt Gire Gulled CommandLine assignments.
- -> The main objective of Commandline assignements are we can Customize the behaviour of the main().

- asigs.length \Rightarrow 3
- exu !- class Test
 - P.S.V.m (Storing[] angs)

 for (PNE i=0; i=congs.length; i+t)

3

 \bigcirc

- S.o.pln (angscis),
- boof Java Test 1

 R.E.I. AIOBE
 - Java test x y L
 - X Y
 - R.E ! AIOBE

```
Cn(2):
```

```
→ with in the main(), Commandline assignments asse available
     in Storing from.
     CX:-
                Class Test
                 P. S. V. M (Storing [1 angs)

S. o. pln (angs [0] + angs [1]);
                 Java Test 10
                0/01- 1020
     Space is the Seperater Blow Command-
3
     Line assignments itself Contain Space Then we should enclose with
      in double codes (")
             Class Test
      ex!
-)
              P·S·V·m (Storing[7 angs)
                S.o.pln(angsto]); Note Book
)
                        Java Test " Note Book"
ر.
   exo!
          class Test
             P. S. V. m (Storing [] args)
                                                     3.0.pn(s):
               Staing[] argh= d"A", "B"?
               angs = angh;
               for (Storing
                            Si 🕆 avgs)
```

Java Test X Y H

OF A

B

Java Test X Y Z H

OF A

B

Java Test H

OF A

B

Java Test H

Note: The maximum allowed no. of Commandline assignments is 2147483647, min. is 'O'

Lava Coding Standards

whenever we are coeffing the Code It is highly recommended to favour Coding Convensions the name of the method are class should or reflect the Parpose of Functionality of that Component.

Class A

d

public int mi(ent x, enty)

d

seturn x+y;

Amoreospet Standard

package Com. dorgasoft.demo;

public class Caluctators

public Static int Sum (int number)

public Static int Sum (int number)

int numbers)

int numbers)

int numbers)

int numbers)

-Hotech-City

 \bigcirc

Coding Standards for Classes:

→ Usually Classnames one Nouns, should strouts with uppercase Letter & if it Contains multiple worlds every inner world should starts with uppercase Letter

```
Student
      Cpl-
             Customer
                              - NOUR
              Storing Buffer,
2) Coding Standards for Enterface: -
  - Usually interface names agre Adjectives should stark with UpperCase
    Letter & if it Contains multiple woords every inner woord Should Starts
    with UpperCase labler.
            Runnable, Sescializable, Cloneable, Movable, Adjactives
   Note:
     Therowable is a class but not intexface. It ack as a 9100t Class
    for all Java Exceptions & Essous.
€ 3) Coding Standands for Methods:
.)
   -> Usually method names asie Citton Verbs on Verb noon Combination
      Should Stark with Lower Case Letter & of It Contains multiple words
)
-)
     "Every innerwoods should starts with upper Case Letter. (camerlase).
<u>.</u>
)
   Col
         9101 ()
)
                               getName()
8 er. Salary() Veb + noon
         Sleep()
)
                   p- Vexbs
         eat ()
          init ()
)
         wait ()
Join ()
()
  (4) Coding Standards for Vasiables:
U
   -> Usually the vascable names asse nouns should starts with
```

Lower Case character & if it Confains multiple worlds, Every Prince world

Shoold Starts with upperase Character (Camel Case).

 Θ

Col Marine /	
91011 DO	
mobile Number - nouns	
Coding Standardsfor Constants;	
-> Usually The Constants are Noons, Should Contain Only Upper Case	
Characters, Zf 96 Contains multiple worlds, These woords are Seperated	
With "-" Symbol	
→ We Can declare Constants by using Static & final modifiers.	o '
epl_ MAX-VALUE	
MIN-VALUE	
MAX - PRIORETY	
MIN - PRIORITY	_
a James Com at 1 1	- · ·
Dava bean Coding Standards	
→ A Java blean is a Simple java class with poivate poroperties &	
Public getter & Setter methods.	
ew. Dublic class Student Bean	
POBLIC CLASS CIERTIFICATION	
Polivate Storing name;	٠.
DOL VILICUI COMPLICA	
Public void set Name (Staing Mame) - From SUN.	
the same name	
this - name;	
9	
Public String getNamel)	0
'd	4
neturn name;	€
1. 9	
i J	1

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```
35
   Syntax for setter method ...
 -> The method name should be prefixe with Set". Compulsary the
   Method Should take some assignment. Dietusis Eype Should be
    Void .
  Syntax for getter method:
 - the method name should be posefixed with get.
  → ZE should be no assignment method.
) -> 91etuantype should not be void.
) Note :.
 -> - For the boolean peroperty the gettern method can be perefixed with
  either get on is . secommended to use is
   exi- possuate boolean empty;
•
          public bookan getEmptyU
)
           return empty;
           public boolean is Empty U
           of Stetuon Empty;
4
O Coding Standards for Listoners:
    To Degister alistener:
  -> method name should be priefix with add,
```

-> after add what even we age taking The anonguement should be Same.

eg: 10 public void add My Action Listener (My Action Listener ()
X @ public void Sigister Myaction Listener (Myaction Listener L)
X 3 public void add myactionlistener (listener R)
To unsiegistes a Listener: The orde is Same as above, Except method name Should be
Perefex with Exemove.
=9!- O public void surmove My Action Listener (ray Action Listener 1)
Depublic Void unsnegisterMyactionListener (rayActionListener U)
X 3 public void delete MyAchion Listener (MyAchion Listener e)
X (2) public void 9temoveMyActionListener (ActionListener L)
Note: 2n Java bean Coding Standards ξ listency Conapt 1 compulsarily.

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book for SCJP.

Operators & Assignments

Incaement Decarment 2 Asithmetic openations 3 Concatarnation 5 Relational operations 5 Equality operations 6 Bitwise operators 7 Shoot-Ciorcuit 9 instance of 6 type Cast Openation 10 Assignment Operator 12 Conditional Operator 13 New operator 13 [] operator 13 Operator priecedence 14 Evaluation Oades of Java Operands. 14 Э

O

Encomment & Decarment Operators:

Pre- Postanchement anchement

Decorment

Post-decrement post-decrement

 \bigcirc

Expaession	Enitial value	-final value	final value	,
J = ++x)	H	6	5	
A = x++;	4	5	4	
A=x)	4	3	3	
A=x;	4	3	4	

i) we can apply incoment and decoment only for variables but not for Constant values.

ii) Nesting of incomment & decomment operations is not allowed otherwise we will get Compile time Edward.

```
Bt x= 4;
                                  C.E:
                                         Unexpected type.
        Pot y=++(++x);
                                      @ found : value
        S.0 pcy);
                                    1) Required: Variable
 iii). We Can't apply incoment & decoment operators for the
      final vasuables.
                                              -final int x= 4;
      Exu) - final int x = 4; X
                                    Ex(2)!-
                                                x = 5
              X++;
                C.E: Carle assign a value to final vasuable 2.
  is). We can apply incoment and Decoment operators for
-
     Every pormitive data type Except Boolean".
         double d=10.5;
             d++;
                                      Ob++;
                                    Sop(4); // b.
         8.0.pcd); 11.5
                  boolean b = taue;
                                        C-E:-
                      ++b;
             X
                                         operator ++ Can't applied to
                  S-0.Pln(b);
                                                               boolean.
      \sqrt{4} int x = 10;
           2++;
          8.0 pln(x); 11
```

)

)

)

-)

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```
Difference blw b++ & b=b+1 :-
                               byte b=10
   O byte b = 10;
                                  b=b+1;
         b++;
                                              C.E: Possible loss of precission
        Sorphich), 11
                                 8-0-P(P);
                                                   -found : int
                                                    Required : byte
      3 byte b =10
                                          Exp: max (Pot, type of a, type of b)
          b = (byk) (b+1)
                                                max (int, byte, int)
          80.blu (P):// 11
                                                 int
                                 Explanation!
    byte a =10;
    byte b = 20;
                                    Max (Pot, type of a, type of b)
                                                                        - )
    byte c = a+b;
                                     Max (int, byk, byk)
                    C.E: PLP
    Sorph (c);
                                       result is of type: int
                       f= nt
                        R = byte
                                               : found is int but
                                                    Dequired is byte
                                             (+,-,*, 1.,1)
                                                                        )
-> Whenever we are performing any arithmetic operation between
                                                                        )
   two variables a & b the result type is always,
                                                                        )
              Max (Int, type of a, type of b)
      byte b=10;
                                                                        ()
      b = (byte) (b+1);
                                                                        4
                                                                        0
       S-0-P(b); // 11
```

```
-> En the Case of Encomment & decomment operations the orequioned 39 type casting)

-type casting automatically performed by the Compiler.
```

byte b++;
$$\Rightarrow$$
 b = (byte)(b+1);
b++; \Rightarrow b = (typert b)(b+1);

Asithematic operators:

-> The Apriltmetic operations agre (+,-,*,1,%)

a and b the result type is always.

Max (int, type of a, type of b)

byte + byte = int

byte + shoot = int

Pot + long = long

long + float = float

double + chase = double

chasi + chasi = int

S.oph (10+0,0); // 10.0.

S. o.pln ('a' +b'); 195

8.0.pln (100+'a'); 197.

Infinity:

9

)

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-> En the Case of integoral a silternatic (int, shoort, long, byte), There

is no way to suppresent infinity. Hence, if the infinity is vesult

We will always get Asithernatic Exception. (AE = 1 by Zero)

Sio-pla (10/0); R.E. A.E. I by Zeono

-> Bot in Case of floating point assithematic, There is always a	
way to Shepshesent infinity. For this float & Double Classes	;
Contains the following two Constants.	
Positive - Enfinity = 2nfinity = -2nfinity = -2nfinity	
- Hence, in the Case of Pearl Floating point Assittematic we won't	() 2 (·
get any Arithmatic Exception.	()
Eg: 0. S.o.pln (10/0.0); Infinity	() ()
(a). S.o.pln (-10/0.0); - Infinity.)
Navn: - (Notan Number)	∂
-> 20 integral agrithematic. There is no way to suppresent undefine	4 0
Stesults. Hence, of the result is undefined we will get A.E	<u>ာ</u>
Po Case of Polegoial Avithematic.	Э
Eg:- S.o.p(0/0); RE: A.E: 1 by Zeono)
But in Case of floating point Assittematic, There is a way to sepreson	o to
undefined nesults for this Float & Bouble classes Contains)
Nan Constana	0
	C O
Hence, Eventhough the result is undefined we won't get any	O
Runtime Exception in flowling point Assittematic.	O

£g:-

S.o.ph(0/0.0); NaN.

```
10
```

```
* S.o.p(0.0/0); NaN
           * S.o.p (-0/0.0); NaN
    En public Static volouble Sgot (double d);
                 8.0. pln (math. Sqot (4)) ;/2.0
                 Sorph (math. Sport (-4)); NaN.
   - too any x value including Nan the below Exponessions always
      neturins false, Except the (!=) Expression returins true.
                    X != NOUN = Tonce
                                                          X> NaN
                                                          X>=NaN
      (10 > foot Nan);
                                          -false
                                                                 & False
 )
                                                          XX NaN
                                                          XX=NaN
_)
                Sop (10 < Float : Nan);
                                          false
                                                          X==NaN
9
              S.o.p (10 = = float. Nan);
                                          flose
.)
               S.o.P (10! = Float. Nan);
          S-o-P (float. Nan = = Float. Nan); false
-)
           S.O.P (Float Nan! = Float Nan); True.
-
   (enclusion!about A.E (Airthmetic Exception):
  -> It is Runtime Exception but not Compiletine Essosi.
    > Possible only in Integral Assithematic but not floating point Arithmet
                                                    ( float , double )
                     Cint, byte, shoot, chasi)
9
  The only openators which couse A-E are I and %.
0
```

3. Storing Concatanation operator (+)

```
-> the only over-loaded operator in Java is '+' operator.
-> Some times it acts as as: thematic addition openiation & Some time
   acts as Storing agrithematic Operators. (or) Storing Concatination Operators.
       Pot a =10, b=20, C=30;
        Storing d = " Shanth ";
                                Go shants
         8.0p (a+b+c+d);
         S.o.p(a+b+d+8); zoshantk30
                                                    d+a+b+C
                                                    Shanth 10+b+c
         S.o.p (d+a+b+c); Shanthio 2030
                                                      Shant51020 +C
                                                       Shanitio2030
         8.0.p (a+d+b+c); coshanth2030.
                                                                          )
                                                                          •
 -> 28 at least one operand is Storing type then '+' operator acts
                                                                          \Rightarrow
                                if both able numbertype)
   as Concatination, otherwise, + acts as asithematic operator.
                                                                          1
                                                                          )
   Here S.o.p() is avaluated from Left to Right.
   <u> 29</u>%-
                                                                          •
         ist a=10, b=20;
                                                                          9
          Storng c = "Shanth";
                                                                          )
       × a = (b+c) rotal Strains

c.E. En compatible type:
                                                                          )
                                                   found: stocking
                                                                          •
                                                    Required: int
       C = Q+C; Blooms
                                                                          )
       / b = a + b
          C = a+b; C.E: - En Compatible type:
                            found : int
                            Required: Stocing.
```

Relational Openations

```
These age 7, 4, >= , <=
we can apply Relational operators for Every pointive datatype
  Except boolean.
                                    5) tonue <= tonue
                    -Palse
                                    6) tave < false/
       3) 'a' < 'b'
                      Taue
                                               Operator <= Can't be
                                           CE:
                     Thue
       3) 10 >=10.0
                                            applied to boolean, botalean
       4) 'a' < 125
                   Torce
we Can't apply nelational openations for the object types.
 _Ego-)" Shanth " < "Shanth" &) "duonga" < "duonga123" X
       CE: openation < can't be applied to Storing, Storing.
Nesting of Relational Operators we are not allowed to apply.
  Egg. S.o.p (10 < 20);
     > 8.0. P (10 < 20 < 30)
                 boolean
    CE!- Operator < Carit be applied to boolean.
        Starg SI = new Starg ("duargo");
        Soing 82 = new Storing (^duonga'),
          Soph(S,=$S2); false (reference)
          S.o.pho(S, equals(Sg)); torue (content)
```

Equality Openations (==,!=)

-> These age == , !=

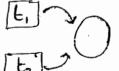
* we can apply Equality operations for Every primitive type including

OLP

→ We Can apply Equality operators over for object oreference also.

-- For the two object references or, and one & on, == one oreturns True)

i.e, Equality operation is always ment for reference address Companions



To apply Equality Operators blw the object references Compalsory

() [either parent to child (or) child to parent (or) Same type otherwise

 \bigcirc

9

0

()

 \odot

 \odot

. }

```
Ma
       Eg ? (3) :-
                  object 0, = new Object(); because object is
                                                              Super Class
                   Thosead E, = new Thosead ();
                   Storing S, = Dew Storing ("shorth");
                                                              Java. Ling
                   S.o.P(t, == Si); CE: En Compasiable types itroead &
    object
                                                           java.larg. Storing
                   Sop( h = =0,)
                   Sop(S, = = 0,); F
   -> foor any object dieference of, if on is pointing to any object
       91 == nuil is always, false, otherwise or Contains null value
            null == null & always True.
O Note .-
  * In General, == operator ment for reference Companision
      where as equals () method ment for Content Compasision.
 .
 )
                                                      (instance of) ~
                         Enstancof operator
)
9
  is By using this openation we can check, whether the given object
 )
      is to a pasiticulasi type on not.
 )
9
          $\tilde{U}_*.-
                     91 Postancof
                                                                 instanceof
                                                                  Hashtore
Strictfp
0
                                      class interface.
      any oreference type
)
En 1.
                 Shoot 5=15;
0
                  Bodean b;
€)
                    b = (s instance of Shoot)
0
                    b = (s instance of Number)
0
```

```
object
                                                                Runnable
    Egg- 1) Thoread t = Dew thread ()
              / S.o.p ( E instance of Thoread); True
              Sop(t Postance of Object);
               S.o.p (t instance of Runnable); Tome
13 To use Postangor operator, Compulsary there should be some
    Inelationship blu asisigument type, otherwise we will get Compile-
    L'ine Esision Saying Inconventable type.
           2) Thosead t = new thread();
               Sop ( t instance of Storing);
                                             C.E.
                                                Enconvertable type
                                               found : Thread
                                                 Required: Stowing
                                                                        )
is when ever we are checking parient object is of child type
                                                                       )
                                                                       -)
   Then we will get false as output.
                                                                        4
                Object 0 = New Objects Triteger (10);
             S.O.P (O instance Of Staing); false
- For any class of Potoxface & X, Dull PostanGof X always
                                                                       0
                                                                        ()
   Dietuoins false"
                                                                        0
                                                                       U
           S.o.p (null Instance of Storing); false.
                                                         else PP(0 instantit Cu)
 _g: Iterator iten = (. iteratori);
                                 Object 0 = itor next=0;
                                                          Apply customer deleted
      while (itan. basnoent())
                                 Pf (D instance of Student)
                                    Apply Student greated function
```

```
oorg:-

→ PP Both openands when Then Result is There
  Bit-wise Openations:-
   (1)
         1 -> OR -> if atleast 1 openand is T
   (§)
          -> x-oR → P Both openands ane deflement
          S.o.pln (T&T);
     -Gi-
          Soplo(TIT);
          S.o.pln (TAT); F
    EDCI)!-
           8.0.pln (4 &5);
            8.0.pln(4/5);
_)
           S.o.pln (4 15);
3
)
     We Can apply these operators Even for integral data-types also.
-)
     OKSO.
)
         Ex. a Sopla (4&5); 4
9
             @) 8.0.pln(4/5); 5
.)
٠,)
             (3) So. PINC4 15); 1
•
()
()
()
```

S.o.pln(NT); CE: Operator N Can't be applied to boolean as we can apply Bitwise Complement Operator only for integral types () but not foor boolean type. (N Torue); C.E: Operation N Can't be applied to boolean. √a) S.o.pln(N4); -5 4 ≥ 0000 0000 1111 NU = [[[]]) 2's Complement) One's Comp 000 0000 ---- 0100 add i to is comp 2's Comp 2's Comp is -ve 5 ·. -5 Note: -> The most Significant bit preparesents Sign bit. O means the no, 0 means -ve no. 0 - the no. will be sreporesented disrectly in the memory. Where as () 0 -ve nois will be Dispassented in 2's Complement from. 0

Boolean Complement Operator (!) !-

→ We Can apply these Operator only for Boolean type book.

Not for integral types.

Ex1- 0 S.O.P (!4);

C.E. operator ! Can't be applied to int.

(1) S.o.p(! False); Tonce

(3) S.o.p(! tome); False

Summany:

•

)

)

)

3

-)

)

()

Ð

٠

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0

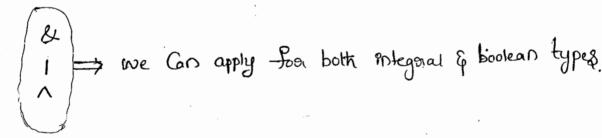
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- N => we an apply only for integral types but not for boolean types.
- ! >> We Can apply only for boolean types but not for integral types.

1) we Can use these operatoria Just to impriore performance of the System.

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2) these agre Exactly Same as nonmal bitwise operators &, 1
Except the following difference.

Ø, 1	$\alpha \alpha_1 \cdots \alpha_n \alpha_n \alpha_n \alpha_n \alpha_n \alpha_n \alpha_n \alpha_n \alpha_n \alpha_n$	
1. Both openands should	1. 2nd operand Evaluation is.	
be Evaluated always.	optional.	
2. Relatively Low-performance	2. Relatively High-performance.	
3. Applicable foor Both	3. Applicable Only foor Boolean	
Boblean & Pritegoral types	types.	
	_	

Ent. of (~ & ~ ~ & ~ ~ ~ e_2 \\

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```
1) x & & y & y will be Evaluated PFF x is Taue.
2) x | y & y coil be Evaluated PFF x is false.
```

QP!

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``,				
	1	X	20	
	81	11	17	
	1	12	16	
	11	12	15	
	&&	11	17	

```
(3)
       int & =10;
        Pf (2++2 <10) && (2/0 >10))
       S.o.pln ("-1kuo");
       else
         8.0.pln ("++;");
  Ano:
       a) C.E
        b) R.E: Asithematic Exception: 1 by Zeoro.
        c) Hello
        d) Hi
Note:
    if we Replace && with &
    Then Result PS (b), that is R.E.
                                                                 \odot
                                                                 •
```

```
Type Cast Openators:
```

- There are & types of personitive type Castings,
 - 1. Emplicit type Casting
 - 2. Explicit type Casting.

```
Implicit Type Casting:-
```

- D Compiler is the susponsible to perform this type Casting
-) a) This Type Costing is Diequiaed when Even we one assigning
- Smaller data type value to the bigger data type vascable.
- 3) It is also known as " Widening (on) UP asting".
- 4) No loss of Portonnation in this type Casting.
-) the following are various possible implicit type Casting

1B 2B

byte -> Shoot 4B 8B

folson -> float -> double

20

()

0 Chouble d=10;

[Complica Converts intoto double automatically]

/ 8.0. pln (d); 10.0

(1) Pot x = a';

[Comprises Converts Chase to lot automotion]

€ / S.o.pln(x);97

0 a=91, b=98 ---

O A = 65, B=66, C=67,

2) Explicit type Casting:-) paggammen is sesponsible to perform this Type Casting a) It is Dequired when ever we are assigning bigger data type Value to the Smalles data-type vasilable. 3) Bt is also known as "Nasisiowing on Clown Casting". 4) There may be a chance of loss of information in this Type--> The following are various possible Convertions where Explicit typeCasting Dreguined. is byte - Shoort - Int - long - float - double chan t &!. byte b = 130 C-E: possible loss of paecession found : int Required : byte byte b = (byte) 130; S.o.p(b); -126 J - when even we are assigning Bigger datatype value to the Smaller data-type variable then the most Significant bit will be lossed 1

```
47
   1) X byte b = 130;
      ~ byte b = chyte) 130;
                                         (32 - biles)
                               10000010
       130 = 0000----
     byte b = 10000000 (8 bit)
                                                 0000010
                                                   111110
                      Ja's Complement
                     1111101
              = 1×26*1×25+1×24+1×23+1×24+0×2°
              = 64+32+16+8+4+2+0
             24e126
               : -126
=
)(<u>a</u>)
     not 1=150;
     Shoot S= Cshoot);
     S.o. PIN(s) $ 150
                                    32 Pik
      150 = 0000 -- -- 0 10010110
)
     Shoat S = 0000 -- , d0010110 -- 2 Bytes = shoat = 16-hits
                           down't apply o's Comp.
<u>ر</u>َ
            +ve
                    1. S=150
_)
       int x = 150;
                                       150 = 0000 - --
       byte b = (byte) &;
                                     byk x = 10010110
0
()
                                                                1101010
        Shoat S = (shoat) x;
0
        So.pln (b); -106
0
       (S.o. pln (x); 150
                                                               = 106
```

```
10/2/11
-> when ever we are assigning floating point datatype values
  to the integeral data types by Explicit type Casting the digiets
   after the decimal point will be lossed.
   Ep1,-
           double d= 130.456;
            int a = (int)d,
            byle b = (byk)d;
             Soph (a); 130
             8.0 pln (b); -126
 Assignment Opeolatoos :-
 - There are 3 types of assignment operators
        1. Simple assignment operators
        2. chained assignment operator
        3. Compound assignment operator.
 1. Simple assignment operation:
          Epi- Pot x =10;
2. chained assignment operator:
            int a, b, c,d;
```

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a=b=c=d= 20;

```
12.48
```

-> We Carit perform chaîned assignment at the time of declaration

En! - int
$$a = b = c = d = 20$$
; } X C.E

C.E: Carit find Symbol

Symbol: Vasciable b

location: Class Test

$$a=b=c=d=20$$

3. Compound assignment operators:

Some times we Can Mix assignment operator with Some other
 operator to form Compound assignment operator.

Ex)- int
$$a = 10$$
; $a + = 30$; $a = 20 + 30$
 $a = 30$; $a = 40 + 30$
 $a = 40 + 30$

-> The following asie vasious possible Compound assignment

Operators in Java.

)

-)

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```
In Compound assignment operators the Diequired type Casting
   Will be performed automatically by the Compiler.
      byte b =10;
                             byte b = 10;
                                              byte b=10
        b = b + 1;
                               b++;
                                                 b+=1;
       S.o.pln (b),
                                               S.o.pln(b) $11
                            Sopho (b), 11
    C.E. PLP
                                               byte b=127;
      -found : int
       Dequired : byte
                                                   b += 3;
            b=b+1;
                                                S.o.pin(b); -126
        int a; b, c,d;
                                                                        a
          a=b=c=d=20;
         a += b *= c+=d/=2;
        8.0.pln (a+"----"+b+"----"+c+"----"+d);
                                                                        --}
                              600
                620
                                                                        •
Conditional Operator ( ):)
                                                                        .)
                                                                        _)
-> The only teanagy operator available in Java is a Ternagy
                                                                        \mathbf{G}
                                                                        •
   Openation (or) Conditional Openation.
                                                      a+b -> binary operator
                                                       ++a - Unagy 4
 Ep! - int a = 10, b=20;
                                                       (atn) ? a:b; - + tomory.
       int x = (a > b) & 40:50;
                                    a>b is T then 40
                                                                        asb is F then 50
        S.o.pho(8) ; 50
```

```
- nesting of Conditional operation 9s possible.
              Pot a=10, b=20;
        -!x3
               int x = (a>50) ? 777 : ((b>100) ? 888 : 999));
               8.0.pln(x); 999
       Ex!-
             int a =10 , b = 20;
               byte C= (Taue) & 40:50;
                                                 ~a<12 T
                                               × acbxCE
              byte c=(False) ? 40:50;
                                                     don't Compasse these variables
              | byte c = (axb) ? 40:50; C.E: PLP
| byte c = (axb) ? 40:50; found
                                                       found : int
                                                        Diequired : byte.
          - final int a=10, b=20;
              byte c = (a < b) ? 40 :50;
byte c = (a > b) ? 40 :50;
)
)
   Dem operation:
) -> We can use This Operator for creation of objects.
2) - En Java there is no Delete operation. because distraction of
     useless object is gresponsibility of Gastbage Collecton.
   [] operator:
  we can use these Openation for declaring & Creating arrays
```

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Operator precidence:

1. Unasy operators:

2. Anithematic Openations:

3. Shift openation! -

4. Comparision operator:

5. Equality operator:

6. Bitwise openators!

7. Shoot - Concuit openatoas:

8. Conditional operators:

9 ASSIGnment operators !-

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Evalution conden of openands:

all operands will be evaluated from left to sight.

```
Class Evaluation Osides Demo
           (2pcs [] privet 2) m. V. 2. 9
            S.o.b (w'(1) + w'(5) xw'(3) + w'(1) * w'(2) \w'(2)).
           p.s. int m, (int i)
            S.o.pln (i);
)
            Actuan i;
()
-)
     0/01.
                           1+2 *3+4 * 5/6
)
          10
)
                             1+6+4*5/6
                               1+6+20/6
                                1+6+3
                                 7+3
                                = (0
```

```
En(2) !-
          Class Test
           P.S. v. m (Staing E7 args)
            Pot x =10;
                                         int x =10;
             2=++8;
                                           x = x + +;
            S.o.pln(x); 11
                                           8.0.pln(x) 210
                                        1st place x = 10
     1st in crement
                                            1. X =10++
       and place that into x
                                      but last operation is
                                                    2=10
     int &=0;
     \Rightarrow
                                                                -
                                                                )
     int x=0;
        x+= ++8 + x++;
       8.0.pln(x);2
        x = x+ ++x + x++;
           =0+1+1
         x = 2
                                                               1
```

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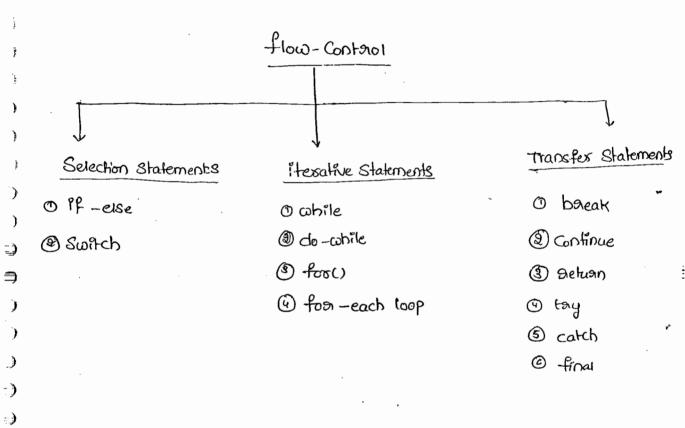
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{ }

-Flow Control :-

-> flow Contaol describes the order in which the Statements will be executed at suntime.



(a) Selection Statements:

Co if -else: -

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Spo: "F(b)

Action "F b "s tome

euse

Action "F b "s fouse

The agguement to the if Statement Should be boolean type.

if we age providing any other type we will get Compiletime Estation.

```
Exe -
    int x=0
                                   int x=10
                                                               Pot & =10;
    if (x)
                                   18(x = 20)
                                                             14 ( = = 20)
     S-o-pln(" Hello")
                                    S.o.PIn("Hello"),
                                                              S.o.pln (Heno);
                                   eise
     Stopin("Hi");
                                                             else
                                   S-0.Pln(" Hi");
                                                            8.0. pln (" Hi");
  C.E. - in Compatable types
        found: int
         Sieguired : boolean
5
                                  3
    boolean b= faisc;
                                   boolean b= faise;
    Pf (b = taue)
                                   if (b = = taue)
    2.0.pln(" He110");
                                      S-o-plac Hello");
    else
                                      clse
                                      8. o.pln ( +47);
    e.o.pln (" Hi");
```

JD !- Hi

%D1.

ttello

```
(2) Cually bases (4.4) and optional and without Charly bases are Can
       take only one Statement & which should not be declarative Statement
   Ext-
if (force)
                            if (torue)
                            S-opin(" Hello");
                            C.E.
   Switch Statement 8-
   -> 2F Severial options are possible then it is never the Commended to use
     if-else, we should go foor Switch Statement.
                Switch (x)
                  Casel:
                          Action 1;
)
)
                   Case 2:
                           Action 27
-)
                   default:
                           defaut Action;
)
()
   He Cuarty barases agre mandatory.
   -) both Case & default one Optional inside a Switch
()
        ey !-
                  int & =10,"
```

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Switch(x)

```
con default. Endependent Statements agre not allowed.
    Ep)_-
             int 2=10;
             Switch (x)
               S.o.p("Hello");
               Case, default on & expected
    until 1.44 The allowed datatypes for switch argument are
      byte
     Shoat
      70t
      Chas
But from 1.5 v onwastes in addition these the Consesponding warapper classes
   (Byte, Short, Character, Integer) & enum types are allowed.
                                                                            •
       1.4 V
                       1.50
                      Byfe
       byte
                  (F)
       Shoot
                     Shoat
                                    (A) Storing
       chasi
                     Chanacter
        int
                     Untegen
                        +
                      enum
-) if we are passing any other-type we will get Compiletime Eggo.
                                                                           0
                                                                           Ð
                                                                           0
```

-> With in the Switch, every Statement Stould be under Some Case

-)

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required : byte.

byte b=10; Chan chia; long 1 = 101; boolean b= toue, Switch (b) Switch (ch) Switch (1) Switch (b) C.E. Possible loss of precision Procompatible types -found : long -bund: boo lean required: into Dequired: PDE -> every case laber should be within the range to Switch argument type other wise we will get Compiletime Essos. eyl byte b=10; byte b=10; Switch (b) byle Switch (b+1) not type Case 10: Case 10: S.o.pin(10%) S.o ph (10); Case 100 % Case 100: S.o.pln ("100"); S.o.pln ("100"); Case 1000: Case 1000: S.0. Pln(" 1000"), S.o.pln ("600"); C.El- possible Loss of presistion found ! byte int

```
taking as vasiable as Case laber we will get Compiletime Entropy.
   B~_
           int 2=10;
           Pot y= 20;
                                                      final int 4=20;
           Switch (x)
                                             Suppose.
                                                       Case y:
                                                           S.0.p10("90");
             Case 10:
                   3.0.pln ("10");
            Case y:
                   S-0. Pln ("20"), X
               Constant Expression orequired.
   If we declare y as final then we wont to get any compiletime
  Energe
-> Expressions asse allowed for both Switch Enguement & Case laber
  but Case label should be Constant Expansion
     ep!-
           int x=10;
            Switch (x+1)
           Case 10:
                    8.0.pln(10°);
             Case 10+20:
                       S.o.pln ("10+20");
                                                                          Û
                                                                          0
```

-> every case laber should be a valid Compiletime Constant, if we are

```
-> duplicate Case labels asie not allowed.
     БÞ;
            int &=10;
           Switch (2)
              Case 97:
                      S.o. pln (4971);
              Case 98:
                        S.o.pin(4 987);
               Case 99:
8.0.pln("99");
                Case 'a':
                       5.0.pln("a");
 )
\Rightarrow
              C.F! duplicate Case label
)
•
     Summary:
)
-)
-)
                                     1 IL Should be Compile time Constant
.)
                                      -2. Supressions also allowed but should be
)
              Case laber
()
                                          Constant Expression
.)
                                      > 3. Value should be with the garge of
Ξ.)
                                          Switch aggument type
                                        4. Duplicates asie not allowed.
()
\odot
(\cdot)
```

() !

```
fall-through inside Switch:
```

with in the Switch Statement if any Case is matched from That case onwards all Statements will be executed until break Statement on End of the Switch. This is Called fall-through in inside Switch.

```
En: Shirth(x)

Case 0:

Sopin ('o');

Case 1:

Sopin ('1');

baleak;

Case 2:

Sopin ('2');

default:

Sopin ('def');

oh!

if x=0;

if x=1:

of definition definiti
```

- fall-through inside Switch is caseful to define Some Common action for Sevesial Cases,

()

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56
```

```
Ep!
                  Switch (x)
                    Case 3:
                    Calc 4:
                    Case 5.
                            S.o.pln ("Summer")
                            bacak;
                    Case 6:
                     Case 7:
                     Case 8:
                     Case 9.
                              S.o.plo ( Ranny );
                               break;
                     Casero:
                     Cale 11:
                     Case 12:
                     Case 1:
\Rightarrow
                     Case 2:
                              S. o.pln ("winter").
                               baeak;
\vec{L}
     default case :-
-)
     → We Can use default Case to define default action.
    This case will be executed iff no other case is matched
 )
     -> we can take default case any where within the Switch but it is
       Convension to take as Last case.
             Switch (x)
     Exi. -
             default: S.o.pln ("def"),
į)
              Caseo:
0
                        8-0.PIn("0");
\Theta
                         Boreak;
              Case 1:
()
                        8-0.pln ("1");
()
               Case 2: 8-0-pin (424)7
```

```
(b) Itemative Statements:
 (1) While :-
- if we don't know the no. of iterations in advance then the best
  Suitable loop is while loop.
  Out while (915. Dext) (3 while (it , has Next))
              Result Set
    3 while (e.has Mode Feltments ())
   Syntax:-
                            boolean type
              Action
to the asignement to the while loop should be boolean type.
                                                                       )
  if we agre using any other-type we will get Compiletime Earston.
                                                                       3
143
         while (1)
          S.o.plD ("Hello");
                             C.E : Procomposible types
                                   -found: int
                                                                       0
                                                                      ()
                                    Dequired: boolean
                                                                      Û
```

```
- Cually bases ase optional and without Charly bases we can
  Take only one Statement which should not be declarative statement.
 Coio
        While (torue)
                            while (tome); while (true) while (true)
int x=10;
          S-o-plo("Hello");
En@ !.
 while (four)
                                  while (false)
                                    S.o.pln ("Herro").
  Soplo ("Hello");
                                     (
S.o.pln ("Hi");
  8.0.pln ("Hi");
                                  CE:- ungreachable statement
   C. E. un greachable Statement
                                ( ) final int a=10, b=20;
      int a=10, b=20;
                                       while (axb)
       while (axb)
                                         S.o.pln (" Hello")
        S.o.pln ( Heno');
        $ .0. pln ("Hi"),
                                       1, ("itt") olg. o. 2
          Hello
                                          Unbreach a ble Statement
           Hello
```

 \Rightarrow

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.)

 $\tilde{\pi}_{n}^{(n)})$

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```
@ do-while :-
- if want to execute loop body at least once then we should go for
    do-white loop.
- Curily breases are optional & without having Evenly breases
  we can take only one Stalement blue do & while which should not be
 declarative Statement.
         do
S.o.pln("Heno"); cohile(tane); int x=10;
while(tane); cohile(tane); int x=10;

\[
\text{Mile(tane)};
\]
\[
\text{Mile(tane)};
\]
      while (true); Compulsary one statement declare (or)
                                                                                9
    X 6.E:
                             take '; "
1 do while (true)
      Sopho("Heno"); while (true)

while (false);

while (false);
     °/D1.
                                                                                0
                                                                               €
```

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()

```
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  601,-0
                                ②
                                     do
                                                           int a 20, b=20;
        do
                                                            do
        & .o. pln ("Hello").
                                    S.o.pln (" Hello");
                                                             8.0. blu ("Henos)
                                  while (false);
        while (true);
                                                             while (axb),
                                   8.0. PID ("Hi"),
    × 8.0.60("##");
                                                             S-0-Pln("H;"),
   (-G)
                                 osp:
                                          Heno
         ungreachable Statemens
                                                              0/P !-
                                          44:
   (Q)
      Int a =10, b=20;
                                     final int a =10, b = 20;
                                                                   final Pot a=10, b=20;
                                                                  do
                                    do
       8.0.pln(" Hello");
                                                                   S-o.pln ("Heno");
                                      S-opin ("Heno")
       while (a >b);
                                                                  while (a >b);
                                     while (axb);
        S.o.pln ("#");
<del>)</del>
                                                                  8-0-Pln ("Hi");
      Heno
                                                                         Hello
                                                                  0/p)-
                                 C.E. - ungreachable
       -Hi
                                                                         44:
                                                Statement
<u>-</u>
-)
9
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```

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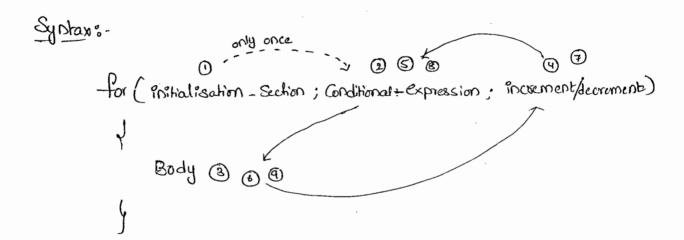
0

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foacs:

-> This is the most Commonly used loop



Only one Statement which should not be declarative Statement.

(a) initialization-Section:

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→ This will be executed only ona.

The variables in this Section.

-> there we can declare multiple variables of the Same type but different datatype variables we can't declare.

@ int i=o, byte b=o; X

@ int i=0, inti=0; X

in the initialization Section we can take any valid journ statement including S.o. puralso

```
exi.
          int 1=0;
         for ( System out print ("Hello U R Sleeping"); ix3; i++)
              S.o pln (" No Boss O only Steeping");
       OP!-
             Hello UR Sleeping
               No Boss U only sleeping
              NO BOSS
                       U only sleeping
               No Boss U only sleeping
 Conditional Expression:
-> there, we can take any java Expression but The Shessit Should
   be boolean type.
-> It is optional and if we agre not specifying then Compiler will
   always places "True".
 Incaement & decarement Section:
-> We can take any valid java Statement including S.o.p() also.
                                                                        )
 Cal.
         int i=0;
                                                                        )
        for (S.O. Pln ("Hello"); ix3; S.O. Pln ("H"))
                                                                        0
```

```
60
-> All 3 pasts of for loop are independent of each other.
-> All 3 pasts of for loop agre ophional
                                       &, Pt is Thue.
                    => Sepresent infinite loop
  Note -
         ; is a Valid Java Statement
est.
for (int i=0; true; i++)
                                                        -for (int 1=0; ; 1+4)
                         for (int i=0; false; i++)
   S.o.plo("Hello");
                               S-o pln ("Hello");
                                                           S.o. pln (" Hello"),
  CopIn (Hi);
                              8-0-pln("++");
                                                          S.o.pln ( +ti"); ~
 Sic-E!- unreachable
                             - C.E. uneseachable
                                                          L.E. unaeachable
 int a =10; b = 20;
                              final intazio; b=20;
for (int 1=0; axb; i+t)
                               for (int 1=0; axb; irt)
                                                True
   Sopin(+kno);
                                S-o-pln ("Hello");
  Sopln ("Hi");
                               · g.o.blu (, H!.);
 O/D! - Hello
                              0/po. C.E!- Unoreachable
Statement.
```

<u>(</u>-

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```
toor-each() Loop: (Enhanced for loop):
-> Introduced in 1.5V. Thes
-> This is the most Convenient loop to retrieve the elements of
   Assays & Collections
Evi- on point elements of Single dimensional Associate by using
       General & enchanad for loops
                   POEEJ a = /10, 20, 30, 40, 50/,
      -for-loop
                                            for -each
  -for (int i=0; ixa.length; i++)
                                         for (int x: a)
      Soph(acij);
                                            S.o.pln(x);
@ pount the element of ap-int Asisiay bying General & for-each loop)
                  int[][] a = //10,20,304, 140,5044;
                                              -for-each
-for (int i=0; ka. length; i++)
                                             for (intil 2 a)
  -for ("or j=0; j razi]. kryth; j++)
                                                 for (int y: x)
   S-0 pln ( a[:7[i]);
                                                  Leophoca);
                                                                       \bigcirc
```

- (i) Zt is not a Generial postpose 10019 (ii) Zt is applicable only for Astronays & Collections
- (ii) By using for-each loop we should thetoive all values of Attracts

 E Collections and Can't be used to thetoived a particular set of
 Values.

(C) Tonansfer Statements:

(1) bosear :-

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1

- -> We Can use break Statement in the following cases
- (1) within the Swith to Stop fall through
 - (2) inside loops to boreak the loop execution based on some Condition
 - (8) inside labeled blocks to boxeak that block execution based on some Condition.

P. S. v.m (-)

(

int i=10;

Li:

So pln (" Halo");

if (i = 210)

boeak li;

So pln (" Hilo");

English

Engl

```
-) If we ask using boreak Statement Any whose else we will get
  Compiletime Essonosi
        Class Test
           P-S. v.m ( -
            int x=10)
             if (x==10) X
            S.o.pln ("Hello");
                                    baeak outside Switch or loop
                                                                     )
 Continue Statement!
 → We Can use Continue Statement to Skip Current itexation and
  Continue for the next iteration inside loops
 En!
        for (ink i=0; ix=10; i++)
           if (1./2==0)
               Continue;
            So.pho(1);
→ If we agre using Continue outside of loops we win get
                                                                    •
                                                                    •
  recovered anthologonas
                                                                    0
```

Epo- int 2=10;

if (x = =10)

Continue;

S.o. pin(" Hello");

C.E!- Continue outside of loop

labeled beneak & Continue Statements:

→ 80 The Case of nested loops to break and Continue a particular loop we should go for labeled break & Continue Statements.

En: l.: .) for (- - - -)) **\$** for (----) .)) -for (- - - -) C \mathbf{C} - baear li; <u>-</u>)) -baeak la; -boreak; :) () -() \mathcal{O} 0

()

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for (90 1=0; 1×3; 1++) -for (int j=0; j<=3; j++) if (i==i) baeak; S.o.pln (1+ ----+ j); baeak:boreak (:-No output Confinue: Continue l, :-

```
do-while Vs Centinue: - (Very hot Combination)
                                                                Y 21
                                                                (h
  Cx!
         int 2 =0;
          do
                                                      2=01
            2++;
            8.0.pln(2)
           if (++x <5)
             Continue;
             2++;
            S.o.p10(2);
            ] while (++ x x 10); i
                                              10
Impgrote!
 -> Compiler will check for unseachable Statements only in the Gase
                                                                        )
  of loops but not in 'of - else'.
                                                                        Э
                                                                        9
 Epl. 0 if (tome)
                                             while (true)
                                       (2)
           S.o.pin (" Helio");
                                              6
S.o.pln("Hi");
           S.o.pln("-11;"),
                                                                        0
                                                                        ()
                                                   Ungreachable
       O/P! Heno
                                                                        0
```

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Statement

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Declarations & Access modifient

D Java Sousse File Structure (1-9)

Prograges - &

- O Class modifiens (10-14)
- 3 member modificas (15-23)
- * 4 Interfaces. (24-31)

Java Source - File Structure :-

A Java program an Contain any no. et classes but atmost one class and be declared as the public. if there is a public class the name of the program & name of public class must be matched oftenuise we will get Compiletime Eason.

Jet there is no public class then we can use any name as Java
 Source file name, there are no restouctions.

Class B

Class B

Cass C

=)

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Source Sour java (9) N Rijava (9) N Dijava .

```
Case(1):
     If there is no public class then we are use any name as
     Java Sousce file name.
  CDI_
         A. java
         Bijava -
         C-java
         Dunga-java
Cases :->
    If class B declared as Public & The program name is Ajava
 Then we will get Compiletime Essos Saying,
     Class B is public should be declassed in a file named B. java"
Case 3:-
    IF we declasse Both A&B classes as public & name of the
   Perogram is Briava then we will get compiletime Eneron Saying.
   Class A 's public Should be declassed in a file named A. java".
 Ex:
     class A
                                                                    •
       P. S. V. m(Storng[] asigs)
          S.o.pin(" A class main method");
      Class B
        p.S.v.m(Storinger angs)
           S-o. PIn (B class main method");
```

```
Class C
     d
      P. S. v.m (Storing [] args)
         Sopho ("c class man multion"),
    Class D
                                            jea.
   Save > Durga java
               Javac Duorga, java
    A. Class
                  B. Class
                            C. class
 1 Java A 4
      A class man method
⑤ java B ←1
      B class main method
3 java C LL
       c class main method
9 JONOR D +1
  A.E: No Such Meltrad Estatos: main
R-E!- NOCIOSS Deffound Enorma : Bunga
```

∌

```
Note 1.
-> It is highly 9recommended to take only one class per Source file
  & name of the file and that class name must be matched. This
  approach improves greatability of the code.
 imposit Statement:
                Test
         Class
          P·S·V·m (Stocky[] args)
            Appraylest 1 = new Appraylest ();
                                               / Asisaylish()
                                                 C-EL — symbol: meltod Association
       C-E!- Cannot find Symbol
              Symbol: Class Assaylish
               Location: Class Test
-> coe an siesolve this problem by using tally Qualified name
             Java·ukl.
-> The pooblem with usage of fully Qualified name every time increases
                                                                           )
   length of the Code & sheadability.
-> We can stessive This paroblem by using imposit Statement
         imposit gava. util. Assaylist;
         Class Test /
                                                                          0
          P-S.V.M(StatingE7 args)
```

AL 1= new AL();

-> Whene even you are Using imposit Statement it is not siequired to use fully Qualified name hence it siedocus imposoves sheadability & steducus Length of the Code.

Case (1) %-

Types of imposit Statements:-

1) There agre 2 types of Proposit Statements

- (1) Explicit class impost
- (9) Implicit class imposit

imposit Statements

Explicit clossimpost:

Epl. imposit java. util. AL;

This type of imposit is highly or the second of the second

→ Best Suitable for Hietch City
Whene Theadability is impartent

Zoplich class impost .

Exi. impost Sova. Util. *;

- → 2t is never the commended to use this type of imposit because it seduce the code.
- Best Suitable foor Amazerpet where typeing is important.

Case 21. difference blu #Producte & impost Stakment :-	
-> En clanguage # notude all the Specified beader-files will be	
loaded at the time of include statement only isomespective of coheather	
We are using the him a second of wheather	r
we age using those header-files one not. Hence This is Static bading	f
But in the Case of Sava language imposse Statement no office will loade	, , ≥d ;
at the time of impose statement, in the noon littles of Code when ever	
we are lading a class at that time only the Coornesponding class	<i>(</i>
file coil be loaded the 1 1 1 1 10 1000 1000 1000 1000 1000 1)
file coll be loaded. This type of loading is called dynamic loading or	}
load on demand on load on fly.	1
Case 3!	Ì
)
Which at the following impose Statements are valid?	. }
X 1) Proposit java util;)
X @ impost java · Util. AL. *,	9
(3) Proposition of the state of	•
3 imposib java · Util·*, "	
@ impost Javar Util. AL;	-)
Case 41.	-
·	3
-> Consider the Code, Class MyRemoteObject Extends Java. 9mi. Unic	ßle →
Remote Object	•
\ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \)
Y	
The Code Compiles fine Eventhough we agre not using imposit	\mathbf{O}
Ollman I am a market of the control	()
Statement because we used fully Qualified Name.	J
Note:	O
- when ever using fully Qualified name it is not nequined to use	Θ

Proposit Statement. When ever we are using imposit statement it is not

```
Trequisied to use fully Qualified name.
                                                            Date avoitable is both year
  Sasample:
         imposit java · util · *;
         imposit java. Sq1. *;
         Class Test
          P.S.V.m (String[] args)
             Date d = Dew Date();
                                    C.E: "Reference to Date is ambiguouse
    Note o.
         even in List Case also we will get the same ambiguity poroblem
    because 9t 9s avoilable in both Util & Sq1 packages.
)
   Casely :-
)
       impost java. util. Date;
                                               oordeon ..
• )
       impost jova. Sq1, *;
                                              ~ 0 Explict class impost
       Class Test
                                                    Classes posesent in Current
         p.s.v.m (Staing[] asgs)
                                                     Woorking Offreetony
                                               (3) emphat class emposit.
            Date d= new Date();
()
    Conclusion: While Resolveing Class names Compiles will always
     gaves the perecedence in the following onder,
```

-) Onder See above

@ Jose default package (current working discutory).

Case 9:-

Proposits Statement is totally compiletime issue if not of proposits increased then compileting will be increased automatically but of them is no effect on execution time.

Static imposit: -

- -> This Concept introduced in 1.5 Version.
- Acroading to SUN Static imposit imposite superiors shedeability of the Code, But accounting to woodle voide perogenaming expests (Like us) Static imposits steduces the greadability of the Code of coneales Conflusion, it is not seconomended to use Static imposit if theore is no specific sequirement

→ Usually we Can access Static members by using class names, but cohen even use ase using Static impost, it is not stequished to use class name and we Can access Static members disectly.

enl Without Static impost

class Test

a

p.s.v.m (Stouge) args)

Sopho(Math. Sque(41);

S.o.pln (Malts, sandom ());

S.o.pln (Math. max (10, 20));

ر ع ب with Static imposit

imposit static java. lang. math. sqst; imposit static java. lang. math. *;

Class Test

P.S.V.m (Stanfiger args)

(S.o.pln (sqorta));

S.o.pin (exandom (s))

S.o.pln (max(10,20));

Date Sql List Util But

```
* Explain about System. out. pointing:-
                                            Class System
  class Test
   & Static Strong name = xyz;
                                            Static Print Stoream Out;
   Test. name . length ();
                                            System. out. paintin ();
                  I RE is a method
                    present in Storing class
                                                            Sat is a method
                                                           present in Auntista
              Static vascable
Be is a class
                                                     It is a Static
              Present in Test class
   Dame
                                       The is a class
                                                      Vaguable of
              of the type Storing
                                       Present in
                                                      type print Stream
                                       Java. lang Package Present in System
Explanation!
-> OUT is a Static Vasiable present in System class hence
   we can access by Using classname.
-> But when ever we are using Static impost it is not required)
 to use class name we can access out voolable directly.
     imposit Static java lang. System out;
     Class Test
       p. S. v.m (Storing [] args)
        out point In ("Hello"); Hello
        out - pount in (" Hir); -Hi
```

```
-ve prospective (Ambiguity)!
```

```
Sol imposit static java lang. Integer . *;

imposit static java lang. Byte . *;

Class Test

P. S. v. m (Stoungs)

So. pln (MAX-VALUE);

CE: Seferers to MAX-VALUE in ambiguity
```

Note:Two classes Contains a vasiable on method with Same
name is Very Common Hence ambiguity problem is also Very
Common in Static impost.

while resolving static members Compiler will always gives the precedence in the following order.

- 1 Guerment class Static members
- @ Explicit Static Emposit
- 8 implice Static imposit.

```
Ex!-
   imposit Static java. lang. Integes. MAX-VALUE; -> @
    imposit Static java. lang. Byte. *;
    Class Test
      Static int MAX-VALUE =999; -> 0
     P·S·V·m (Storng[] asys)
         S.O.PIn (MAX-VALUE);
 -> If we agre Commenting Line of then Explicit Static
 împost winget porosity Hence we win get antegen class
 MAX_VALUE is % &147483647
-> 2f we gove Commenting Lines O & 1 Then Byte Class
  MAX-VALUE will be Considered & we will get 127 as 0/P.
                                                            )
                                                            .)
(-Ve Point):
                                                            :
-> Strictly Speaking usage of Class Name to access
                                                            9
 Static vasiables & methods improves treadability of the Code.
                                                           0
 Hence it is not necommended to use Static imposis
                                                           0
```

Ð

 \bigcirc

- (3) which of the following imposit Statements asse valid.
 - X (1) imposit java. lang. math. *; (we should not use * after the class).
 - X@ imposit java-lang. math. Sgot. *; (we should not use * after the method).
 - X 3 imposit Static java. lang. math;
 - imposit java lang matt;
 - VB imposit static java lang. math. x;
- 76 impost Static java. lang. math. 8901=();
- Timpost Static java lang, multi. Squt;

Dogmal Proposit Vs Static imposit:-

- of a package when even we are using general impost it is not nequired to use fully abolitied Name & we can use Short names directly.
- of a class when even we are using Static impost then it is not nequired to class name to access Static impost then we can access Static impost then we can access Static impost then it



7846 Classes are Their in Java according to 1.60

Package:

Facinge,-)
-> It is an Encapsulation mechanism to good related Cl	asseg
and interfaces into a Single module. The main puriposes of	
O To Siesolve naming Conflicts,	
@ To parovide Security to the Classes & interfaces. So	That
OUt Side Porson Can't access directly	
3 24 improves modularity of the application.	
> These is one Universally accepted Convension to name pack	ages:
Te to use internet domain name in sieverse.	
Comicicibant Joan housingloan Account domain name module Submodule Class name in sneverse	- -
Ex:-	
Package Com. duagajobs. itjobs;	
Public class Hallobs	
P· S. v.m(Storing [] angs)	
S.o.pln (" Getting Jobs is kniery easy"),) (H)

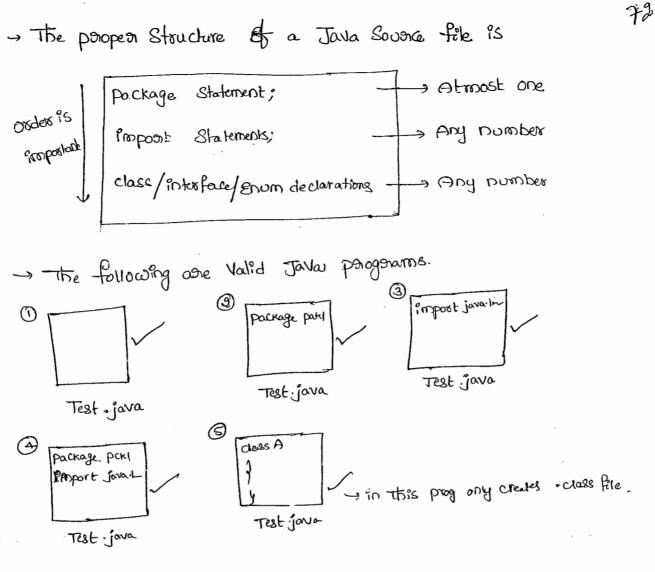
→ If the Specified destination is not available then we will get Compile — time Enough

Hydrobs. class

Ept Javac -d Z: HydJobs.java U > 28 Z: is not already available then we will get Compiletime Estatos.

```
Java Com. duagasobs. itjobs. HydJobs
  of Getting Job is Very easy.
Conclusions: -
D Pr Any Java perogenam there should be only Otmost 1 packge
  Statement. If we are taking monethan one package Statement we
  Will get Compiletine Essos.
         Package packi;
         -> package packs; <
             class A
                       C.E:- Class, interface or enum expected.
@ Zn Any Java perogenam the first non Comment Statement
  Should be package Statement (if it is available).
    CMI- ~ impost gava. Util. >+;
         - package packi;
             Class A
                Class, interface or enum Expected.
```

 \bigcirc



- An Empty Source file is a Valid Java Perogeram.

** Class modifiens **

when even you we agre cognifing over own java class Compulsar	1
We have to posovide some information about over class to the JVM	2
Like,	,
(a) Wheather over class accessable from any where or not	,
@) Wheather child class Careation is possible for over class or not.	
@ cuhealther instanciation is possible or not e.t.c.	\h_2 \h_2 \h_2 \h_2 \h_2 \h_2 \h_2 \h_2
→ we can specify this information by declaring with appropriate modifien.)
-> The only applicable modifiers for top-level Classes agre	
n Public	·)
2) <default></default>)
3) final)
4) abstract	Ć
5) Staictfp	-)
→ 2f we are using any other modifier we will get Compiletime Ensura.	-)
Saying "modifien mounted here".	<i>)</i>
Saying modifien mount not allowed here".)
Cp1 Paivate class Test	
Private class Test Private class Test)
p.s.v.m())
₹)
int x=0;	O
for (Int 9=0; 4<3; 4++)	Ð
(S.O.PIN(X); C.E!- modified powder not allowed	Ð
(S.o.Pln(2U)	0
۶ <i>)</i>	1 >

-> But for the Inner classes the following modifiens are allowed

- (1) public
- (2) < default>
- (3) final
- (4) abstract
- (5) Starictfp
- (3) Parvate
- (7) porotected
- (8) Static.

H: possuate Class A

B: of

12: y

H: Static class B

15: of

P: S. v. m (-)

S: opin (Hi);

because the man class is

note obeclare.

28/04/11

access Specifiers Vs access modifiers:

→ En old languages like C & C++ Public, private, protected & default are Considered as access Specifiers. & all the remaining like final, Static are Considered as access modifiers.

-> But in Java there is no such type of division all are Considered as access modificals.

Public classes:-

If a Class declared as the public then we can across that class from any where.

Ex:.

Package packi;

Public class A

' Public void m()

8-0.PIO("+KUO");

Jovac -d. A. Java

DOCK | A. Class

```
Package packe;
         impost Packi.A,
         Class B
          P·S·V·M (Stanger args)
               A a = new A();
               a.m.();
comp. javac -d · Bjava <
       java packe. B +
-> Of we agre not declassing class A as public, Then we will get
 Compile time Essos while Compaling B class, Saying packI.A
   is not public in pack 1 . Cont be accessed from outside Package
 default classes:-
-> 2f a class declared as default then we can access that class only
  With in that Cuspent package. i.e from outside of the package
                                                                    )
  we Can't access.
                                                                    1
```

```
- Final modifier :-
```

 \Rightarrow

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()

```
exi-
            Class P
              Public void paoperty ()
               S.o.pin(" money + Goold + Land");
             Public final void massing()
               S.o.pln (" Subba laxmi").
C.E
           Class c extends
             public void massing
              S.o.pln (" Kajai | 3sta | atora");
```

C.E!- massayes in a Cannot overside massayes in p; oversider method is final.

→ 8°P a class declassed as the final then we Gan't Coneate child class

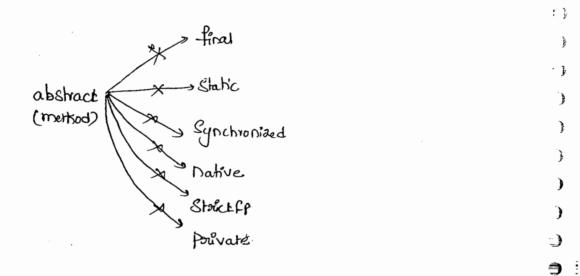
On! - final class P class C extends P

ex! - final class p	
₹	
9	
Class C extends P	ī,
. }	ì
y	,
C.E.I Carit inhealt from final p.	<i>y</i>
→ Every method present inside a final class is always final bydefault	; ;
but Every Variable present in final class need not be final.)
The main Advantage of final keywoord is we can achieve Security)
as no one is allowed to change over implementation.)
→ But the main disAdvantage of final keywood is we are missing	.) =)
Key benefits of Oop's Enheoritana ε polymosiphism (oversiding).	.)
	•
Hence, if there is no specific nequinement never necommended to)
Use final keywoord.	
abstract modifier:)
)
- abstract is the modifier applicable for classes & methods but	•
not for variables.	Ç
abstract method :-	e
unoril about)
Eventhough we don't about, implementation Still we Can declare a multiple)
coff abolical - be)
,	0
	Û
Compulsory Ends with ; ,	U

```
Exi.
          D Public abstract void milly
          (a) public abstract void m2();
   -> Child classes agre gresponsible to parovide implementation for pasient
      class abstract methods.
     Co:
             abstract class Vecticle
                public abstract int gett Noof wheels ();
            Class Bus extends Vehicle
               Public int getNoOfWheeels()
٥
                 Sietuain 6;
)
             Class Auto extends Vehicle
                Public int getNoof Wheek
                  Sietum 3;
    → By declassing abstract methods in pasient class we can define
     Guidelines to the child classes which describes the methods those
      are to be Compulsory implemented by child class.
```

.)

- → abstract modifier never tacks about implementation, if any modifier tacks about implementation than it is always illegal Combination coult obstract.
- The fonousing are Vasious illegal Combinations of modifiers for methods



abstract class:-

-> -foor any java class if we don't want instanciation then we have I to declare that class as abstract i.e., for abstract classes instancialing (coveration of object) is not possible.

Sp:- abstract class Test

Test t = Dew Test();

C.E? Test is abstract; Cannot be instantiated

Test t= pew Test(),

O

 $\mathbf{0}$

En). Storing S1 = New Storing ("dwoga");

Storing S2 = New Storing ("dwoga");

Storing S2 = "dwoga";

Storing S4 = "dwoga";

Storing S4 = "dwoga";

Say

dwoga

Say

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O. C.

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1) notify() & notifyall()	~1) Larguage fundamentals	
a) Collection & Collections	(a) Openators & Assignants	٠.,
<u></u>	3) flow-Contorol -	
3) equals() & ==	4) declaration & Access modifier	
4) Compagable & Compassatog	s) oops concepts	1
5) Staing & Storing Ruffer	Exception Handling	、) 、少
6) Stowng Ruffer & Stowng Buildern	9) Java-lang package	
Throw & Throws	10) dava. io. package	· j
8) Thorows & Thorown	11) Seovialization	
9) Hashmap & Hashtable	12) Java-Util-packaga (Collection (frame	. , , , , , , , , , , , , , , , , , , ,
10) enum, Enum, Enumeration	14) Regular Expressions	ji L
11) final, finally, finalizer	(5) 6	*
	1) IIBN =)
	18) erum -	•
	19) developement)
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oell	Chaitanyavobr@gmail Gm	.)
	Satishydwoodo 4	•
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	Chaithanya-Anumanchi@dell. Gm)
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	ON pungajobsInfo 9870807070	•
	Old boold or 251140 1 4 1000 1010)
)
)
	ϵ) ,

- If a class Contains at least one abstract method then Compulsary

 That class should be declassed as abstract otherwise over will get

 Compiletime Essavor because, the implementation is not complete & hence

 we can't oxeate an object.
- → Eventhough Cas class clossnot Contain any abstract method still ace

 Can declasse the class as abstract i.e, abstract class Con Contain Zerro

 O no of abstract method.

Ex!- HTTP Servelet, This class doesn't Contain any obstract method but Still it is declassed as abstract.

```
Class Test

O Class Test

Public Void MIC;
```

C.E:- missing method body, or declare abstract

```
② Class Test

Public abstract void mill

G.F.:- abstract methods Grit have a body
```

(3) Class Test

/

public abstract void m.();

Cit! Test is not abstract and doesn't overbride abstract method mil)

Co-41. abstract class Test	
j	,
Public abstract void mic);	
Public abstract void mach	į
G)
class Subtest extends test	÷
4	þ
public abstract void mill / 6	Je of
· •	ţ
'	•
Ct!- Subject is not abstract and does not oversude abstract	•
method mach in Test	}
C30&5	<i>)</i>
-> We can handle these compiletime Esision Citian by declassing subject))
as abstract on by providing implementation for moll.	. J
	÷ ⇒
Note:	Э
→ The usage of abstract methods, abstract class & interfaces are)
Gecommended & it is always good perogenaming peractice.	Э
9.01.32	•
abstract Vs final:	.)
)
→ abstract methods are have to oversuide in child classes to provide)
implementation, where as final methods carlt be oversudden, Hence,)
	, Э
abstract final Combination is illegal Combination for methods.)
-> for abstract classes we should coneate Child classes to pavide proper)
simplementation but for final classes ove Can't Coneate child class. Hence	-
	o
Obstract final Combination is illegal for classes,	Ð
	O

-- Final class Girls have abstract methods where as abstract class can Contain final methods.

-final class A

abstract class A

poblic final void mic);

y

y

Starctfp (an Lower case) modifier !- (Starctfloatingpoint)

)

)

- → Stoucker is the modifier applicable for methods & Classes but not for variables.
- ⇒ if a method declared as &tocctfp all floating point Caluctations in that
 Method has to follow IEEE 754 Standard So, that we will get
 Platform independent nescutts.
 - > Stouchfpes, always trailes about Proplementation where as abstract
 method never talks about Implementation. It has stouchfp-abstract method
 Combination is illegal Combination for methods.
- → 28 a Class declared as Structsp Then every <u>Concerente method</u> in

 That class has to follow IEEE 754 Standard So, that we will get

 Allabform independent shoults.
- abstract Stocktp Combination is legal for classes but illegal for meltiods

 exi- abstract Stocktp class Test

Membea (vaisciables & methods) modificas:-

```
1 public members:
```

The we declare a member as public then we can access that member from anywhere but Corresponding class should be visible (public) i.e., Before checking member visibility we have to check class visibility.

eni-

```
Package pack1;

Class A

public void mic)

2.0.pln("Hi");
```

```
Package pack2;
Emposit poek1.A;
Class B

P. S. V.m( ____)

A a = new AU;

a.m.();
```

Eventhough 10, (1) method is public, we can't access 10, (1) from outside of packs because the Coornerpording class A is not declared as public. If both asie public then only we can access.

@ default members:

That member declared as the default, then we can access that members only with in the Current package & we can't access them outside to the package. Hence, default access is also known as a package level access.

O

-)

)

O

3 parvate members:

- → 2F a member checlased as pouvate then we can access that member only within the Current class.
- → abstract methods should be visible in child classes to passide.

 Implementation where as powde methods are not visible in child

 Classes. Hence portate-abstract Combination is illegal for methods.
- 4) Parotected members: (The most misunderstood modifier in Java):
- → 2f a member declared as protected then we can access that member with in the Current package any where but outside package only in child classes.

Parotected = < default > + kids of an another package (only child Deference).

- with in the Current package we can access perotected members either by Parent selement or by child selements.
- But from outside package we can access priotected members only by using Child reference. if we are trying to use parent reference we will get C.5

Go!. Package pack1;
public class A

)

Porotected void on ()

S.o.pln ("The most misunderstood modifier in Java");

Class B entends A

P. S. v. m (____)

/ A a = new AL)	→ The most sestricted modifier Ps 'possuate'
S B b = new BU	→ The most accessible modifier
/ b.m.();	is public.
A a, = new BC) // a, m, c);	parvate < default < porotected <public></public>
Packege packes	Vasiables is powate
Propose Packi. A.	→ The DieCommended modifier for
public class c entends A J	methods is public
p-s.v.m()	
A a = Dew A()	E
× a.m,();	3
C c = new C()	
~ cm,();	•
A α, = D(w) C()	
a _i ·m _i ();	3
}	• • • • • • • • • • • • • • • • • • •
· ·	$\mathbf{\epsilon}$
Pastage 2	package3
+ Protected void mily - B extends A	D extends B
o c entends B	
I - CANTARDE	9
The most restoucted	• • • • • • • • • • • • • • • • • • •
THE THOUSE STRUCTED	Ü

1910 - 11 To 1910

•

visibility .	pouvate	<default></default>	protected	public.
0 with in the Same Class	~	1	~	~
@ form child class of Same Package	×			1
Same package	×			
(9) from child class of outside Package.	×	×	CBUE WE Should W	æ
6 from non-child class of Outside package.	×	\times	reference ×	

= Final Vasuables:-

-) -> In Generial for instance & Static variables it is not Dequired
) to perform initialization Emplicity JVM will always Provide default
) Values.
-) -> But for the local vascables JVM won't to psovide any default Values Compulsary we should provide initialization before using that) Vascable.

- Final instance Vasiables :-

()

- Initialization Explicitly JVM will powide default values.
- Use should persform instalization wheather we are using or not otherwise

we will get Compiletime Eason epol -Class Test Class Test final intx; C.E. Vascable & might have not been mitalized. Role: for the final instance vascables we should perform initialization Defore Constructor Completion. -> i.e, the following are vasious places for this, 1) At the time of declasiation ep!class Test final int 2 = 10; (3) inside instance Block. Class Test 3 inside Construction. Class Test . x=10;

→ Other Than these of we are perform énitialization any where else we will get Compiletime Esparon.

Class Test

final int x;

Public Void mil)

x=10;

C.E.I. Cannot assign a value to

final Vasiable X.

-final Static Vasciables:-

- of the normal Static Vasiables it is not sequisited to perform initialization Explicitly, Jum will always provide default values.
- \rightarrow But for final static variables we should perform initialization Explicitly Otherwise we will get C.E.

Class Test

Static int x;

Class Test

final Static int 2;

C:E! Vasuable x might not have been initialized.

Rule :

(<u>)</u>

- For the final Static Vasiables we should perform initialization Before Class loading Complition.
- i've, the following asie various places to perform this.

```
(1) At the time of declaration
         ep! class Test
                 final Static int x=10;
      Inside Static Block
      Cn!-
               Class Test
                   final static int x;
                   Static
-> If we are performing initialization any where else we will
                                                                     )
        Compiletime Estatos.
  Jek
         class Test
          final Static int x;
           Public Void mil)
               2=10;
                                           a vasiable to final
                      C.E!- Can't assign
                               vasiable &.
                                                                    0
```

```
ii) final Local Vasiables:
```

For the local vasciables JVM won't to poweride any default values Compulsary we should perform initialization before used that vasciable.

```
Octass Test

Publicisation math ()

Publicisa
```

Eventhough Local vasicable declared as the final it is not sequired to Perform initialization if we are not using that vasicable.

```
Go! Class Test

P. S. v.m()

final int x;

S.o.pln ("Hello Soi");

f

Of! Hello Sai.
```

The only applicable modifien for local variables is final. If we able using any other modifien coe will get Compiletime Estatos.

```
Public intx20; y final int x240; ~
```

```
-> formal parameters of a method Simply access as Local vasciables of that method hence, a formal parameter can be declared as final.
```

→ If we declare a formal parameter as final within the method we Cart change its value otherwise we will get Compiletime Espect.

```
exi_
               Class
                      Test
                 p.s v.m( ----)
                     m_1(10, 20);
                               P.S.V.MI (-final int x, int y)
                                          -formal parameters
                           / Can't assign a value to final vascable ac.
                  9-8000 >
                 8.0.pln (x+"---"+4)
                                                                            .)
             y
                                        Static - class level
                                         instante -> Object level
 Static modificar:
                                                                            _,}
                                                                            .)
    Static is the modifien applicable for variables & methods but not for
                                                                            •
   Classes (but innerclass Can be declared as Static).
                                                                            )
>>> If the value of a variable is varied from Object to Object theo
                                                                            )
                                                                           ()
  we should go for instance vasuable. In the ase of Enstance vasiable for
                                                                           0
* Every Object a Seperiale Copy Will be Created.
                                                                           U
-> Biff the value of a variable is Same for all Objects then we should
                                                                            Ð
 To for Static Vasciables. In the Case of Static vasciable only one copy will be
   Greated at class Level and Shane that Copy for every object of that class.
```

fixe Static variable is cooled at when was is weaked.

EN. Class Test

int 2=10;

Static int y=20;

P.S. v.m (---)

Test t, = new Test();

t, x = 888;

ti y = 999;

N 210 yzao

y=20

Test to = New Test ();

S.o.pln(6.x+"---"+ 6.4);

lo 999

for every object a Seperate copy will be

Coexted.

Static members Can be accessed from both instance & Static areas () where as instance members and be accessed only from instance area directly. i.e, from Static assea coe Carit access instance members directly atterwise We will get Compiletime Essona.)

- @ Consider the following declarations
 - I. Int x =10;
 - P. Static int x=10;

- 1. Public Static void mil
- II. Public void mil) S.o. pln(x);

S.O.PID (X);

()

()

 \Rightarrow

)

()

```
- which of the above we can take Simultaneously with in the
   Same class.
  A). I & 11
  X B). I dil . CE!- non-Static vasiable & Connot be accessed
                                   from Static Contenses
  SO ILAD
  V B)
        I & TO
 XBIUD
 X F) II & D
 -> for Static methods Compulsary implementation Should be available
  where as for abstract methods implementation should not be available Herce
  abstract-Static Combination is illegal for methods.
-> for Static methods overloading Concept is applicable Hence with in
 The Same class we can declare 2 main methods with different agreements
            Class Test
             P. S. V.M (String[] args)
                Sopln(" Staliges");
             Public Static wild main (inter args)
              S.o.pln (" PDE []");
                                                                        1)
                    %. Stacyer
                                                                       \mathbf{O}
```

→ Enhesütance Concept is applicable for Static methods including main() method hence while executing child class if the child doesnot!

Contain main method then the parent class main method will be exerte

P-S-V.m (Statige 1 args)

Q-0-pln (" pasunt class");

Class c extends p

Javac p. java

P. Class C. Class

% Java p % Java C
Patient Class pointent Class.

Duk it is not oversiding, it is method hiding.

Class P

ز

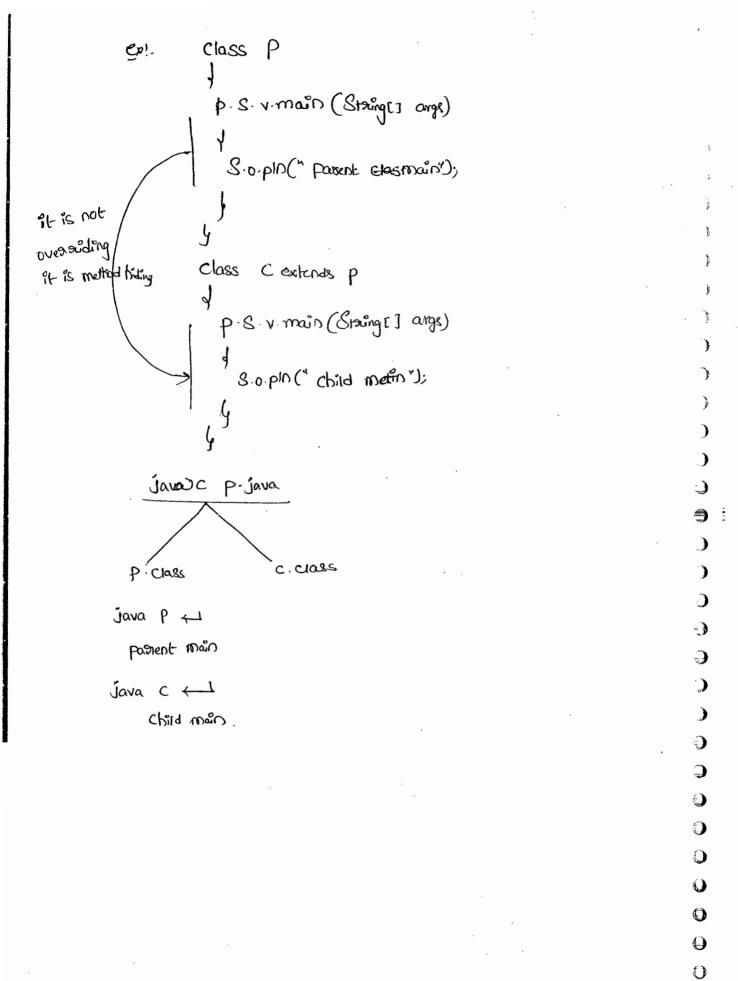
(,,

(;

()

€) >

P.S.V.m (-)



enementalist appreciation

 \mathbf{O}

- notive is the modified applicable any for methods but not for variables and classes.
- The native methods ase implemented in Some other languages like CEC++ hence native methods also known as "foreign methods."
- The main Objectives of native Keywood asset to improve performance of the System
 - 10 TO imposove performance of the System.
 - (3) To Use already existing Legacy non-Java Code.

```
PSoudo Code :. -
```

```
-> To use notive keyword
```

ex: class Native d Static

1) Load Sustem

native library | System-load Library ("native Library")

Decrare Public native void mil);

Class Child

p.8.v.m(____)

(1) 20 Notive n = Dew Nahve();

Notive metted n.m.v;

ا با با

-> For notive methods implementation is already available in other languages	
and we are not responsible to provide implementation. Hence native	
Method declaration should Compulsary Ends with ";"	
evi. O class Test	
Public Native void mi())
4 ×	
(C.E! - native methods Can't have a body.)
(3) public native void m1();)
The native methods implementation should be available in some other)
Languages where as for abstract methods implementation should not be available	ر اعاد (
Hence abstract-native Combination is illegar Combination for methods.)
Dative methods Cannot be declared with Structfp modifier because))
These no guarante that old language fallows IEEE 754 Standard.)
Hence abstruct native - Struct p Combination is illegal for methods.)
the main disadvantage of native keyworld is to it breaks platform) }
independent nature of Java. because we ask depending on Siesuit of platform)
dependent languages.	0
	0
	O
	U
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	O

alled and the malane

:

Synchroniaed modifier: :-

- → Synchronized is the modifier applicable for meltiods & Blocks.

 We can't declare class & variable with this keyword.
- ightharpoonup 3 f a method (091) Block Cleclasted as Synchronized then at a -time only one thread is allowed to operate on the given Object.
- The main advantage of Synchronized Keywoord is the Can Diesolve data in Consistancy peroblems. But the main des-Advantage of Synchronized keywood is it increases coaiting time of thread and effects performance of the System.

 Hence, If there is no specific Diequirement it is never Diecommended to use Synchronized Keywoord.

o transient modified:

- → transient is the modifier applicable only for variables & we Contempty for methods & classes.
-) At the time of Servialization, if we don't wont to Save the value of a posticular variable to meet Security Constraients, then we should go for transient keywoord.
-) -> At the time of Seovalization JVM ignories the obliginal value of transient variable & default value will be Seovalization.

○ o Volatile modifies :-

0

- > Volatile is the modifier applicable only for variables but not for methods & classes.
- O —) If the value of a vasiable keep on Changing Such type of vasiables we have to declasse with volatile modifier.

3 7f a vasilable declared as volatile then for every thread a Seperate	
locar Copy will be Concated.	
- Every intermediate modification - Enformed by that thread will takes place	
in local Copy instead of master Copy.	÷.
+> OnGe the value got finalized just before terminating the Thread the master	()
· · · · · · · · · · · · · · · · · · ·	.)
Copy value com be opdated conth local stable value.	
→ The main advantage of volatile Keywood is we an actione resolve data)
in Consistency paroblems)
But the main disadvantage of volatile Keywoord is, Escating & maintaining)
a Seperate Copy for every thread, increases Complexity of the programme	Ē
& effects performance of the System. Hence, of there is no Specific requirent) T
if is never recommended to use volatile Keywood, & It is almost)
Outdated keywoord	∌
)
Volatile variable means it's value keep on changes where as that)
Vasilable means 9ts Value Neves Changes. Hence final-volatile Combination	•
is illegal Combination for Variables.)
Conclusion:)
)
The Only appliable modifier for local variables is final	
-> The modifieous cohich are appliable only for Vascables, but not for classes &))
methods age. Volatile & transent)
-> The modifiens which ask applicable only for methods but not for classes &)
Vaguables native & Synchownisted.	
-> The modifiers which one applicable for top Level classes, methods & variables) }
ase public, idefault>, final)

1. 1.3

•

enom Gastaudas	>	>	>	> ×	X	×	×	X	×	×	×	×	
in testaces er))	X	× ×	×	\ \	X	X	X	>	×	×	
blocks	×	×	X	X	×	×	}	>	×	K	×	×.	
Variables	\	\	>	>	>	X		×	×	X	>	` `	
methods	/	\))	\)		}		>	X	×	
C. Dosso))	\	7)))	×	X)	×	×	
Outen	/	}	×	×	7	7	X	X	×	\	×	×	
Modifien	Public	<default></default>	Powake	Pstoteches	Phod	abshact	Static	Synchronized	Dahve	८१५८८१	translent	Volatile	

<u>:</u>

.-

:

The modifiers which are applicable for Inner classes but not for Outer Classes are Painate, protected, Static

Interfaces

ひ	2ntonoduction	.)
(3)	Enterface declaration & Emplementation)
(3)	(a) extends vs implements. Enterfale methods	•
(4)	anterface Vascables) (
(5)	Enterface Naming Conflicts)
	(1) method naming Conflicts)
	(9) Vadiable u 11	Š
(6)	masikesi Enterfac	3
	Adapter class)
	Abstract class vs Concerete class vs interface.	•) •
(9)	diff. blw abstract class & interface	e e
		.)
		6 c
		. .
		9
		· ·
		O
		O
		e O
4	:	•

Interface:

- De Any Service Dequirement Specification (SRS) is Considered as Enterface.
- from the client point of view Gen an Interface defines the Set of Schwigs what is expecting.
- -> forom the Seavice porovides point of view an interface defines the Set of Seavices what is offering.
- Hence an Interface Considered as Contract blw Client & Sequise provider Ex.
 - By using Bank ATM GUI Scaren, Bank people will hightlate the Set of Seavices what they are offering At the Same time the Same Scaren describes the Set of Seavices what End-Usean is Expected.
- Hence this GUI Scareen acts as Contack blue the bank people & customers
 - That in the Interface we Can't ownite any implementation, because it has to highlight Just the Set of Senvices what we are offering or what you are Expecting. Hence every method present inside interface. Should be abstract. Due to this interface is Considered as 100% pure abstract class

Vuhat is an Interface:

- Any Service requirement Specification (SRS) (00) Any Contract blu

 Offent & Service provider (00) 100% pure abstract class is nothing

 but an Interface.
 - -> The main Advantages of Enterfaces asse.

- (i) coe an acheive Secusity, because we are not highlighting our internal implementation.
- (ii) Enhangment will belome Very Casy, because with out effecting outside person we can change our internal implementation.
- (iii) Two different Systems Con Communicate via Interface.

 (A Java application Can talk with Main France System through Interface).

Declaration & Implementation of an Interface:-

Implement an Enterface by using Interface keywood, we an

Col. interface Interf

Void m1(); / by default public abstract void m1();

Void m2();

abstract class Seavice Provider implements Interf

-> Public void m1()

Af a class implements an intersface Compulsary we should provide in mplementation for every method of that interface otherwise we have to declare class as abstract. Violation leads to Compile-time Espron.

0

-> when ever we are implementing an interface method Compulsory it should be declared as public otherwise we will get Compiletime Exercise

Extends Vs implements:

- 1. A class can extend only one class at a time.
- a. A class an implement any no. to intexfaces at a time.
- 3. A class an extend a class and an implement any no. of interfaces Simultanuiously.
- A. An Enterface Can extend any no. of interfaces at a time.

Exi: intextac A

f

y

intextac B

f

y

intextac C extends A, B

 \Rightarrow

- a) which of the following is Torue?
- (1) A class can extend any no. of classes at a time. X
- (2) A class Can implement only one Interface at a time. X.
- (3) A class Can extend a class and an implemptate an interface but not both Smultaneously x
- (4) An Enterface Can extend only one Poterface air a time X
- (s) An Enterface can implement any nort classes at a time X
 - (6) none of the above ~

Q) Consider the expression	
X extends 4 for which of the following possibilities	1
This Exponession is Thue?	<u>ļ</u>
10 Both should be classes	2
@ Both should be in kofaces) }
8 Both Can be either classes or interfaces	ر الا
)
4 No Restauction.	•
) ()
O X extends 4, Z)
(e) X, Y, Z should be intexfaces)
② X extends y implements Z	() (<u>)</u>
×, y -> classes	9
Z - noterfaces	•
3 X Porplements y extends Z))
$C \cdot \epsilon$	9
	•
))
	•
	9
	9 0
	o o
	Q
	9
	Ú
	Ú

100

•

Interface methods:

wheather we are declaring or not, every interface method is by - default, public & abstract

en!- interface Interf d Void 100, ();

abstract:

Belause interface methods Specifies requirements but not implementation.

Hence the following method declarations are equal inside

(1) void mill;

- (2) public void mi();
- (3) abstract void m1(); ~
- (4) Public abstract void mil); ~

of As every interface method is by default public & abstract the following modifies are not applicable for interface methods.

- (1) parvate
- (3) paotected
- (B) <default>
 - (4) final

- (5) Static
- (6) Stack fp
 - (3) Synchronized (2)
 - (8) native

```
-> which of the following method declaration are valied inside interfact
  (1) public void mic) 16 x
  (9) public Static void mill; x
  (3) public Synchronized void m, (1; X
  (9) pourate abstract void mill; X
  (5) public abstract void mil);
interface variables:
 -> An interface can Contain vasuables the main purpose of these "
  Variables is to Specify.
 Constants at Dequiarment Level:
 -> Every interface variable is always public, static, final cohetter
  we are declaring (or) not.
     intesface Enter
      int & =10;
  Public: To make this variable available for every implementation
  Static !- without existing object also implementation closs can access
         this variable.
  final: implementation class Can access this variable but Can't modify
  -> Hera înside interface The following declaration are valid & equal.
                                 4) public static final intacio;
                                                                       ()
    D int &=10;
                                                                       ()
                                 5) public Static "int x=10;
    2) public intacio;
                                                                       4
                                 6) final 9pt 8/210;
    of public static int x 210;
                                                                       0
                                 # public final interclo;
```

- B) Static final int x=10;
- As stressage variables are public static & final we Got declare with the following modifiers.
- (1) posivate (3) <default> (5) volatile.
- (a) ponotected (y) transient
- -> for the interface variable Compulsary are should perform initialization at the time of declaration only otherwise will get Compile time Estador.
- Phterface Interf

int a; X C.E:- = Expected.

) -> which of the following vasuable declasiations are allowed inside) interface.

(1) Pot x=10;

(5) transient int x=10; X

(2) Pota; x

- (6) volatile int x=10; X
- (3) paivate int x=10;X
- (1) public static final int x=10; ~

(public sot x=10; ~

(,;

-)

 \mathcal{C}

)

- -

· .

-> Enside implementation Classes we can acress intextac vasuables but we Can't Modify These values. Go! interface Interf int & =10; Class Test implements Class Test implements Interf Interf P·S·v·m (Starge7 args) p.S.v.m (String[] args) x=888; int x = 88; S.o.phn(x); 8.0.pln(4);88 Ciej. 20 kosface naming Conflicts: 1 method naming Conflicts: • Casel: -**)** → 2f Two interfaces Contains a method with Same Signature & Same netuan type in the implementation class we can provide implementation for Э \mathbf{j} only one method. Û ۔ اور interface Left 0 interface Right € Public void miles; Public void milly 0

```
Class Test implements Left, Right

Public void mil)
```

Case 2:

→ 2f Two interfaces Contains a method with Same name but differentange then, in the implementation class we have to possible implementation for both methods & these methods are Considered as overloaded methods.

```
Co!- interface Left interface Right

Public void mi();

public void mi(int i);
```

Class Test implements Left, Right

Public void mi()

Public void mi(int i)

methods

__)

() ()

4

→ ZF Two intexfaces Contains a method with Same Signature but different setupintypes. Then it is impossible to implement both intexfaces at a time.

Ex!
interface Left | interface Right

Ep)interface Left interface Right

Public void m1(); public int m1();

-> We Cont woulte any Java Class which implements both interfaces Simultanewy.

interfaces Simultaneously.

#) yes, Except 24 Two intextaces Contains a method with Same Signature but different Dietan types.

)

()

٦

(2) Vascable naming Conflicts:

intexfac Left

int x=888;

9 not x = 999;

interface Right

1

```
Class Test implements Left, Right

P. S. v.m (-----)

d

Sopln(x);
```

C.E:- sieference to a is abéquaces.

There may be a chance of a Interfaces Contains available with Same name & may sise vasiable naming conflicts But we can sessive these naming conflicts by using Interface names.

Sop(Left.x); 888 Sop(Right.x); 999

Masikesi Poterfale :-

<u>_</u>)

:)

()

()

()

()

Ex: Kenya

-> 2f an interface wont Contain any method & by implementing that interface if own objects will yet ability Such type of interfaces are Called Masikes interface (or). Tag interface (or) ability interface.

En. Seaializable, Clonable, Random Access, Single Thread Modie.

- These interfaces are marked from some ability.
- En!- By Proplementing Seonalizable interface we an Send Objects.

 across The N/w and we an Save. State of Objects to a file.

 This extra ability is provided through Seonalize intexface.

```
# By implementing Cloneable interface our Object will be
    in a position to poovide exactly duplicate Objects
9) Masikes interface count Contain any method then how the Objects
  will get that Special ability?
     JVM is sesponsible to paovide sequined ability in maskes
    interfaces.
   why TVM is peroviding enequined ability in marker interface?)
     To Seduce Complexity of the pozogonaming.
Q) Is it possible to Careate Our Own Marken Enterface?
    Yes, But Customization of JVM is acquired.
 Ex: Sleepable, Estable, Jumpable, Lovable, Funnable.
                                                                    )
Adapteon Class :-
 - Adapteon class is a Simple java class that implements an
                                                                    -)
                                                                    .)
 interface, an interface only with Empty implementation.
                                                                    )
                                                                    •
                          abstrack class Adapter X implements X
     interface x
                                                                    )
                                                                    1
                                             If we coneate an object
                            mico dy
       mio;
                                                                    Э
                                             for this Empty 9 result
                            wall gh
       mach
                                             So for this class side
                                                                    ()
                                              declare as abstract.
                             missoil 44
                                                                    0
       m1000 ();
                                              by default abstract
                                                                    \Theta
```

-> If we implement an interface directly consulsary we Should provide implementation for every method of that interface, whether we are introded or not & whether it is required (or) not. It increases length of the Code, so that readability will be reduced.

Class Test implements X

min of b

If we extends adapted class instead of implementation interface disectly then we have to possible implementation of only for required method but not all this apponach reduce Length of the Code & improves needability.

>> Class Tesk extends Adapter x

d

mu()

d

Concerete class Vs abstract class Vs interface :

→ we don't know any thing about implementation Just we have nequinements Specification, then we should go for interface Ext. Seavlet.

()

()

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-> we agre talking about implementation but not completly

(Just partially implementation) Then we should go for abstract

Class.

En: Generic - Seavlet

HTTP_Seavlet

-> We ask talking about implementation Completly & sheady to posovide Seswice, Then we should go foor Concrete Class.

Ext. Ovar won Seawlet.

Défference b/w interfaces & abstract class:-

D If we don't know any thing about implementation Just we have requirement specification. Then we should go for interface.

intesface

- 2) Every method present inside interface is by default public & abstract.
- 3) The following modifiers are not allowed for interface methods:

Stockfp, porotected, Static, native pourvate, final, Synchronized,

abstract class

- implementation but not Complety

 Chartially implementation) Then

 we should go for abstract class.
- 2) every method present inside of abstract class need not be public of abstract. We can take concrete of methods also.
- 3) There are no restrictions for U

 Obstract class method modified U

 i.e., we can use any modified.

 $\mathbf{0}$

- 4) every variable present

 Proside interface is public, Static

 Final, by default wheather we are

 declare (201) not
- 5) for the interface variables
 coe Grit declare the following modifier
 private, protected, transient, volatile
- 6) for the interface variables

 Compulsary we should perform

 Philialization at the time of declaration

 Only
-) 7) Inside interface we ant take) instance & Static blocks.
 - 8) Inside Interface we cart take Construction.

- 4) abstract class variables need not be public, final Static.
 - 5) There are no restriction for abstract class variable modifieds.
 - 6) for the abstract class variables
 There is no restriction like performing
 initialization at the time of
 declaration
 - 7) Enside abstract class we Can take Static block & instance blocks.
 - 8) Inside abstract class we an take Constructor.

(Q)	Znsid	le abst	ract	class	we	<u>an</u>	take	Constructor	but	
								s, what is		

- De abstract Class Construction will be executed whenever we agree Coneate Child class Object to perform initialization of parent class instance variable at parent Level only and this Construction meant for child object Coreation only
- as in abstract class also we can take only abstract methods then what is the need of interface.

•

)

)

•

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Û

- A) ... Interface purpose we can replace abstract class but it is not a good perogenaming peractice we are miss using the role of abstract class.
 - we should boung abstract class into the picture whenever coe asie talking about implementation.

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Oops Grapt

- Data Hiding 2
- 9) Abstraction 2
- 3) Encapsulation 2
- 4) Tightly Encapsulated class 3
- 5) Is-A Relationship 3
- 6) Has-A Relationship 5
- 7) Method Signature 6
- * 8) Over loading =
- 9) Over riding "

 \Rightarrow

:)

.)

- to) Method hiding 14
- 1) Static Control from 18
- 12) Enstance Control from . 22
-) 13) Constructions 24
 - 14) Coupling us
 - 15) Cohesion 43
 - 10 Type-Casting -40

polymorphism - 17

Type = cashing = 40

```
1 Data - Hiding:
 -> Hiding of the data, So that out Side peason Cant access own data
   distrectly.
 -> By using portrate modifier we can implement Data Hiding.
    En!.
          Class
                  Account
             Polivate double balance = 1000;
- The main Advantage of Data Hiding is we Can acheive Secusity
)-Abstraction :-
 -> - Hiding internal implementation details & just highlyate The Set
                                                                          3
                                                                          of Services what we are observing, is called "Abstraction.".
                                                                           )
  ep!_
 By Bank ATM machine, Bank people will highlate the Set of Services )
    what they are offering without Highlating internal implementation.
    This Concept is nothing but Abstraction.
                                                                          9
  → By Using interfaces & abstract classes we Can acheive abstraction.
  -> The main Advantages of Abstraction age.
                                                                          0
      1) We can acheive Security as no one is allowed to Know own
                                                                          0
                                                                          \mathbf{0}
        Potesinal implementation.
                                                                          ()
     2) With out effecting outside person we Can change over internal
                                                                          ()
       implementation Hence Enhancement will become Very Easy.
                                                                          €
```

-> The main disadvantage of Encapsulation is it in Coreases The Length	
of the Code & Slows down Execution.	
Tightly Encapsolated class:	
→ A Class is Said to be tightly encapsulated iff every data member declared as the polivate.	;
wheather the class Contains getter & Setter methods are not &)
wheather those methods declared as public or not these are not)
Dequired to check.	}
ex:- Class A)
Z Cass A	()
Portale int balance;	\odot
Public int gela Balance U) (*)
4)
oretuon balanacs;	•)
· · · · · · · · · · · · · · · · · · ·	•
,	9
En: which of the following classes ase Tightly Encapsulated.)
Class A	J
~ 4	9
porvate int x =10;)
Class R subode A	9
Class B extends A of the y=20;	0
10 t A=80;	O
, Class c entends A	0
Class C entends A Posivate int 2=30;	O
bb Pouvate int 2=30;	0

```
3) It improves modulatily of the application meaning?
 3) Encapsulation :-
 -> Encapsulating data & cossesponding methods (behavious) into a Single
    module is Called "Encapsulation".
 → 2f any Java class fallows Data Hiding & Abstraction such type
    of class is Said to Encapsolated class.
             Encapsulation = Data Hiding + Abstraction
    Ewi.
             Class
                    Account
              Parivate double balance;
                                                                welcome durge
               Public double get Balance ()
                                                                 Gel- Balance
                  / Validate Useon
                                                                 Withdraw
                  neturn balance;
                                                                  GUI Scoon
               Public Void SetBalance (double
                                               balance)
                   // Validate Uses
                   this balance = balance;
  -> - Hiding data behind methods is the Centeral Concept of Encapsulating
-> The main advantages of Encapsulation are 10 we can acheive Security.
  @ Enchangment corn become Very Casy.
  1 imposives modulability of the application.
```

```
ex3! Which of the following classes are Tightly Encapsulated.
```

```
Class A

on by default what modifies to variables?

N

int x=10;

Class B extends A

postvate int y=20;

Class C extends A

postvate int z=30;

y
```

Conclusion ?-

a

→ 2f pasient class is not tightly Encapsulated then no child class is Tightly Encapsulated.

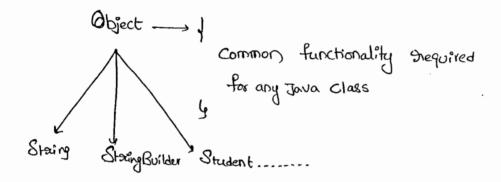
```
5) IS-A Relationship :-
```

- -> RE is also known as Inheritance
- -> By using extends Keywood we Can implement Is-A Relationship
 - -) The main advantage of Is-A Relationship is Reusability of the Code.

```
Class Test
       P. S. V. M (String (1 args)
              P P = Dew P();
                 P. m. ();
                 P. Ma(); X -> c.e :- Cannot find Symbol
                                       Symbol: method m2()
                                        location: class P
    Case 2!
             C c = Dew C();
                 c.m,U;
                 C.mg(); ~
   Cases !.
             P PI = Dew C Cs;
                 Pingo;
                  Parmau, X JCE.
    Case 41.
             C C, = new P C); X C.E. in Complable types
                                          found : P
Conclusion (1)!
                                           Stequired . C
                                                                   ാ
What even the pasient class has by default available to
                                                                   •
                                                                   9
   The child Hence and child class steference both Can Call both
                                                                   .
   Passent & child class methods.
                                                                   Ð
                                                                   ()
@ what even the child has by default not available to the powent
                                                                   0
                                                                   0
  hence on the Pasient class she fevence use Can Can only pasent
                                                                   O
 Class methods & we Cant call child Specific methods.
                                                                   \Theta
                                                                   0
```

- B pasient class sufference can be used to hold child class Objects by using that sufference we can call only pasient class methods but we can't call child specific methods.
- 1 We Can't use child class sufference to hold pasient class Objects.

En:
1) The Common functionality which is shequipped for any java classes is defined in Object class and by keeping that class as super class it's functionality by default available to every Java classes.



Epist the Common functionality which is snequioned for all Exceptions & Errors is defined in Throwable class as Throwable is parent for all Exceptions & Essavers, 2tis functionality will be available automatically to every child not slequired to shewrite. a) to 'Throwable' has 'object' as parent class?

Java won't parovide suppart for multiple inheritance but through interfaces it is possible.

Class A extends B, C

But

interfac A extends B C

C.E.:

(

:)

-> Every class in Java is the child class of Object.	
-> 2f own class doesn't extend any other class then only it is the	
discet child class of Object.	
Cp!- Class A Object = A	٠.
y	-
→ 8° Oug Class extend any other class then oug class is not	
dispectly child class of Object. Object	
ept- class A extends B	٠
d _	j
(mostificate interchance)	-1 15
→ Cyclic inhesitance is not allowed in Java	ت و-
CACILI HIMESTIANA 13 THE CHICAGO III SOLVA	7)
Class A extends B	j
A	
Class B entends A X)
d)
C.E? - Cyclic inhesitance involving A	-)
J.)
(1) Class A entends A	0
J Chas FA Chickes 14	0
	U
J ,	0
	U
	_

121 - 131 - 138

U

6) Has - A Relationship:-

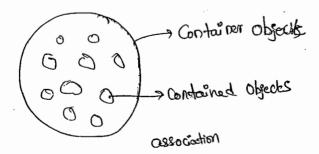
- -> Has-A Relationship is also known as Composition on Aggregation."
- There is no specific Keywood to implement Has-A Relationship the mostly we are using new keywood.
- -> The main advantage of Has-A Relationship is Reusability or (code Reusability)

Class Gog has Engine Deference.

The main disadvantage of Has-A Relationship is it increases depending blue the classes and Caeates maintains possilems.

Composition Vs Aggoregation:

→ En the Case of Composition cohenever Container Objects is destroyed all Contained Objects will be destroyed automatically. i.e., without Existing Container Object There is no chance of existing Contained Object there is no chance of existing Contained Object the Contained Object having Strong association



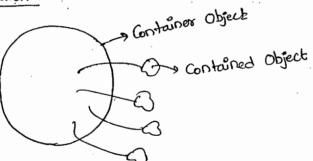
EN 8.						
\Rightarrow	university	is	Compossed	of	Sevenal	departments

whenever your are closing university automatically all departments will be closed the relationship blw University Object & department object is strong association which is notting but Composition.

- Aggoregation ; <

→ Whenever Container Object destroyed, There is no governmenty of Clestruction of Contained Objects i.e., without Existing Container Container Object there may be a chance of Existing Contained Object i.e., Container Object Just maintains References of Contained Objects.

This Delationship is Called Weak association which is nothing but —Aggregation.



Exi:
Several proofices will work in the department

Whenever one are closing the department Still there may be a Chance of existing parolesess. The Stellationship blu department & profession

is Called weak association which is nothing but Aggregation.

Ð

1)

```
public void mi (int i)
     2.0.pln (" int-arg");
   public void mi (floot F)
      8-0-pln (" float - arg");
P. S. V.m (____)
 Test & = new Fest ();
      t. m1(); / no-arg
      t. m, (10) / / int-arg
      t. m. (1058); // Ploat - arg
```

En overloading method sesolution always takes case by Compiler based on seference type. Hence overloading is also Considered as Compiletime polymosphism (ex) Static polymosphism (e) Carley Binding In Overloading Sufference type will play very important stole & Suntime Object will be during.

Casel :-

()

(

"Automatic premmotion in overloading:

To overloading method resolution, if the Matched method with Specified assignment type is not available then Compiler work suise

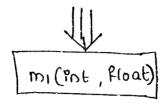
```
any Esson immediately. Fish it promotes that assignement
   to the next Level and checkes for matched method.
-> If the matched method is available then it will be Considered and
  if it is not available then Compiler once again promoted this as aguement
  to The next Level.
-> This process will be Continued untill all possible promotions after
  Completing all promotions Still it the matched method is not available
  Then only we will get C.E.
-> Thes form's Cared Automatic paramotion in overboading.
- The following are various possible pramotions in Overloading.
      byte -shootk
                         ent --- long --- float --- double
             chasi
Casel?-
2×1-
        Class Test
          Public wold mon (int i)
             S.o.pln(" not-arg");
          public void m, (flook f)
             8:0. Pln (" froat -avg"),
          P S. V. m (String [] args)
                                                                       Ð
                                                                       1
             Test t = new Test();
```

Method Signatuae :-

()

-> method Signature Consists of name of the method & assignment.
List,.

Ep: Public void mi (int i, float f)



-> En Java Sieturn type is not past of method signature.

-> Compile will always use method signature while heasolving method Cans

-> With in the Same Class Two methods with the Same Signature hot allowed Otherwise we will get Compiletime Espaces.

Col Class Test ∌ .) (mi (int) is the method signature) public void milint i) \bigcirc \cdot -) Public int milling i)) Test to = new Test() Dieturn 10; t, m, (10); C.E! - m, (int) has already defined in Test

Overloading

```
Overloading:
     Two methods are Said to Overloaded iff method names are same
     but assignements are different.
 -> Lack of overloading in'c' increases Complexicity of the program.
    In C, language : I there is a change in method assignment type.
    Compusary we should go for new method name.
          expl.
                 absco
                 labs c)
                        ----- long
                 fabs () -> float
-> But in Java -1700 methods having the Same name with different
  as guernents is allowed & these methods one Considered as overloaded
                                                                          .)
  meltrods.
                                                                          )
            epl-
                 abs ( int)
                                                                          9
                 abs ( long)
                  abs (float)
-> Having overloading Concept in Java Simplifies The perograming
    En!
                 Class Test
                                                                          •
                                                                          ()
                                                                          0
                    Public void mi()
                      8.0.pln ( no -argy);
                                                                          0
                                                                          ()
```

 \bigcirc

```
104
 Gose! - In Overloading mosts specific version will get highest povoquity.
                                          what does it mean?
   Case & 3 -
         ex!
                   Class Test
                      Public Void m, (Stocky Buffer Sb)
                          S-o.pln C" Stady Buffer - a sing");
                     Public void mi (Storing s)
                            S.o.pln ("s Ading -version");
                         Public . S. v.m (
 By default Exirg
     eurit ( tom ( new SB ( "dwga")); / Stowng Buffer-angs

et type

et type

integral constant of "the type

integral planting literal"

A L. M. ( "dwga"); // Stown

Abouting literal"

Xo L.
constant of string closs
Dra. Erkefal carstant of Ent
                        Xp E·m, (null); p // C.E!- seference mills is
                                                                   ambiguity.
                                                                                                      U
                                                                                                     0
```

```
E.m, ('a'); // int-arg
                     t.m, (101); / float -arg
                      t. m, (105); x c.e.
                                           Cannot find Symbol
                                          Symbol: melthod m; (doolde)
                                          location: class Test
  * Gse 2:-
  → In overloading method resolution child-argument get more privily
     Than parrant arrowment.
        801.
            Class Test
             Public void mi (Object o)
3
               S-o.pln("Object version");
             Public void mi (Storing s)
       3
                 S.o pin(" Staing Versions");
)
                                                          Object
             P. 8. v.m (-)
                                                         Storing
               Test t = new test ();
                 t.m, ( new Object()), / object-vension
                                                         SUPPOSE @ Statement takes //
                 t.m. ("dwarga"); // Stacing-version (Stacing the off is Objects
                 to micholi); & altoing - resiston
```

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()

- Hence Overseidling is also Known as "Trunking playmosuphism (07)

 Cynamic polymosuphism (09) Late binding".
- Oversuiding method Diesolution is also known as Dynamic method dispatch.

Rules for Overroudding:

- 1.e, method Signatures must be matched.
- De overseiding seturn type must be matched. But this stule is applicable until 1.4 version. from 1.5 version and according to this, child method stetuen types as parent method seturn type. its child classes also allowed.

```
Class P

Deblic Object mill;

Suctum null;

Class C excends p

Public Stocing mill

At is valid in 1.50,

Stetum null;

But invalid in 1.40
```

Spring double

Object Number String double

Child multiple String Theger Object int

- → Co-vaorient oneturn type Concept is applicable only for object type but not foor porimitive types.
- @ we Can't oversuide parente class final method. But we can use it as it is
- (4) poivate meltods are not visible in Child classes there oversaidding Concept is not applicable for posvate meltods.
- ⑤→ Based on Ovor Grequessement we can devere the Same

 Parent class portrate method in Child class also it is valid

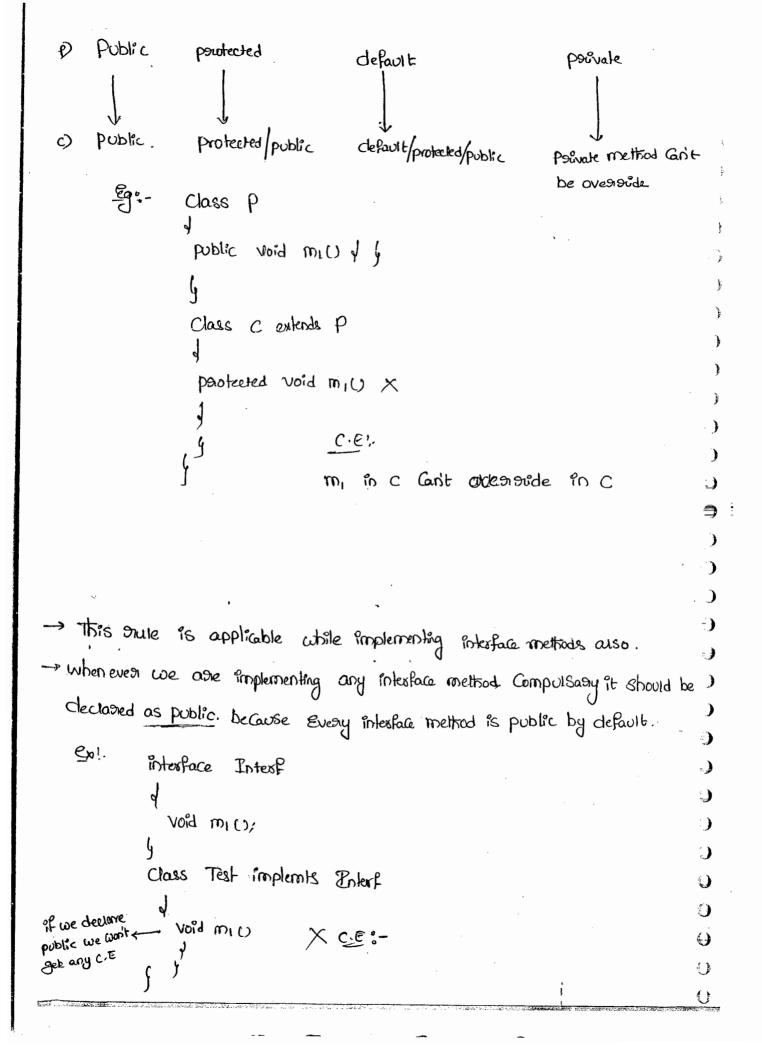
 bull it is not oversiding.

- for parient class abstract methods we should overide in thild class to perovide implementation.
- We can Overside Parsent Class non-abstract method as abstract in Child class to Stop parent class method implementation ovailability to the child classes.

```
Cp1 -
                 Class P
                   Public void PC)
                abstract class c extends p
                   Public abstract void pc);
   -> The following modifiers won't play any sestocitions
     Oversiding
                 1) native.
                 @ Synchronized
                 1 Struct Pp
                                               Synchroniaed
                                                                  native
                   non-final
                               abstract
                                                                 non-native
  Col non-final
                                               non-Synchroniaes
                     fina
                                non-abstract
        Statep
       non-Strictfp
      while oversiding we ant decrease scope of the modifier
\bigcirc
      but we an increase the following are various acceptable oversidings
```

Parivate < default < parotected < public

0



```
Pf child class method Throws Some Checked Exception Then Compulsary
     Parsent Class method should throw the Same Checked Exception on its
           class Exception.
     Passent, otherwise we will get C.E.
   -> BUE There is no Sile for unchecked Exaption.
   Ex:0
            Class
               Public void mic)
             Class C extends P
               Public void mich throws Exception
_)
                             FIFT
                                     MICO in C Cart overside mIl) in p.
Э
                                Oversidden method doesnot throw Exaption.
_)
   En@:
         P: Public void mi() throws
                                          IDEXception
)
          C: Public void mil)
)
          P: public void mic)
. )
          S: Public void mil) throws
                                         TOEXCEPTION
.)
ું)
         P: public void mil throws
          C: public void mIL) throws IDEXCEPTION
         P: public void MIC) throws to Exaption
         C: Public Void , mil) throws
```

```
P: public void mil throws IDEXCEPTION
 6
       C: Public void mics throws fileNotfound Exception, EOF Exception
       p. public void mil throws lockaption
 6
       C: public void mil throws EOFException, Protonoupted Exception
X
(7)
     P: public wid mic) throws JOEXCEPTION
     C: Public void mic) throws AE, NPE
      p : public void MIU
     C: Public void mic) throws AE, NPE
Oversiding w. s. t Static method :-
-> We Can't overlaide a Static method as non-static.
          Class
   ON! -
             Public Static Void mil)
                                                       Static
                                                        non-Static
           Class Centends P
             public void mi()
                                                                     0
                                                                     O
                                   mics is Can't overall mich in p;
                           C.El.
                                                                     0
                                   oversiden method is Static.
                                                                     \bigcirc
                                                                     \Theta
                                                                     \bigcirc
```

 \mathbf{O}

- -> Similarity, we Can't overside non-Static method as static
- → 2f both pasient & child Class method class asie Static Then we wont to get any C.E it Seems to be oversuiding is happen, but it is not oversuiding. It is "Method Hiding".

Public Static void mill)

Public Static void mill)

Class C extends P

Oversaiding

Public Static void mill)

Ptiding

Public Static void mill)

Method Hiding :-

→ All Shules of Method Hiding are Exactly Same as Oversolding Except the following difference.

method hiding

OVerrouding

- 1) Both methods Should be Static
-) 2) Meltod Die Solution takes Care
 by Compileon based on Reference
 type.
 - 3) It is Considered as Compiletime Polymosphism on Static polymorphism On early Binding

- 1) Both methods should be non-Static
- 2) Metho Diesolution always takes Care by Jum based on Runtime Object.
 - 3) 8th is Considered as Runtime
 Polymosophism on dynamic polymosophism
 or Lake Binding.

```
En!.
              Class P
               Public Static void MI()
                 S.o.pin(" parent");
             Closs C extends P
welfor Herg
                Public Static void mIC)
                  So.pin (" child");
            Class Test
              p. s. v.m ( ---)
                P P= New P();
                   p.mi(); -> porrent
                    c= new C();
                     c.miUi - child
                 P P, = new · C ();
                   P. mill; parent
                                                                         •
                                                                         \bigcirc
```

```
109
```

→ If both methods are non-Static then it will become oversolding in this

Case the olp is: Postent

Child

Child

Oversiding w. or. t Van-any methods:

 \rightarrow (we Gn't overloade a Var-arg method with general method. If we are taying to overload it will become overloading but not overloading. \rightarrow A var-arg method should be overloaden with blar-arg method only.

```
Class P

Poblic void mi (int... i)

Soplin (" pasient");

bute not bute not poblic void mi (int i)

Poblic void mi (int i)

Soplin (" child");
```

:_)

)

 \bigcirc

C c= new Cc); c.m, Clo); // lehald Prenew Co;
Prenew Co;
Prenew Co;

```
→ 28 both pagent & child class methods agre Vagraged Then it will becomes oversiding in this Case of is pasent child persent oversiding co.gr. t Vagriables ?-

Oversiding Concept is not applicable for vagrables.

Type. Gruntime object coon't to play any grole in vagrable gresdution?

Class P

Int x=888;
```

```
type. Druntime object won't to play any stole in vasicable session
       Class C extends P
           POL 2 =999;
         Class Test
           p.s. v.m (_____)
               P P = Dew PC);
             S.o.ph(p. x); // 888 -
              C canew C();
                 8.0.pln (c.x), /999~
                P P, = new C();
                   S.O.PID (P. 2); 888.
                                                                 0
                             do the static / both instancy one static & once instance
                                                                 0
                                                                 0
                                            °/p 888
                                                                 0
```

 \bigcirc

-> coheather The Vasicables are Static or non-static There is no Change in result.

difference blu Overloading & oversiding:

Paopearty	Overloading	Oven ou ding
10 method names	must be Same	must be same
) © asingwements	must be different (at least Oordern)	must be same (including onder).
® method Signature	must be different	must be same.
9 9 return type	no siestanchions	must be same until 1.47
		but from 1.50 onwasids Co-vasient Sietum types are
5) porvate, stake & final methods	Canbe overloaded	Can't be oversudden
) @ access modifiens	no déstructions	Scope we Can't decrease The Scope.
) Thorows Classe	no Diestructions	Size & level of checked
)))		Exceptions we Can't increase But we can decorease. But
) © method gresolution	-Always takes Gare by Compiles based on selemena type	No Diestructions for unchecked CAGptions. Always takes Gre by JVM based on Suntime Object
9 Also known as	Compile-time polymosiphism (03) Static polymosiphism (01) Easily binding	eruntime polymosphism (cs) dynamic polymosphism (r) Latebidius.

Note:

- → In Overloading we have to check only method names (must be Same) & assignments (must be diff.) All Gremaining teams like (Gretuan type, Throws Clause, Acressmodificans e.t.c.) asse not required to check.
- But in Oversiding coe have to check each & every Thing.
- O) Consider The following method cleclaration in parent class which of the following methods allowed in child class?

P: public void mi (int i) throws IDException

Owneridades public void mb (int i)

overloading public void mic) throws Exception

overlooding & public Static int mi (double d) throws I DException

C.E X @ public int mo (int i)

CE X & Public Synchronized void mi (int i) throws Exception

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overloading & public Static void mi (int... i) throws Exception

C.E x 9 public native abstract void m1(). throws Exception.

Polymosiphism

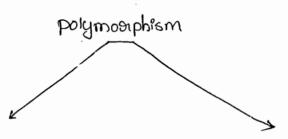
poly - many mosiphs means - Posims

> many froms ire polymosphism means

- Live Can use Same name to Shepsiesent multiple forms in Polymosiphism.

En oversiding we an have a method with one type of Amplementation in passent, but different type of implementation in child class.

-> Theore ane & types of polymosphism.



Compile-time polymosiphism

Run-tême polymorphism

Oveehoading

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Oversuiding.

Method Hiding

3 Pillages of Oops :-Znheoutana (Reuseability Oops

Encapsulation (Security)

Polymoorphism : (Plexibility)

funny diffination of polymosuphism:

A boy uses the wood FRIENDSHIP to Starts Love, but given uses the Same wood to gods. Same wood but different attitudes. This behaviour is nothing but polymorphism.

```
3 Static Conterol flow:
 En: -
                     Base
             Class
              Static int 2 = 10;
              (Static)
         (2)
                 m, (); ) 8
                (8.0.Pln (" FSB")) 10
               Public Static void main (Storing[] args)
         (3)
                 (m/U) 13
                                                                            3
                 (8.0.pln ("main method"); )15
                                                                            -
            (Public Static void mIL)
                                                                            )
                                                                           •
                (Soph (y);
                                                                            )
                                                     x = 0[ RI WO]
                                                                            )
                                                     Y=0[RIW0]
              Static
        (3)
                                                                            1
                                                      x=10[R&W]
                 S.o. Pln ( " SSB");
                                                                            0
                                                      y = 20[ R& W]
                                                                            Static int (y = 20;
                                                %!-
                                                       mainmetter.
                                                                            O
```

Paocess:

- then even we are trying to execute a Java class first that class file should be loaded, at the time class loading the following actions will be performed automatically.
- 1) identification of Static members forom Top to bottom. (1 to 6)
- (3 to 12)
- (3) Execution of main method. (13 to 15)

Read Indisactly write only state (RI WOS)

→ RF a Vasciable is in Read indisnectly would only state then we Can't perform head operation disnectly otherwise we will get Compile—time Enonon Saying "Illegat—forward—Reference."

 \Rightarrow

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..)

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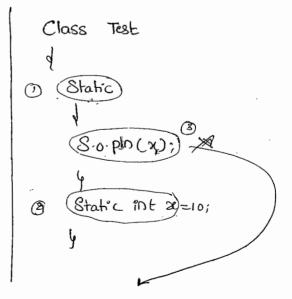
_)

:_)

()

()

()



C.E'- Illegal-forward sheftena,

Static block 8-

-> At the time of class loading if we want to perform any
activity we have to define that activity inside Static block because
Static blocks will be executed at the time of class loading.
- 2 cett en a class con con take and on a la all all the

70. of Gla Static blocks but all These Static blocks will be Executed forom top to bottom.

Ex(1): -

- After loading JDBC douber class we have to siegister doubler with deriver manager but every Driver class Contains a Static methology to perform this activity at the time of Dorver class loading automaticity) We ask not responsible to perform skegister Explicitly.

Class Derven Q. Register this Doubles with DM

Ex(0) & - Advantage:

Compulsary - At the time of class loading we have to Though The Coasesponding we an define this skep inside Static block.

۔ ِامِعَ Class Native Static System. load Library (native Library Path);

٠

.)

```
Static Control flow in parents child classes:
```

```
Class
         Base
   Static int @=10); 13
     Static
(D)
        M,(); (3)
        S-0-pln (" Base SB"); (S)
    Public Static void main ( )
      S.o.pln (" Base main");
    public static void mi()
        8-0-pln (y); (ii)
6 Static int (9=20) (6)
   Class Dealived extends Base
     Static int ()=100); (7)
6
      Static
3
         ma(); (B)
        8.0.pln (" DFSB'); @
    Public Static void main ( )
      $0-p("Desirved main"); &
```

 \Rightarrow

)

```
Public Static void ma()
 (9)
        8.0.pln(j); 1 20
                                                    2=0 [RIWO]
       Static
 (10)
                                                     Y=O [RIWO]
         Sop ( DSSB); (1)
                                                      1=0 [RIWO]
                                                      j=0 [RIWO]
       Static 12 (1) = 200)
  (1)
                                                       X =10 [R W]
                                                       y=20[R & w]
      > Java Derived
                                                        1 = 100 [R & W]
                                                                             •
     %P!-
             O
                                                        J = 200 [R & W]
                                                                             ď
             Base SB
             0
                                                    > Java Base
             DFSB
ossb
             200
                                                       Base SB
                                                        20
             Degrived main
                                                        Base main.
PROGRE:
         > Java C Desirved Java
                                                                            - )
                         Desived. class
    Base. class
                                                                            )
    > Java Desilved
    Identification of Static members from parent to child [1 to 1]
                                                                            ÷)
                                                                            )
     Execution & Static Vasilable assignments & Static blocks from
                                                                            .)
                                                                            •
    papent to child [12 to 22]
                                                                            )
                                                   [23 to 25]
*(3) Execution of only child class main method
                                                                            ()
                                                                            9
     (because mains) method of passent class is overeriding in child class, then, child
       -Class man() method executed)
                                                                           \bigcirc
```

Ponocess :-

- -> cohenever we are trying to load child class Then automatically parient class will be loaded to make parent class members available to the child class. Hence whenever we are executing child class. The following is the flow with respect to Static members step.
 - (1) Identification of Static members from parent to child
 - (9) Execution of Static Vasiable assignments & Static blocks from parent to child
 - (3) Execution of only child class main method. [IF the child class work Contain main method Then automatically parent class main() method will be executed]:

Note &

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)

When even we are loading child class automatically powent class will be loaded. But when even we are loading parent class child class wont be loaded.

Instance Control flow:-

```
Pagient
      class
       Pot (2)=10) 9
                                                 x=0 [RIWO]
                     6
         m; ();
                                                  4=0 [RIW]
        S.O.P ("FIIB"); (3)
                                                  2 = 10 [R w]
                                                   y = 20 [R W]
        Pagent U
    6
          S.o.pin(" Constructor"); (5)
    1) Public Static void main (Stochgest args)
        Pagent p = Dew Parent ();
        Sopla (" main');
        Public void mil)
    (6)
          S.o.phn (y); 1
                      instanceblock
    Ŧ
         S.o.pln(" SIIB"); (3)
         int y=20; @
    (8)
0/01_
         FIIB
        SJIB
         Copstructos
```

- Sequence of events will be performed automatically.
- (1) Identification of instance members forom top to bottom
 [1 to 8]
- (2) Execution of instance variable assignments & instance blocks from top to botom [9-14].
- (3) Execution of Construction [15]

* Note:

-> Static Control frow is only one time activity and it will be performed at the time of class loading But instance control flow is not one time activity for every object Coreation it will be executed.

Instance Control flow forom pasent to child:

Class Pament

3 int x=10; (5)

© m(1); © S.o. pln ("pasent"); ®

Sopho (" parent Constructori");

```
1) Public Static void main (-----)
                                                        parent
     Patient p=new parent();
                                                        Parent Constructor
(2)
     S-o-pin (" chita main");
 @ Public void mi()
      8-0-pln(y); (3)
 (f) (int y=20; (f)
    Class Child extends Pasent
   (int 1=)100;
  (11)
                                                        0
        m_2();
                                                        CIIB
        S.o.pln ("CIIB"); @
                                                       CSIIB
                                                        Child Constructor
        Child ()
  (12)
                                                        Child main.
          S.o.pln (" child Construction"); (8)
                                                                     )
       Public Static void main (_____)
                                                                     )
          child c = new child();
           8.0.pln (" child main"); 29
                                                                     \mathbf{O}
        public void ma()
                                                                     \bigcup
           8.0 pln (1); @
```

(int j=200; (3)

Porocess %-

- The following Sequence of execute events will be performed automatically.
-) (1) adentification of instance members from parent to child.
- (a) Execution of instance vasifable assignments & instance blocks
 ⇒ only in pasient class.
 - (3) Execution of Pasient class Constructor.
- Only in child class.
 - (3) Execution of child class Constructors.

>java parent

_)

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Constauctoors :-

Col.

```
-> Object Coreation is not enough Compulsary we should perform initialization then only that Object is in a position to perovide Diesponce peroperty.
```

```
The Code coil be executed automatically to perform initialization this place of Code is nothing but Construction. Here the main objective of Construction is to perform initialization
```

Cocchull.

```
for the newly Coneated Object.
```

```
Class Student

O Pot Siolino;
```

Student (Storing name, int shollno)

```
this . noting = name;
this . noting = rolling;
```

```
this . Tollno = rollno;

Public Static void main (Storing 17 args)
```

- At the time of object Caealion if we want to perform instalization of instance variable then we should go for Constructor.
- -> Other than initialization activity if we want to perform only activity at the time of Object Caeation then we Should go for instance block.
- → We Can't Preplace Constructions coith instance block because Constructor Can take agguement cohere as instance block Can't take agguements.
- Similarly we Can't Suplace instance block with Construction because a class Can Contain more than one Construction. If we want to deplace instance block with Construction then in every Construction we have to write instance block code because at suntime which Construction will be Called we can't Expect. It she suits deplicate & Coneater maintaince.

Class Test Ey! -•) **.**) Static int Count = 0;) ONA ove Test () Count ++1 if we (geste Test (int i) gratance Count ++) ;) p. S. v.m (---TOSE 6 = new TOSE(); if Test to = new Test (10);

 \Rightarrow

```
Rules to define Constauctors:
```

```
The name of the class & name of the Construction must be matched.
 2) Sieturn type Concept is not applicable for Constructor even void also.
   By Mistake if we declare return type for the Construction we wont
   get any Compiletime (as suplime Essoss, because Compiles treats
   96 as method.
        Class Test
   ,cXJ_
                               It is a noormal method but not Constructor
           void Test ()
  It is legal (for Stuppid to have a method whose name is exactly
    Same as class name).
(3) The only applicable modifiens for Constructions and
     Public, pouvate, porotected, colefault> [PPPD], if we asse toying
  to use any other modifies we will get Compile-time Essos Saying
                       is not allowed here.
     modifies
               XXXX
                    3 Static/final Staich-fp - --
       Class Test
 CXI.
          final Test()
                              mother
                                        Final
                                             is not allowed here
```

Singleton classes :-

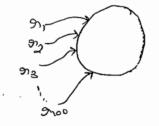
-) for any java class if we ask allowed to Coreate only one Object Such type of class is Called Singleton class

Runtime, Action Servick (Structs 1.x)

Business Deligate (EJB), Segvice Locaton (EJB) ---- e-t.c

-> The main advantage of Singleton is, instead to Coneating a Separate Object for every requirement we can create a Single Object and reuse The Same Object for every requirement this approach improves memory Matilization & performance of the System.

> Runtime on = / Runtime. get=Runtime () Runtime 912 = Runtime. gue Runtime() La Glass Static method



Runtime 2000 = Runtime. get Runtime()

Coneation of our own Singleton Class:

-> We can Create own own Singleton classes also for this we have to use paivate Construction & factory method.

Class Test Ep!-

parvate static Test t:

Paivate Test ()

Public Static Test get. Instance()

```
if (t = = Dull)
          t = New Test ();
        Detuan t;
    Public Object clone()
       Detuan this;
  4
    Test t, = Test.getInstana ();
    Test to = Test get Instance();
        tion = Test get Instance ();
    Test
         Go1 = Test. Clore();
tactory method:
  -> By using class name if we call any method & Gretusin
                                                                        9
   Same class Object. Then That method is Consider as factory
                                                                        )
                                                                        )
   method.
                                         -> factory method
  Ex! -
        Runtime on - Runtime, get-Runtime();
                                                                        )
        Date Format df = Date format get Instance ();
                                           I factory method.
                                                                        0
               t = Test. get Instance();
         Test
                                                                        U
                               - factory method
                                                                        \Theta
```

```
- Similarly we can Create Doubleton, Threableton ---- xxxx ton 96
 Classes.
-How to Coneate Doubleton class :-
Ex: -
         Class Test
           Porvate Static Test tis,
           parvate Static Test to;
           Paivate Test()
           Public Static Test get Instance ()
            if (t, == na1)
              t, = new Test();
             seturn ti;
            else
            if ( ( = = Dall)
              ty = New Test ();
              Detuoin to;
             4
           else
             if (math. enandom () < 0.5)
              Dietuoin ti;
             else
```

actuan to;

Rule 8-

Default	Constructor	Ÿ-

- -) If we are not writing any Construction then Compiler will always generate default Constauctor.
- -> 27 we are writing atleast one Constauctor Then Compiler won't generate default Construction.
- Henc a class Can Contain esther programmen whiten Constructor (69) Compiler generated Construction but not both Simultaniously.

Perototype of default Constructor 8-

- t) It is always no asycement Construction.
- a) The access modifier of default Construction is Same as class modifien but this sure is applicable public & < default>.
- 3) It Contains only one line, It is a no asycument Call to Super class Construction.

Super 12?

```
class Test
(1)
    public class Test
(3) Class Test
```

```
Class Test
U
       Testu
      Superici;
```

(3)

void test () > It is not a constructed mather It is a round mathed (3)

Class Test TESE () Super(); void Test()

public class Tests

public Test()

Superic);

Class Test (4) Test()

Class Test (y) Test() Superic);

Class Test T86() This (10): ... Test (int i). Class Test Test() This (10); Test (int i) Superic);

```
(6)
        Class Test
                                       Class Test
                                  6)
          Test (int i)
                                          Test (int i)
              Super ();
                                             Superic);
  Supear & This :-
-> The finish Line inside a Construction Should be either Supercial
  thise).
-> If we are not consisting any thing Compiler will always places super()
 Case (i):-
    We have to keep either Super ()(on this () only as the frost Line
  of the Construction.
          class Test
             TOSE ()
              S-o.p(" Hi");
                                      Call to Super must be first
              Superici; x
                                       Statement in Construction.
Case (ii):-
    with in the Construction we Can use either super() or this)
 but not both Simultaniously.
                                                                         ુ
      Class Test
        TESEL
                                                                         7
                                  Call to this must be frast statement
         this (); x -
                                   in the Constauctor
```

- [)

()

we can use Super & this only inside Construction of we are using any where else we will get compiletime estation.

Gol. Class Test

| public void mi()
| Supericly; X C.E! - Call to Superi multiple
| Supericly; X Frost Statement in the Constructor
| frost Statement in the Constructor Statement i

Superic) must be used only in Construction

Superic) os the fixe statement only

this() but not both symultaniously.

this():- To Call Cuspent class Constructions

Super():- To Call Parent Class Constructions

Compiler provides default super() but not this().

Superco	Super this
thisc	COS
(1) These agre Constauctor	D these agre key woods to greferse. Super & Chass instant members
(2) we should use only in Constauchous	2) We Can use any where Except in Static assea.

```
Ex:-
          Class Test
            p.s. v.mil)
             S. o.pln ( Super. hash Code ()); X
                         GE! NON-Static Vasicable Super Carit be
                                Dieferenced from a Static Context
Construction overloading :-
 A class Can Contain moone than one Construction with Same
name but with different arraquements & these Constructors are
 Considers as overloaded Constructors.
 en.
          Class Test
           Test (double d)
             this (10);
              S.o.pin (" double - asigs");
           Test (int i)
              this ();
              s.o.pln(" int-aggs");
          Test ()
          S.o.p (* No - asgs");
          P.S. v.m (____)
                                                                       0
```

Test E, = Dew Test (10.5); > No-0388 int - angs double - asys Test to = new Test (10) - No-angs Test to = new Test() - No-args

- -> Enhercitance & oversiding Concepts are not applicable 9091 Constructions.
- ** Every class in java including abstract class also can Contain Construction. But interfaces can't have the Constructions.

→ Case(<u>i)</u>:-

)

)

-> Recupisive method Call is always mustime Exception where as Diecusisive Construction innocation is a Compiletime Essassi.

ep!-Class Test. P.S. v. m1 () 4 m2 (); p. s. v. m2 () かいい m2() mill p.8. v.m(mac) MIC S.O.p(" Hello"); RE!. Stack Over Flow Examp class Test TESE () this (10) Test (Int 1) this (); p.s.v.m (---) S-o-pln (" Heno");

C.E !- Recursive Constructor innocation.

Case (ii) :-Cp. Class P Class P Class P P() P Cint i) Class C extends p Class C extends P Class C extends p C.E. Superco, Grit find Symbol Reprobbol: Constructor P() location: class P. Note: - if the passent class contains some Constructions then while woulting Child class use have to take special asse about Constauctions. → when ever we are coeffing any asignement Constructor it is highely • _) Decommended to work no arguement Construction also.) Caselin) :--) -> if parient class construction throws Some Checked Exception Compulsary Child class construction should throw Same Checked Exaption or its parent other wise the Code won't Compile.) class P class C extends p () **()** PC) throws IO Exaption C.F. unduported Exception Java. so. Idealeption in delaute Constant

By: - Class P

P() throws IOException

Class C extends p

CL) throws IOException Exception

y

Q) which of the following is Taue?

- @ Every class Contains Constructions
- and Concacate Classes Can Contains Constructions but not abstract Classes:
 - 3 the name of the Construction need not be same as class name X
 - 4 Preturn-type is applicable for the Construction X
-) & the only applicable modificas for Constructions are public & default X
-) © 8f we are trying to declare return type for the Constructor we could get Compiletime Estavor X.
-) @ Compiler will always generate default Construction X
 - 18 The acress modifien of the default construction is always default.
- (9) The first Line inside every Construction should be Super X.
- The fierst line " Should be Superior thisu,

if we are not worthing any thing compiler will always place thises X.

- (1) Enterface Can contains Construction
- 10 Both overloading & oversuiting Concepts are applicable for Constructor X.
- 3 Enhantana Concept is applicable for Constructor X

Type - Casting

Type-Castings.

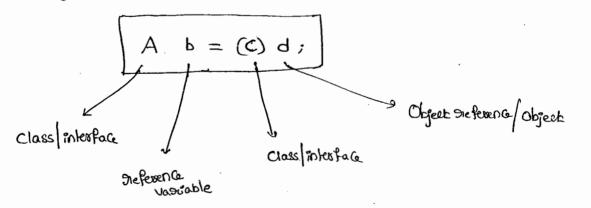
-> Pasient Class sieference Can be used to hold child class object

EDI- Pasient b = Dem CHIPO:

→ Similabilly, interface preference can be used to hold implemented class Object.

en, Runnable on = New Thosead();

Syntax:



Compiler onle (1):

C & type & d must have Some stellionship (either parent to)

Child Gos Child -> passent Gos Sometype) Other wise we will get Compiletime

Esonor Saying "inConvertable-types-found d type but sequired C type".

```
125 40
```

```
an:
              Object o = new String ("duago"),
              Storing Buffer Sb = (Storing Ruffer) 0;
              Storing 8 = New Storing ("dworga");
                       Sb = (SB)s;
              in Convextable types
               found : java-lang. Storing
               Dequired : Java-long - SB
   Compilerchecking sule 2 !-
        C must be either same on drawed type of A otherwise we
Cutill get Compiler time Estados Saying "incompatable types"
                                            found: C
                                             Irequired: A
- )
        GKili-
                    Object o = new Stocing ('doorga");
                     Storing S = (Storing) 0 ;
                  Storing & = new Storing ("dwonga");
                   Staing Buffer Sb = (Object)8;
                     in Compatable types
                    found: Object
                     Diguind: SB
```

.)

```
Runtime Checking
Rule 3:
```

-> The under leging object type of 'd' must be either same or désouved type of C. otherwise we will get suplime Expeption Saying " Class Cast Exception".

En: Object o = New Strong ("duarga");

SB sb=(SB)0; X

Rule 1 ~

3 × (R.E):- CCE

@ Object 0 = new String ("dwga");

Storing s = (Storing)0;

Rule 10 L

£81.

Basel

Der 2

Object

Dex3

Base2

Egyo Base a b = new Dery ();

(Object 0 = (Base 2) b;

x 3 Object 0 = (Baser) b;

(Base & b, = (Base 2) 0;

y & Basel b3 = (Dean) (New Derzen).

C: [: Incorrection ble types

tound : Bases

Dequired: Bases

(C.E.)inconvertiable type

found! Denz

Diequired : Den) \bigcirc

3

 Θ

-> Strictly Speaking in type-Casting just we are Converting only
type of object but not underlying object itself

Go!-

Ep1.

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.)

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```
हिन्दी :-
             A -- public Static void mic)
                      8.0.pln('A');
                     Public Static Void mic)
                       8.0.pln("B");
                    → public Static void mil)
                      8.0.pln(4c4);
     » C c = Dew C();
              c.m(U) // c
         ((B)c)·m,(); // B
     *)
          ((A)c). m, (); // A
Eg 3 %
            \rightarrow int x=777,
                                              c = Dew CC);
                                          8.0.pln(c.x); 999
             → int x=888;
                                          Sopho(((B)c) a); 888
                                                                            \Theta
                                                                            U
                                          & o. pin ((B)c)). x), 777
           + 10k &=999;
                                                                            \Theta
                                Checause The oversaiding Concept is not applicable
                                                                            0
                                                              for variable),
```

if we declare all Variables as Static Then There is no chance & change the O/p.

Note!

- wheather the vosciable & Static or instance vasciable sessition Should be done based on she ferrence type but not based on shunting - Object.

Coupling

Coupling :-

-> The degree of dependency blo The Components is Called Coupling

Ep! $\stackrel{\frown}{=}$

Class A

Static int i=B.j;

Static int j=C.m.();

P.S.v.int m.()

I shetwon D.k;

Class D Static int K=10;

-> The above Components are Said to be tightly Coupled with each other. Tightly Coupling is not Decommended because it has Several Servious disAdvantages.

(1) With out effecting semoining Component we Can't modify any Component

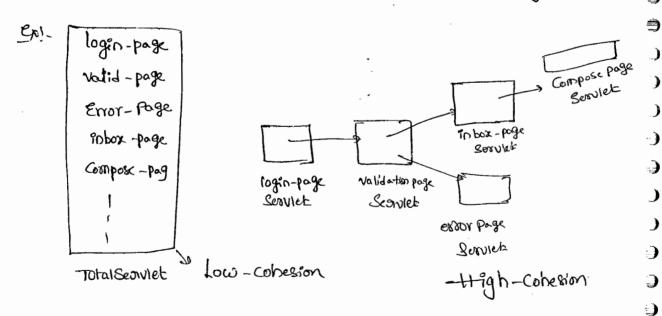
Hence, enhancement cutil become défficault.

- 2) it steduces maintainability.
- 3> ZE doesn't promote section lifty.
- Hence it is highly ene Commended to maintain loosely Coupling & dependency blu the Components should be as less as possible.

Cohesion

Cohesion :-

→ for every Component a clear well-defined functionality we have to define, Such type of Component is Said to be fallow High-cohesion



- → High-Cohesian is always a good paragraming practice which has Several Advantages.
 - (1) with out effecting semaining Components we can modify any Component Hence enchangement will become Very easy

7)

J

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- (2) ZL imposoves maintainability of the application
- (3) The posmotes sieuscability of the Code.

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-> cuhene even validation is nequired use Can ne use The Same Validate Servlet willhook negociting.

Note: -

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Loosely Coopling & high-Cohesian are good penagraming penactices.

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this:

to use the current class reference

means with out Creating multiple Objects, directorly on Object's Created Then
Can those values from the Current Class.

Kondalu-T@yahoo. Win

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     Collection forame work (1-43)
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- 1 Regular Exposessions (77-82)
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- 9 Development (96-101

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I would try to update our site JavaEra.com everyday with various interesting facts, scenarios and interview questions. Keep visiting regularly.....

Thanks and I wish all the readers all the best in the interviews.

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